

# Meta Data Management

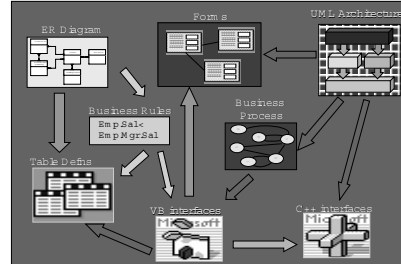
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Microsoft Research

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# Meta Data Management

- 1 Meta data = structural information  
DB schema, interface defn, web site map, form defns, ...



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# Meta Data Problems

- 1 Many data management applications primarily involve transformations of structured data
- 1 Data translation
- 1 Schema evolution
- 1 XML message translation
- 1 Application integration
- 1 Data warehouse loading
- 1 ER/UML design tools
- 1 Wrapper generation for SQL
- 1 UI/4GL generation
- 1 Dependency tracking
- 1 Lineage tracing
- 1 Info resource mgmt
- 1 Binding, renaming
- 1 Software build (make)
- 1 Configuration mgmt

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# Outline

- 1 Introduction
- 1 Meta data problems
- 1 Design patterns
- 1 Solution templates
- 1 Wrap up

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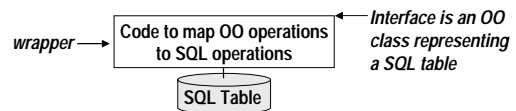
# Why Meta Data is Important

- 1 Many DB problems are easier to solve by manipulating meta data  
Instead of writing code  
Instead of manipulating data directly
- 1 Meta-data-based solutions all involve models (schemas) and mappings  
Mappings - data transformations, queries, dependencies, ...  
Model, manipulate, and generate them  
Usually, generate code from them

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# Example: Object-Oriented Wrapper for SQL Tables

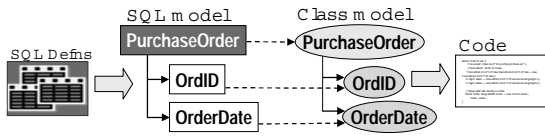


- 1 Manually program a wrapper for each table
- 1 This is very repetitive work
- 1 So you write a program to generate a wrapper for each table

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## OO wrapper for SQL (cont'd)



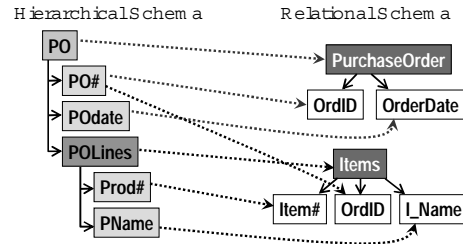
- The wrapper generator does the following:
  - Imports each table definition into a model
  - Generates a model for the class wrapper
  - Generates a mapping from table to class
  - Generates code from the class model and mapping.

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## Example - Data Translation

- Translate data from one data model to another
- Either write a program or generate it



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## Meta-data-Speak

Meta-data-Speak	English Example
meta-meta-model = meta-meta-meta data	Built-in types (usually hard-coded)
meta-model = meta-meta data	Schema for "Table," "Column," "Key," ...
model = meta data	Schema for the Employee Table
data	Employee Table

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## Outline

- Introduction
- Meta data problems
- Design patterns
- Solution templates
- Wrap up

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## Meta Data Solution Template

- Get a data manager form models and mappings

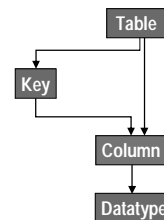
- Usually, it's an object manager
  - OO programming language
  - OODB
- Hence, meta-meta-model is the object manager's built-in types
  - Classes, attributes, methods, objects
  - Plus operators to manipulate them, such as New Class, New Attribute, New Object, Write Attribute

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## Meta Data Solution Template

- Get a data manager form models and mappings
- Design meta-model(s) (e.g., for SQL schemas)



