CSE P 501 – Compilers

Parsing & Context-Free Grammars

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Administrivia

- Project partner signup: please find a partner and fill out the signup form by noon tomorrow if not done yet (only one form per group, please)
  - Watch for spam from CSE GitLab as repos are set up (save and ignore for now)
- Written HW2 out now, due in a week
- HW1 solution posted in a couple of days
- First part of project – scanner – out later this week, due in two weeks
  - Programming is fairly simple; this is the infrastructure shakedown cruise

Agenda for Today

- Parsing overview
- Context free grammars
- Ambiguous grammars
- Reading: Cooper & Torczon 3.1-3.2
  - Dragon book is also particularly strong on grammars and languages

Syntactic Analysis / Parsing

- Goal: Convert token stream to an abstract syntax tree
- Abstract syntax tree (AST):
  - Captures the structural features of the program
  - Primary data structure for next phases of compilation
- Plan
  - Study how context-free grammars specify syntax
  - Study algorithms for parsing and building ASTs

Context-free Grammars

- The syntax of most programming languages can be specified by a context-free grammar (CGF)
- Compromise between
  - REs: can't nest or specify recursive structure
  - General grammars: too powerful, undecidable
- Context-free grammars are a sweet spot
  - Powerful enough to describe nesting, recursion
  - Easy to parse; restrictions on general CFGs improve speed
- Not perfect
  - Cannot capture semantics, like "must declare every variable" or "must be `int`" – requires later semantic pass
  - Can be ambiguous (something we'll deal with)

