

Chapel: Base Language

"Hello World" in Chapel: Two Versions

- Fast prototyping

```
writeln("Hello, world!");
```

- "Production-grade"

```
module Hello {
  proc main() {
    writeln("Hello, world!");
  }
}
```



Chapel Influences

C, Modula: basic syntax

ZPL, HPF: data parallelism, index sets, distributed arrays

CRAY MTA C/Fortran: task parallelism, synchronization

CLU (see also Ruby, Python, C#): iterators

Scala (see also ML, Matlab, Perl, Python, C#): type inference

Java, C#: OOP, type safety

C++: generic programming/templates
(but with a different syntax)



Outline

- **Introductory Notes**
- **Elementary Concepts**
 - Lexical structure
 - Types, variables, and constants
 - Operators and Assignments
 - Compound Statements
 - Input and output
- **Data Types and Control Flow**
- **Program Structure**



Lexical Structure

- Comments

```

/* standard
   C style
   multi-line */

// standard C++ style single-line

```

- Identifiers:

- Composed of A-Z, a-z, _, \$, 0-9
- Cannot start with 0-9
- Case-sensitive



Primitive Types

Type	Description	Default Value	Currently-Supported Bit Widths	Default Bit Width
bool	logical value	false	8, 16, 32, 64	impl. dep.
int	signed integer	0	8, 16, 32, 64	64
uint	unsigned integer	0	8, 16, 32, 64	64
real	real floating point	0.0	32, 64	64
imag	imaginary floating point	0.0i	32, 64	64
complex	complex floating points	0.0 + 0.0i	64, 128	128
string	character string	""	N/A	N/A

- Syntax

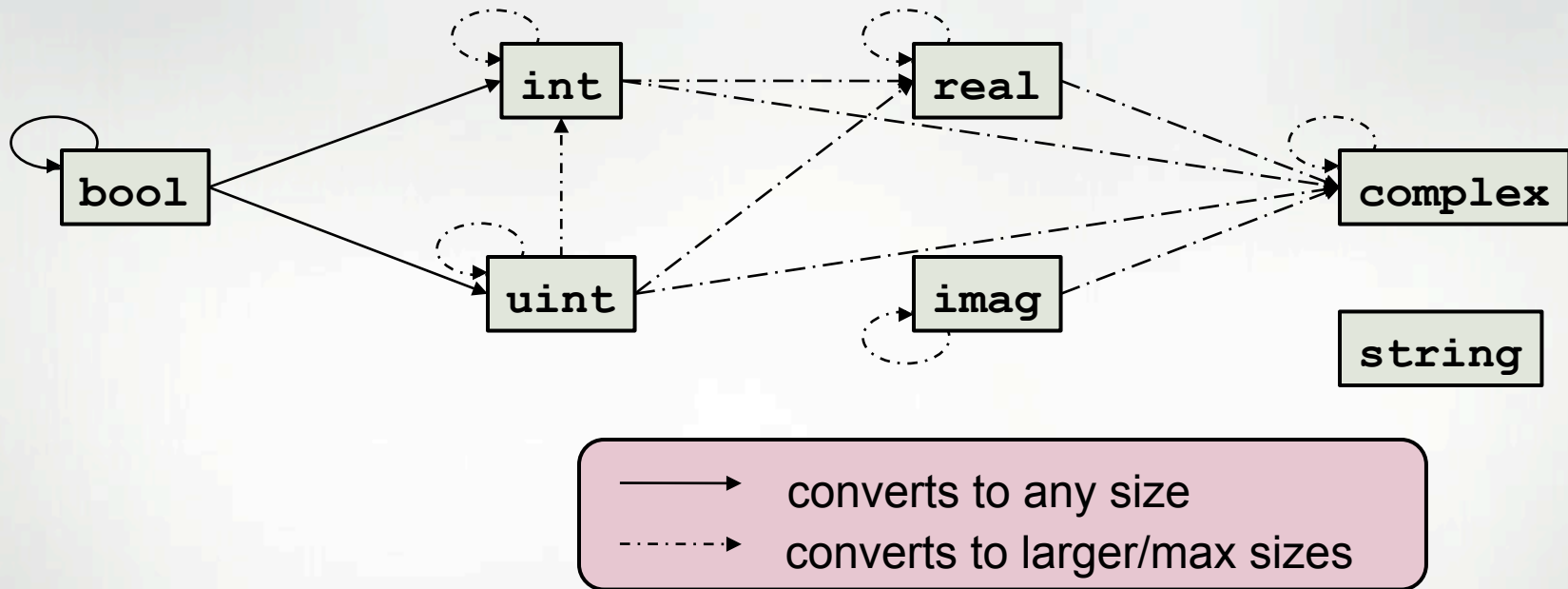
```
primitive-type:
  type-name [( bit-width )]
```