If the meta-meta-model is OO, then the meta-model consists of class definitions

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## Meta Data Solution Template

- 1 Get a data manager from models and mappings
- 2 Design meta model(s) (e.g., for SQL schemas)
- 3 Build a model in porter for each meta model

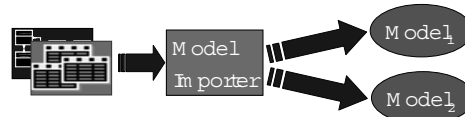


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## Meta Data Solution Template

- 1 Get a data manager from models and mappings
- 2 Design meta model(s) (e.g., for SQL schemas)
- 3 Build a model in porter for each meta model
- 4 Invoke model in porter(s)



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## Meta Data Solution Template

- 1 Get a data manager from models and mappings
- 2 Design meta model(s) (e.g., for SQL schemas)
- 3 Build a model in porter for each meta model
- 4 Invoke model in porter(s)

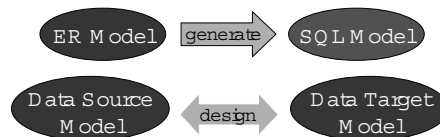
Problem	Model <sub>1</sub>	Model <sub>2</sub>
Data translation	source schema	target schema
Msg translation	source format	target format
App integration	source interfaces	target interfaces
DW loading	source schema	DW schema

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## Meta Data Solution Template

- 1 Get a data manager from models and mappings
- 2 Design meta model(s) (e.g., for SQL schemas)
- 3 Build a model in porter for each meta model
- 4 Invoke model in porter(s)
- 5 Generate or design mappings



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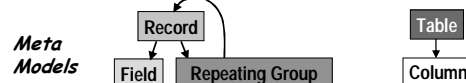
## Meta Data Solution Template

- 1 Get a data manager from models and mappings
- 2 Design meta model(s) (e.g., for SQL schemas)
- 3 Build a model in porter for each meta model
- 4 Invoke model in porter(s)
- 5 Generate or design mappings
- 6 Generate code: data /msg translation script, app wrapper, ETL script, view defn's, etc.

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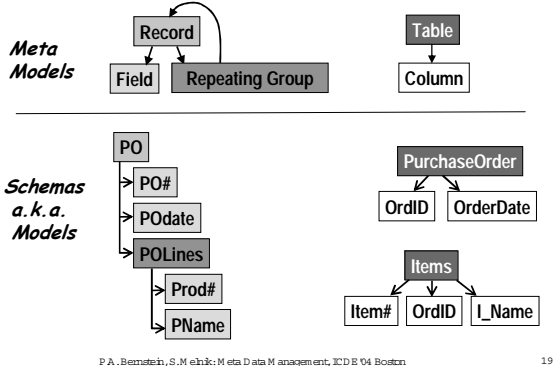
## Example - Data Translation



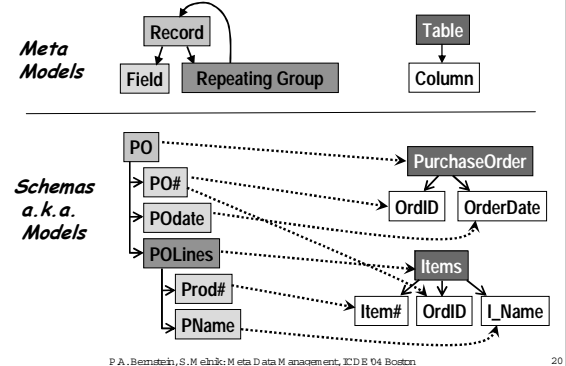
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## Example - Data Translation



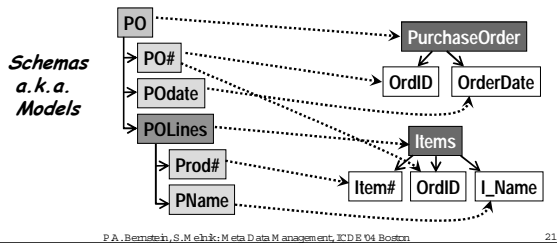
## Example - Data Translation



### Generated data translation script

```

Foreach [po#, poD, poL] in PO
  Insert [po#, poD] into PurchaseOrder
  Foreach [prod#, pN] in poL
    Insert [prod#, po#, pN] into Items
  End
End
  