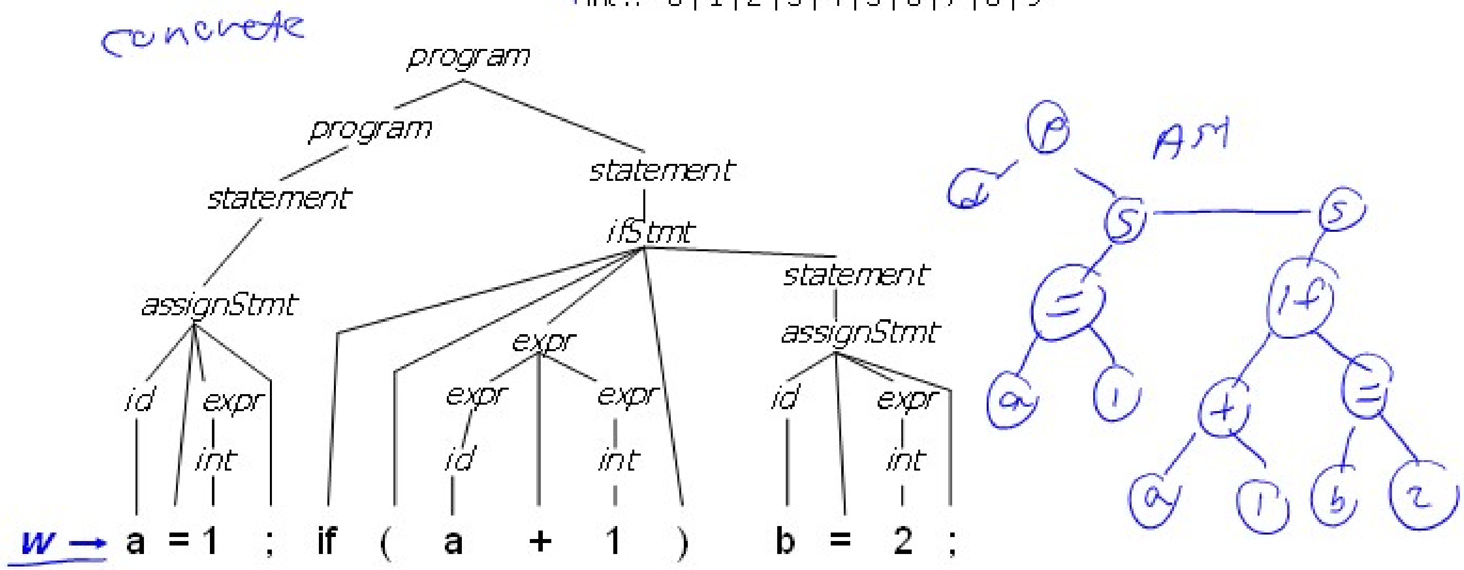
Derivations and Parse Trees

- Derivation: a sequence of expansion steps, beginning with a start symbol and leading to a sequence of terminals
- Parsing: inverse of derivation
  - Given a sequence of terminals (aka tokens) recover (discover) the nonterminals and structure, i.e., the parse tree (concrete syntax)

Old Example

G

program ::= *statement* | *program statement*
statement ::= *assignStmt* | *ifStmt*
assignStmt ::= *id* = *expr* ;
ifStmt ::= if (*expr*) *statement*
expr ::= *id* | *int* | *expr* + *expr*
id ::= a | b | c | i | j | k | n | x | y | z
int ::= 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9



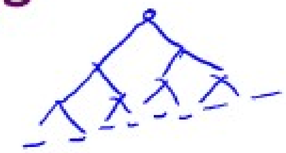
Parsing

- Parsing: Given a grammar *G* and a sentence *w* in *L(G)*, traverse the derivation (parse tree) for *w* in some *standard order* and do *something useful* at each node
  - The tree might not be produced explicitly, but the control flow of the parser will correspond to a traversal

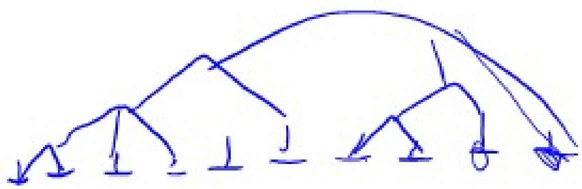
“Standard Order”

- For practical reasons we want the parser to be *deterministic* (no backtracking), and we want to examine the source program from *left to right*.
  - (i.e., parse the program in linear time in the order it appears in the source file)

Common Orderings



- Top-down
  - Start with the root
  - Traverse the parse tree depth-first, left-to-right (leftmost derivation)
  - LL(k), recursive-descent



- Bottom-up
  - Start at leaves and build up to the root
    - Effectively a rightmost derivation in reverse(!)
  - LR(k) and subsets (LALR(k), SLR(k), etc.)

“Something Useful”

- At each point (node) in the traversal, perform some semantic action
  - Construct nodes of full parse tree (rare)
  - Construct abstract syntax tree (AST) (common)
  - Construct linear, lower-level representation (often produced in later phases of production compilers by traversing initial AST)
  - Generate target code on the fly (done in 1-pass compilers; not common in production compilers)
    - Can’t generate great code in one pass, – but useful if you need a quick ‘n dirty working compiler

Context-Free Grammars

- Formally, a *grammar G* is a tuple $\langle N, \Sigma, P, S \rangle$ where
  - N is a finite set of *non-terminal* symbols
  - Σ is a finite set of *terminal* symbols (alphabet)
  - P is a finite set of *productions*
    - A subset of $N \times (N \cup \Sigma)^*$
  - S is the *start symbol*, a distinguished element of N
    - If not specified otherwise, this is usually assumed to be the non-terminal on the left of the first production

Standard Notations

- a, b, c elements of Σ
- w, x, y, z elements of Σ^*
- A, B, C elements of N
- X, Y, Z elements of $N \cup \Sigma$
- α, β, γ elements of $(N \cup \Sigma)^*$
- $A \to \alpha$ or $A ::= \alpha$ if $\langle A, \alpha \rangle \in P$

Derivation Relations (1)

- $\alpha A \gamma \Rightarrow \alpha \beta \gamma$ iff $A ::= \beta$ in P
  - derives
- $A \Rightarrow^* \alpha$ if there is a chain of productions starting with A that generates α
  - transitive closure

Derivation Relations (2)