- Examples

```
int (16) // 16-bit int
real (32) // 32-bit real
uint // 64-bit uint
```



Implicit Type Conversions (Coercions)



- Notes:

- reals do not implicitly convert to ints as in C
- ints and uints don't interconvert as handily as in C

Type Aliases and Casts

- Basic Syntax

```

type-alias-declaration:
    type identifier = type-expr;

cast-expr:
    expr : type-expr
  
```

- Semantics

- type aliases are simply symbolic names for types
- casts are supported between any primitive types

- Examples

```

type elementType = complex(64);

5:int(8) // store value as int(8) rather than int
"54":int // convert string to an int
249:elementType // convert int to complex(64)
  
```



Variables, Constants, and Parameters

- Basic syntax

declaration:

```

var   identifier [: type] [= init-expr];
const identifier [: type] [= init-expr];
param identifier [: type] [= init-expr];

```

- Semantics

- **var/const**: execution-time variable/constant
- **param**: compile-time constant
- No *init-expr* \Rightarrow initial value is the type's default
- No *type* \Rightarrow type is taken from *init-expr*

- Examples

```

const pi: real = 3.14159;
var count: int;           // initialized to 0
param debug = true;      // inferred to be bool

```



Config Declarations

- Syntax

```
config-declaration:
  config type-alias-declaration
  config declaration
```

- Semantics

- Like normal, but supports command-line overrides
- Must be declared at module/file scope

- Examples

```
config param intSize = 32;
config type elementType = real(32);
config const epsilon = 0.01:elementType;
config var start = 1:int(intSize);
```

```
% chpl myProgram.chpl -sintSize=64 -selementType=real
% a.out --start=2 --epsilon=0.00001
```



Basic Operators and Precedence

Operator	Description	Associativity	Overloadable
:	cast	left	no
**	exponentiation	right	yes
! ~	logical and bitwise negation	right	yes
* / %	multiplication, division and modulus	left	yes
<i>unary</i> + -	positive identity and negation	right	yes
+ -	addition and subtraction	left	yes
<< >>	shift left and shift right	left	yes
<= >= < >	ordered comparison	left	yes
== !=	equality comparison	left	yes
&	bitwise/logical and	left	yes
^	bitwise/logical xor	left	yes
	bitwise/logical or	left	yes
&&	short-circuiting logical and	left	via isTrue
	short-circuiting logical or	left	via isTrue



Assignments

Kind	Description
=	simple assignment
+ = - = * = / = % = * * = & = = ^ = & & = = << = >> =	compound assignment (e.g., $x += y;$ is equivalent to $x = x + y;$)
< = >	swap assignment

- Note: assignments are only supported at the statement level



Compound Statements

- Syntax

```
compound-stmt:
  { stmt-list }
```

- Semantics

- As in C, permits a series of statements to be used in place of a single statement

- Example

```
{
  writeln("Starting a compound statement");
  x += 1;
  writeln("Ending the compound statement");
}
```