```



## Meta Data Problems

- 1 Data translation
- 1 OO or XML wrapper generation for SQL DB
- 1 User-Interface / 4GL-program generation
- 1 Design tool support (DB, UML, ... )
  - Model generation, reverse engineering
  - Round-trip engineering
- 1 Schema evolution (applies to all scenarios)
- 1 XML message translation for e-commerce
- 1 Integrate custom apps with commercial apps

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## Meta Data Problems (cont'd)

- 1 Data warehouse loading (clean & transform)
- 1 Lineage tracing (provenance)
- 1 Information resource management
- 1 Dependency tracking
  - Impact analysis
  - Navigation between tools
- 1 Binding, renaming
- 1 Software build (make)
- 1 Version and configuration management
  - Release management
  - Product data management

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## Meta Data Solutions

- 1 They strongly resemble one another
- 1 We characterize that resemblance
  - Prototypical problems, or design patterns
  - Solution specifications, or solution templates
  - Primitive solution steps, or operators
- 1 Goals
  - A methodology to solve meta data problems
  - Ultimately, operator implementations to turn solution templates into solution programs

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## Outline

- 1 Introduction
- 1 Meta data problems
- ➔ Design patterns
- 1 Solution templates
- 1 Wrap up

## Meta Data Design Patterns

- 1 Design pattern - a problem description consisting of
  - Input models and mappings
  - Output models and mappings
  - Criteria for the output to be correctAn application specializes it to meta models and mapping languages
- 1 Solution template - a sequence of operators producing the desired output
- 1 Operators - a single step that computes a model and/or mappings

## Operators

- 1  $m_{ap} = Match(M_1, M_2)$   
Return a mapping between the two models
- 1  $\langle M_2, m_{ap_{12}} \rangle = ModelGen(M_1, meta\ model_2)$   
Return a model  $M_2$  that is expressed in meta model<sub>2</sub> and is equivalent to model $M_1$
- 1  $\langle M_3, m_{ap_{13}}, m_{ap_{23}} \rangle = Merge(M_1, M_2, m_{ap})$   
Return the union of models  $M_1$  and  $M_2$
- 1  $m_{ap_3} = Compose(m_{ap_1}, m_{ap_2}) = m_{ap_1} \circ m_{ap_2}$   
Return the composition of  $m_{ap_1}$  and  $m_{ap_2}$ , which is a mapping from  $m_{ap_1}$ 's domain to  $m_{ap_2}$ 's range.

## Operators (cont'd)

- 1  $m_{ap_3} = Confluence(m_{ap_1}, m_{ap_2}) = m_{ap_1} \cap m_{ap_2}$   
Return the "merge" of mappings  $m_{ap_1}$  and  $m_{ap_2}$
- 1  $\langle M_2, m_{ap_{12}} \rangle = Extract(M_1, m_{ap})$   
Return the sub-model of  $M_1$  that participates in the mapping  $m_{ap}$
- 1  $\langle M_2, m_{ap_{12}} \rangle = Diff(M_1, m_{ap})$   
Return the sub-model of  $M_1$  that does not participate in the mapping  $m_{ap}$

## Design Patterns

- 1 Meta Modeling
  - 1 Model Mapping
  - 1 Model Generation
  - 1 Model Integration
  - 1 Mapping Composition
  - 1 Mapping Alignment
  - 1 Change Propagation
  - 1 Model Reintegration
- } Single Operator Solutions