- $w\,A\,\gamma \Rightarrow_{lm} w\,\beta\,\gamma$ iff $A ::= \beta$ in P
  - derives leftmost
- $\alpha\,A\,w \Rightarrow_{rm} \alpha\,\beta\,w$ iff $A ::= \beta$ in P
  - derives rightmost
- We will only be interested in leftmost and rightmost derivations – not random orderings

Languages

- For A in *N*, *L*(A) = { w | A =>* w }
- If *S* is the start symbol of grammar *G*, define *L*(*G*) = *L*(*S*)
  - Nonterminal on left of first rule is taken to be the start symbol if one is not specified explicitly

Reduced Grammars

X → abc
program → Foo
A → Abc
A → x
A → ε

- Grammar *G* is *reduced* iff for every production A ::= α in *G* there is a derivation
  - S =>* x A z => x α z =>* xyz
  - i.e., no production is useless
- Convention: we will use only reduced grammars
  - There are algorithms for pruning useless productions from grammars – see a formal language or compiler book for details

Ambiguity

- Grammar G is *unambiguous* iff every w in L(G) has a unique leftmost (or rightmost) derivation
  - Fact: unique leftmost or unique rightmost implies the other
- A grammar without this property is *ambiguous*
  - Note that other grammars that generate the same language may be unambiguous, i.e., ambiguity is a property of grammars, not languages
- We need unambiguous grammars for parsing

Example: Ambiguous Grammar for Arithmetic Expressions

expr ::= *expr* + *expr* | *expr* - *expr*
| *expr* * *expr* | *expr* / *expr* | *int*

int ::= 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9

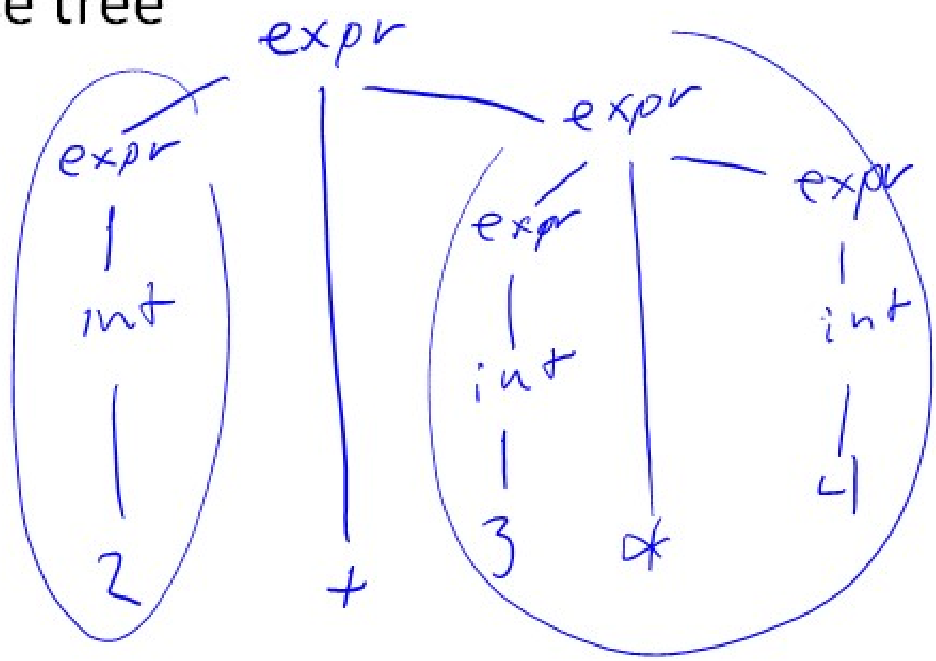
- Exercise: show that this is ambiguous
  - How? Show two different leftmost or rightmost derivations for the same string
  - Equivalently: show two different parse trees for the same string