Console Input/Output

- Output

- `write(expr-list)`: writes the argument expressions
- `writeln(...)` variant: writes a linefeed after the arguments

- Input

- `read(expr-list)`: reads values into the argument expressions
- `read(type-list)`: reads values of given types, returns as tuple
- `readln(...)` variant: same, but reads through next linefeed

- Example:

```
var first, last: string;
write("what is your name? ");
read(first);
last = read(string);
writeln("Hi ", first, " ", last);
```

```
What is your name?
Chapel User
Hi Chapel User
```

- I/O to files and strings also supported



Outline

- Introductory Notes
- Elementary Concepts
- Data Types and Control Flow
 - Tuples
 - Ranges
 - Arrays
 - For loops
 - Other control flow
- Program Structure



Tuples

- Syntax

```
heterogeneous-tuple-type:
  ( type, type-list )
```

```
homogenous-tuple-type:
  param-int-expr * type
```

```
tuple-expr:
  ( expr, expr-list )
```

- Examples

```
var coord: (int, int, int) = (1, 2, 3);
var coordCopy: 3*int = coord;
var (i1, i2, i3) = coord;
var triple: (int, string, real) = (7, "eight", 9.0);
```

- Purpose

- supports lightweight grouping of values
(e.g., when passing or returning procedure arguments)
- multidimensional arrays use tuple indices



Range Values

- Syntax

```
range-expr:
  [low] .. [high]
```

- Semantics

- Regular sequence of integers

low <= *high*: *low*, *low*+1, *low*+2, ..., *high*

low > *high*: degenerate (an empty range)

low or *high* unspecified: unbounded in that direction

- Examples

```
1..6           // 1, 2, 3, 4, 5, 6
6..1           // empty
3..            // 3, 4, 5, 6, 7, ...
```



Range Operators

- Syntax

```

range-op-expr:
  range-expr by stride           range-expr align alignment
  range-expr # count            range-expr[range-expr]
  
```

- Semantics

- by**: strides range; negative *stride* \Rightarrow start from *high*
- #**: selects initial *count* elements of range
- align**: specifies the alignment of a strided range
- [] or (): intersects the two ranges

- Examples

```

1..6 by 2 // 1, 3, 5
1..6 by -1 // 6, 5, 4, ..., 1
1..6 #4 // 1, 2, 3, 4
1..6[3..] // 3, 4, 5, 6
  
```

```

1.. by 2 // 1, 3, 5, ...
1.. by 2 #3 // 1, 3, 5
1.. by 2 align 2 // 2, 4, ...
1.. #3 by 2 // 1, 3
0..#n // 0, ..., n-1
  
```



Array Types

- Syntax

```
array-type:
  [ index-set-expr ] elt-type
```

- Semantics

- Stores an element of *elt-type* for each index
- Array values expressed using square brackets

- Examples

```
var A: [1..3] int = [5, 3, 9], // 3-element array of ints
    B: [1..3, 1..5] real,      // 2D array of reals
    C: [1..3][1..5] real;     // array of arrays of reals
```

Much more on arrays in data parallelism section later...



For Loops

- Syntax

```
for-loop:
  for index-expr in iteratable-expr { stmt-list }
```

- Semantics

- Executes loop body serially, once per loop iteration
- Declares new variables for identifiers in *index-expr*
 - type and const-ness determined by *iteratable-expr*
 - *iteratable-expr* could be a range, array, or iterator

- Examples

```
var A: [1..3] string = [" DO", " RE", " MI"];

for i in 1..3 { write(A(i)); }           // DO RE MI
for a in A { a += "LA"; } write(A);     // DOLA RELA MILA
```



Zipper Iteration

- Syntax

```
zipper-for-loop:
  for index-expr in zip( iteratable-exprs ) { stmt-list }
```