## Meta Modeling

- 1 Design pattern - develop a representation (i.e. meta model) for models and mappings
- 1 Applications - they all depend on this
- 1 Solution template
  - Design a meta model
  - Write Import & Export functions
    - Import SQL, Import XSD, Import ERD, ...
  - Today, it is manual engineering design
  - Design once and reuse often

## Meta Modeling (cont'd)

- The Import function for models
  - Parse text
  - Copy elements of the parsed form into a model that conforms to its meta model
- The Import function for mappings
  - Same as models but may require more semantic analysis
  - E.g., program dependencies, data lineage
  - For some languages and mapping meta models, Export is hard (e.g., XSLT)

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## Model Mapping

- Design pattern - Design a mapping between two models and generate code from it

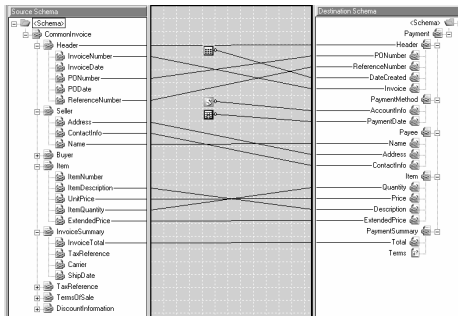


- Applications
  - Data translation
  - XML message translation fore-commerce
  - Integrate custom and commercial apps
  - Data warehouse extract, transform & load
- Solution templates
  - map = Match ( $M_1, M_2$ ); Export (map)
  - Mapping reuse: Compose, Confluence

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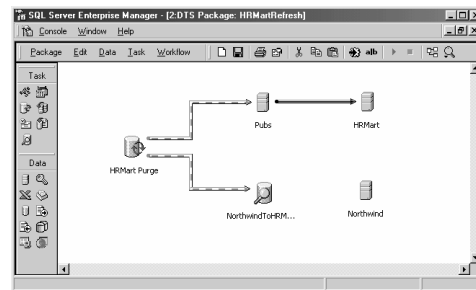
## An XML Mapping Tool



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## A Data Warehouse Loading Tool



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## Model Generation

- Design pattern - Given a model, generate an equivalent model in another meta model



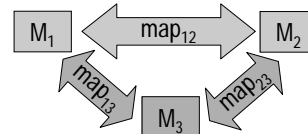
- Applications
  - Wrapper generation (SQL file or XML)
  - Design tools (ER file SQL, SQL file ER)
  - UI/4GL generation
- Solution template
  - $\langle M_2, \text{map} \rangle = \text{ModelGen}(M_1, \text{meta model}_2)$ ; Export( $M_2$ )
  - ModelGen often needs human guidance

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## Model Integration

- Design pattern - Given two models, develop a model that subsumes both of them



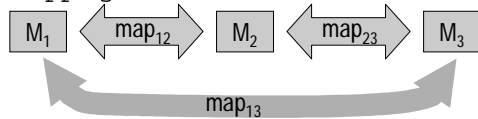
- Applications
  - View integration
  - Data integration
- Solution template
  - $\langle M_3, \text{map}_{13}, \text{map}_{23} \rangle = \text{Merge}(M_1, M_2, \text{map})$ ; Export( $M_3, \text{map}_{13}, \text{map}_{23}$ )

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## Mapping Composition

- Design pattern - Compose two given mappings



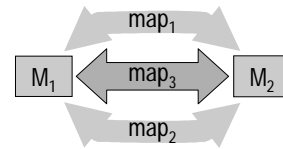
- Applications
  - Processing queries on views
- Solution templates
  - $map_{13} = \text{Compose}(map_{12}, map_{23})$
  - Answering queries using views (LaV), Query modification (GaV), GLaV

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## Mapping Alignment

- Design pattern - Align two mappings between the same pair of models



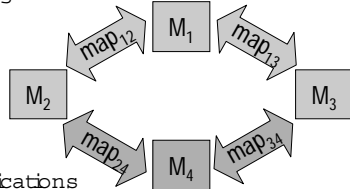
- Applications
  - P2P query processing, mapping design
- Solution template
  - $map_3 = \text{Confluence}(map_1, map_2)$

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## Model Reintegration

- Design pattern - Given a model and mappings to two modified versions of the model, produce a merged model



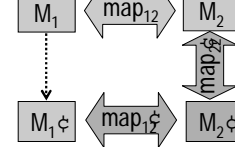
- Applications
  - Parallel development
- Solution template
  - Multistep application of many operators

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## Change Propagation

- Design pattern - Given two models and a mapping. One model changes. Fix the mapping and other model.