Example (cont)

expr ::= *expr* + *expr* | *expr* - *expr*
| *expr* * *expr* | *expr* / *expr* | *int*

int ::= 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9

- Give a leftmost derivation of 2+3*4 and show the parse tree

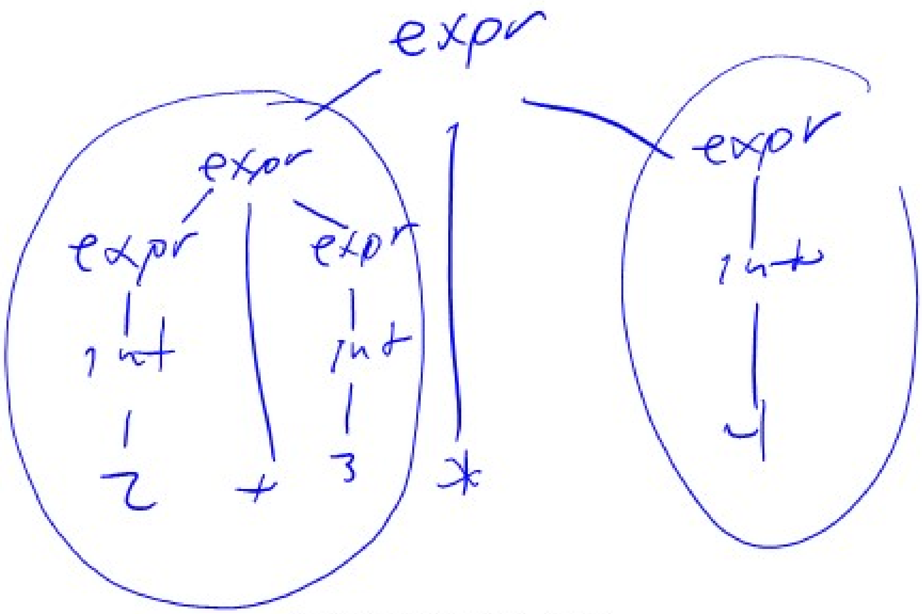


Example (cont)

expr ::= *expr* + *expr* | *expr* - *expr*
| *expr* * *expr* | *expr* / *expr* | *int*

int ::= 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9

- Give a different leftmost derivation of 2+3*4 and show the parse tree

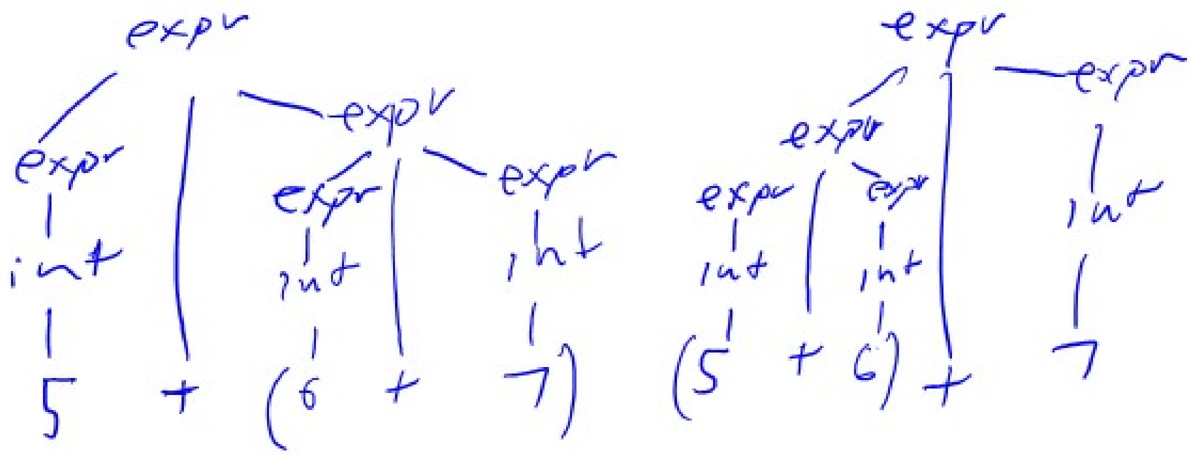


Another example

expr ::= *expr* + *expr* | *expr* - *expr*
| *expr* * *expr* | *expr* / *expr* | *int*

int ::= 0 | **1** | 2 | 3 | **4** | 5 | 6 | **7** | 8 | 9

- Give two different derivations of 5+6+7



(a - b) - c   a - (b - c)

What's going on here?