- Semantics

- Zipper iteration is over all yielded indices pair-wise

- Example

```
var A: [0..9] real;

for (a,i,j) in zip(A, 1..10, 2..20 by 2) do
  a = j + i/10.0;

writeln(A);
```

```
2.1 4.2 6.3 8.4 10.5 12.6 14.7 16.8 18.9 21.0
```



Other Control Flow Statements

- Conditional statements

```
if cond { computeA(); } else { computeB(); }
```

- While loops

```
while cond {  
  compute();  
}
```

```
do {  
  compute();  
} while cond;
```

- Select statements

```
select key {  
  when value1 { compute1(); }  
  when value2 { compute2(); }  
  otherwise   { compute3(); }  
}
```

Note: Chapel also has expression-level conditionals and for loops



Control Flow: Braces vs. Keywords

Most control flow supports keyword-based forms for single-statement versions

- Conditional statements

```
if cond then computeA(); else computeB();
```

- While loops

```
while cond do  
  compute();
```

- For loops

```
for indices in iteratable-expr do  
  compute();
```

- Select statements

```
select key {  
  when value1 do compute1();  
  when value2 do compute2();  
  otherwise do compute3();  
}
```



Outline

- Introductory Notes
- Elementary Concepts
- Data Types and Control Flow
- Program Structure
 - Procedures and iterators
 - Modules and main()
 - Records and classes
 - Generics
 - Other basic language features



Procedures, by example

- Example to compute the area of a circle

```
proc area(radius: real): real {
    return 3.14 * radius**2;
}
```

```
writeln(area(2.0)); // 12.56
```

```
proc area(radius) {
    return 3.14 * radius**2;
}
```

Argument and return types can be omitted

- Example of argument default values, naming

```
proc writeCoord(x: real = 0.0, y: real = 0.0) {
    writeln((x,y));
}
```

```
writeCoord(2.0); // (2.0, 0.0)
```

```
writeCoord(y=2.0); // (0.0, 2.0)
```

```
writeCoord(y=2.0, 3.0); // (3.0, 2.0)
```



Iterators

- **Iterator:** a procedure that generates values/variables
 - Used to drive loops or populate data structures
 - Like a procedure, but yields values back to invocation site
 - Control flow logically continues from that point
- **Example**

```

iter fibonacci(n) {
  var current = 0,
      next = 1;
  for 1..n {
    yield current;
    current += next;
    current <=> next;
  }
}

```

```

for f in fibonacci(7) do
  writeln(f);

```

```

0
1
1
2
3
5
8

```



Argument and Return Intents

- Arguments can optionally be given intents
 - (blank): varies with type; follows principle of least surprise
 - most types: **const**
 - arrays, domains, sync vars: passed by reference
 - **const**: disallows modification of the formal
 - **in**: copies actual into formal at start; permits modifications
 - **out**: copies formal into actual at procedure return
 - **inout**: does both of the above
 - **ref**: pass by reference
 - **param/type**: formal must be a param/type (evaluated at compile-time)
- Return types can also have intents
 - (blank)/**const**: cannot be modified (without assigning to a variable)
 - **var**: permits modification back at the callsite
 - **type**: returns a type (evaluted at compile-time)
 - **param**: returns a param value (evaluated at compile-time)



Modules

- **Syntax**

```

module-def:
  module identifier { code }

module-use:
  use module-identifier;
  
```

- **Semantics**

- all Chapel code is stored in modules
- `use`-ing a module makes its symbols visible in that scope
- module-level statements are executed at program startup
 - typically used to initialize the module
- for convenience, a file containing code outside of a module declaration creates a module with the file's name



Program Entry Point: main()

- Semantics
 - Chapel programs start by:
 - initializing all modules
 - executing main(), if it exists

```
M1.chpl:
use M2;
writeln("Initializing M1");
proc main() { writeln("Running M1"); }
```

```
M2.chpl:
module M2 {
  writeln("Initializing M2");
}
```

```
% chpl M1.chpl M2.chpl

% ./a.out

Initializing M2
Initializing M1
Running M1
```