- The grammar has no notion of precedence or associatively
- Traditional solution
  - Create a non-terminal for each level of precedence
  - Isolate the corresponding part of the grammar
  - Force the parser to recognize higher precedence subexpressions first
  - Use left- or right-recursion for left- or right-associative operators (non-associative operators are not recursive)

Classic Expression Grammar

(first used in ALGOL 60)

expr ::= *expr* + *term* | *expr* – *term* | *term*

term ::= *term* * *factor* | *term* / *factor* | *factor*

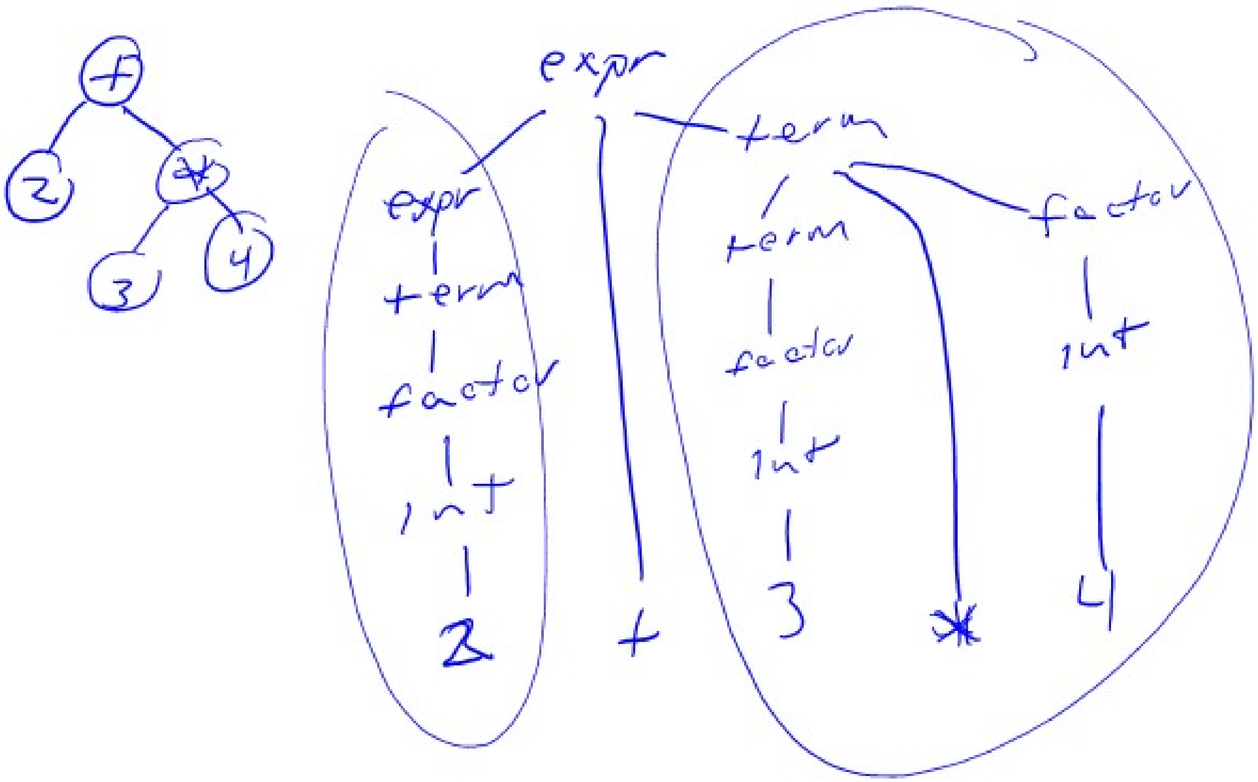
factor ::= *int* | (*expr*)

int ::= 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7

Check:

Derive 2 + 3 * 4

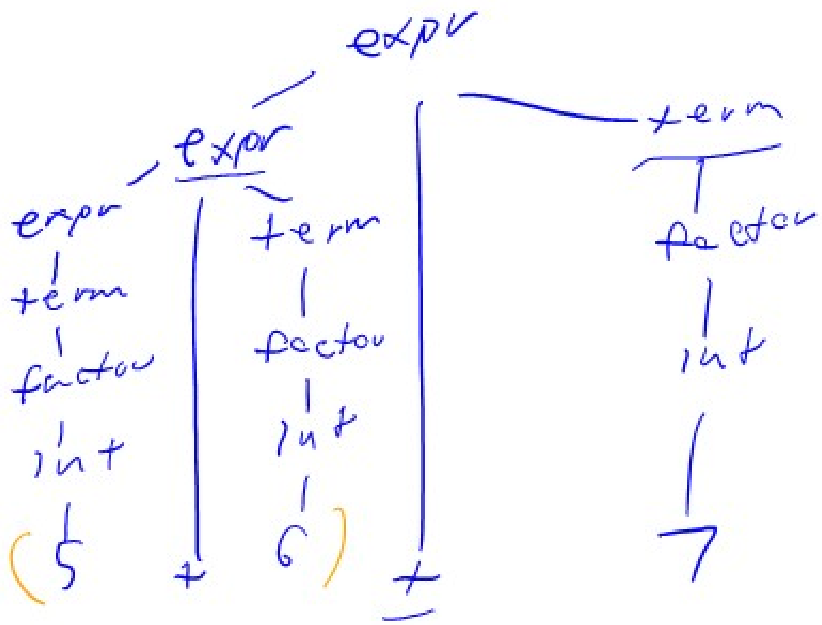
expr ::= *expr* + *term* | *expr* – *term* | *term*
term ::= *term* * *factor* | *term* / *factor* | *factor*
factor ::= *int* | (*expr*)
int ::= 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7



Check:

Derive 5 + 6 + 7

expr ::= *expr* + *term* | *expr* – *term* | *term*
term ::= *term* * *factor* | *term* / *factor* | *factor*
factor ::= *int* | (*expr*)
int ::= 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7



- Note interaction between left- vs right-recursive rules and resulting associativity

Check:

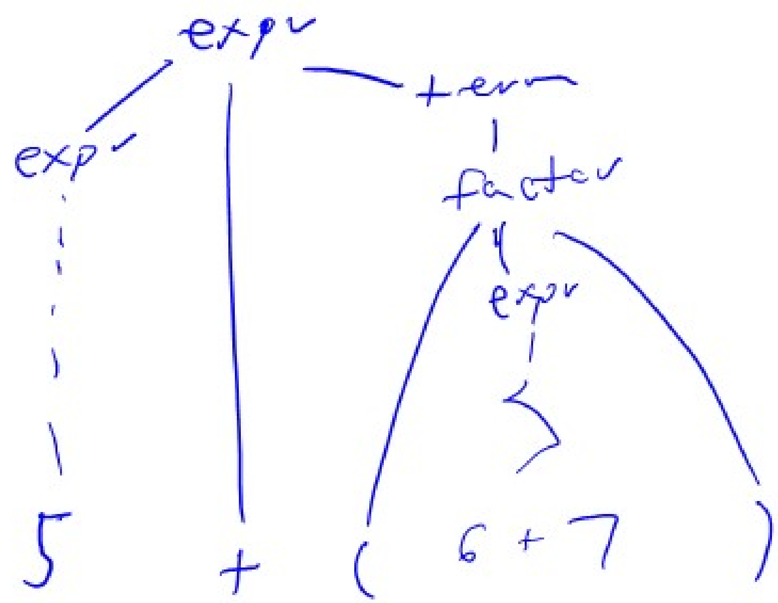
Derive 5 + (6 + 7)