Revisiting "Hello World"

- Fast prototyping

`hello.chpl`

```
writeln("Hello, world!");
```

`==`

```
module hello {
  writeln("Hello, world!");
}
```

- "Production-grade"

```
module HelloWorld {
  proc main() {
    writeln("Hello, world!");
  }
}
```

Module-level code is executed during module initialization

main() executed when program begins running

Records and Classes

- Chapel's struct/object types
 - Contain variable definitions (fields)
 - Contain procedure & iterator definitions (methods)
 - Records: value-based (*e.g.*, assignment copies fields)
 - Classes: reference-based (*e.g.*, assignment aliases object)
 - Record : Class :: C++ struct : Java class
- Example

```

record circle {
  var radius: real;
  proc area() {
    return pi*radius**2;
  }
}

```

```

var c1, c2: circle;
c1 = new c1(radius=1.0);
c2 = c1;           // copies c1
c1.radius = 5.0;
writeln(c2.radius); // 1.0
// records deleted by compiler

```



Records and Classes

- Chapel's struct/object types
 - Contain variable definitions (fields)
 - Contain procedure & iterator definitions (methods)
 - Records: value-based (*e.g.*, assignment copies fields)
 - Classes: reference-based (*e.g.*, assignment aliases object)
 - Record : Class :: C++ struct : Java class
- Example

```

class circle {
  var radius: real;
  proc area() {
    return pi*radius**2;
  }
}

```

```

var c1, c2: circle;
c1 = new c1(radius=1.0);
c2 = c1; // aliases c1's circle
c1.radius = 5.0;
writeln(c2.radius); // 5.0
delete c1; // users delete classes

```



Method Examples

Methods without arguments need not use parenthesis

```

proc circle.circumference {
    return 2* pi * radius;
}

writeln(c1.area(), " ", c1.circumference);

```

Methods can be defined for any type

```

proc int.square() {
    return this**2;
}

writeln(5.square());

```



Generic Procedures

Generic procedures can be defined using type and param arguments:

```
proc foo(type t, x: t) { ... }
proc bar(param bitWidth, x: int(bitWidth)) { ... }
```

Or by simply omitting an argument type (or type part):

```
proc goo(x, y) { ... }
proc sort(A: []) { ... }
```

Generic procedures are instantiated for each unique argument signature:

```
foo(int, 3);           // creates foo(x:int)
foo(string, "hi");    // creates foo(x:string)
goo(4, 2.2);          // creates goo(x:int, y:real)
```



Generic Objects

Generic objects can be defined using type and param fields:

```
class Table { param size: int; var data: size*int; }
class Matrix { type eltType; ... }
```

Or by simply eliding a field type (or type part):

```
record Triple { var x, y, z; }
```

Generic objects are instantiated for each unique type signature:

```
// instantiates Table, storing data as a 10-tuple
var myT: Table(10);
// instantiates Triple as x:int, y:int, z:real
var my3: Triple(int, int, real) = new Triple(1, 2, 3.0);
```



Other Base Language Features not covered today

- Enumerated types
- Unions
- Type select statements, argument type queries
- Parenthesis-less functions/methods
- Procedure dispatch constraints (“where” clauses)
- Compile-time features for meta-programming
 - type/param procedures
 - folded conditionals
 - unrolled for loops
 - user-defined compile-time warnings and errors



Status: Base Language Features

- Most features working well
- Performance is currently suboptimal in some cases
- Some semantic checks are incomplete
 - e.g., constness-checking for members, arrays
- Error messages could use improvement at times
- OOP features are limited in certain respects
 - generic classes w/ subclassing, user constructors
- Memory for strings is currently leaked



Future Directions

- Error handling/Exceptions
- Fixed-length strings
- Interfaces (joint work with CU Boulder)
- Improved namespace control
 - private fields/methods in classes and records
 - module symbol privacy, filtering, renaming
- Interoperability with other languages (joint with LLNL)