expr ::= *expr* + *term* | *expr* – *term* | *term*

term ::= *term* * *factor* | *term* / *factor* | *factor*

factor ::= *int* | (*expr*)

int ::= 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7



Another Classic Example

- Grammar for conditional statements

  stmt ::= if (*expr*) *stmt*

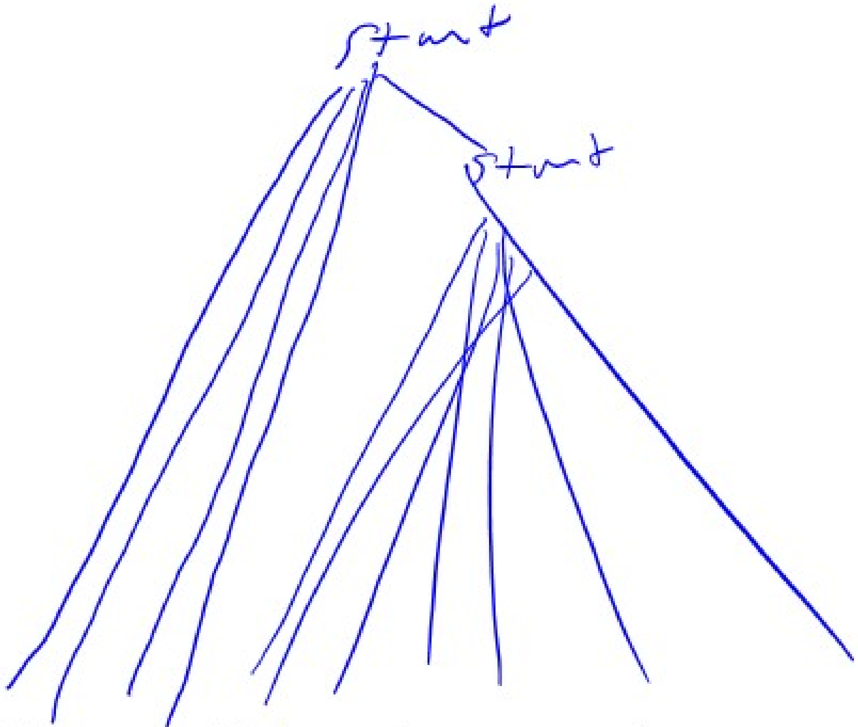
  | if (*expr*) *stmt* else *stmt*

  (This is the "dangling else" problem found in many, many grammars for languages beginning with Algol 60)

  – Exercise: show that this is ambiguous
    - How?

stmt ::= if (*expr*) *stmt*
| if (*expr*) *stmt* else *stmt*

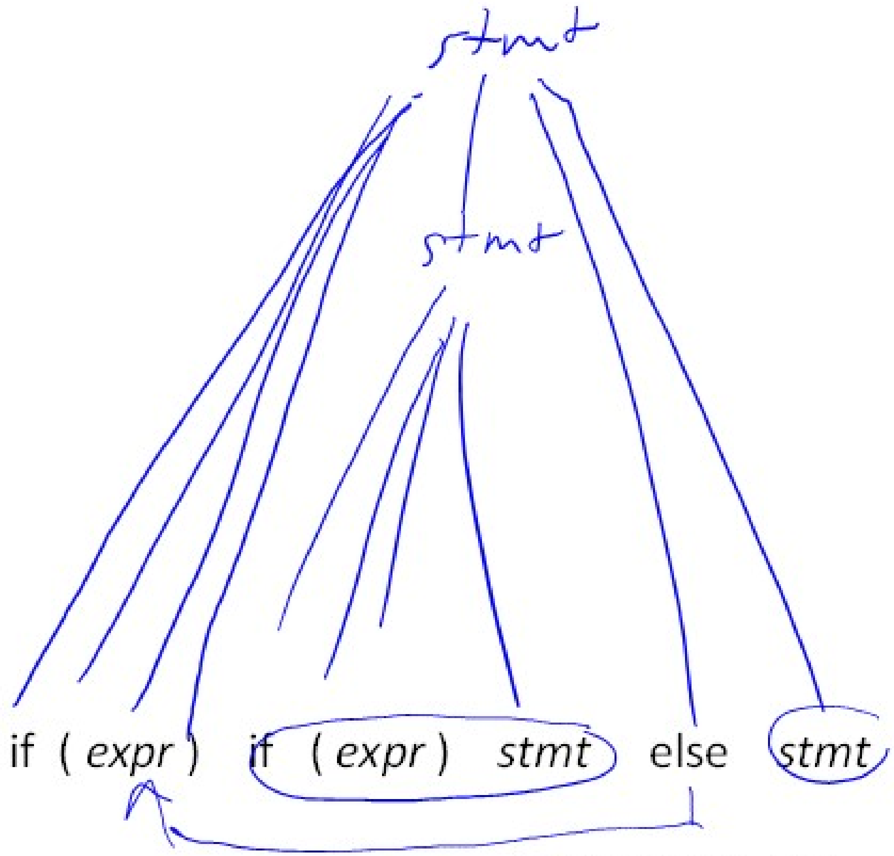
One Derivation



if (*expr*) if (*expr*) *stmt* else *stmt*

stmt ::= if (*expr*) *stmt*
| if (*expr*) *stmt* else *stmt*

Another Derivation



Solving "if" Ambiguity

- Fix the grammar to separate if statements with else clause and if statements with no else
  - Done in Java reference grammar
  - Adds lots of non-terminals
- or, Change the language
  - But it'd better be ok to do this – you need to "own" the language or get permission from owner
- or, Use some ad-hoc rule in the parser
  - "else matches closest unpaired if"

Resolving Ambiguity with Grammar (1)

Stmt ::= MatchedStmt | UnmatchedStmt

MatchedStmt ::= ... |

  if (Expr) MatchedStmt **else** MatchedStmt

UnmatchedStmt ::= ... |

  if (Expr) Stmt |

  if (Expr) MatchedStmt **else** UnmatchedStmt

– formal, no additional rules beyond syntax

– can be more obscure than original grammar

Check

Stmt ::= MatchedStmt | UnmatchedStmt
MatchedStmt ::= ... |
    if (Expr) MatchedStmt **else** MatchedStmt
UnmatchedStmt ::= **if** (Expr) Stmt |
    if (Expr) MatchedStmt **else** UnmatchedStmt

if (*expr*) if (*expr*) *stmt* else *stmt*

Resolving Ambiguity with Grammar (2)

- If you can (re-)design the language, just avoid the problem entirely

  Stmt ::= ... |
  if Expr **then** Stmt **end** |
  if Expr **then** Stmt **else** Stmt **end**

  – formal, clear, elegant
  – allows sequence of Stmts in then and else branches, no { , } needed
  – extra end required for every if
    (But maybe this is a good idea anyway?)

Parser Tools and Operators

- Most parser tools can cope with ambiguous grammars
  - Makes life simpler if used with discipline
- Usually can specify precedence & associativity
  - Allows simpler, ambiguous grammar with fewer nonterminals as basis for parser – let the tool handle the details (but only when it makes sense)
    - (i.e., expr ::= expr+expr | expr*expr | … with assoc. & precedence declarations can be the best solution)

Parser Tools and Ambiguous Grammars

$A \to \alpha \mid \beta$
$A \to \alpha$
$A \to \beta$

- Possible rules for resolving other problems:
  - Earlier productions in the grammar preferred to later ones (some danger here if grammar changes)
  - Longest match used if there is a choice (good solution for dangling if)
- Parser tools normally allow for this
  - But be sure that what the tool does is really what you want
    - And that it's part of the tool spec, so that v2 won't do something different (that you *don't* want!)

Coming Attractions

- Next topic: LR parsing
  - Continue reading ch. 3