- Applications
  - Schema evolution, interface evolution, ...
  - Required maintenance for all meta data problems
- Solution template
  - Requires all of the operators

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## Outline

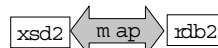
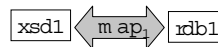
- Introduction
- Meta data problems
- Design patterns
- ➔ Solution templates
  - Change propagation
  - Model reintegration
  - Change propagation revisited
- Research background
- Wrap up

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## Change Propagation

- Given
  - $map_1$  between xsd1 and SQL schema rdb1
  - xsd2, a modified version of xsd1
- Produce
  - rdb2 to store instances of xsd2
  - a mapping between xsd2 and rdb2

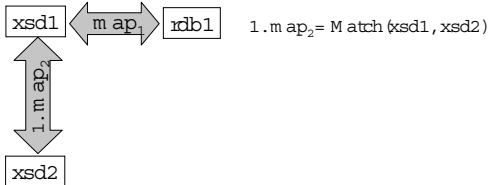


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## Change Propagation

- Given
  - $m_{ap_1}$  between  $xsd1$  and SQL schema  $rdb1$
  - $xsd2$ , a modified version of  $xsd1$
- Produce
  - $rdb2$  to store instances of  $xsd2$
  - a mapping between  $xsd2$  and  $rdb2$



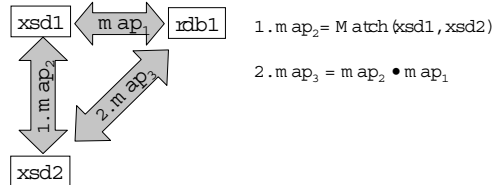
1.  $m_{ap_2} = Match(xsd1, xsd2)$

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## Change Propagation

- Given
  - $m_{ap_1}$  between  $xsd1$  and SQL schema  $rdb1$
  - $xsd2$ , a modified version of  $xsd1$
- Produce
  - $rdb2$  to store instances of  $xsd2$
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1.  $m_{ap_2} = Match(xsd1, xsd2)$

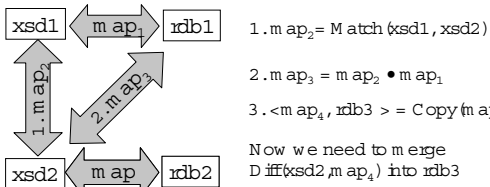
2.  $m_{ap_3} = m_{ap_2} \bullet m_{ap_1}$

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## Change Propagation

- Given
  - $m_{ap_1}$  between  $xsd1$  and SQL schema  $rdb1$
  - $xsd2$ , a modified version of  $xsd1$
- Produce
  - $rdb2$  to store instances of  $xsd2$
  - a mapping between  $xsd2$  and  $rdb2$



1.  $m_{ap_2} = Match(xsd1, xsd2)$

2.  $m_{ap_3} = m_{ap_2} \bullet m_{ap_1}$

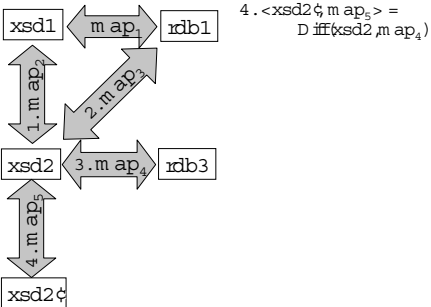
3.  $\langle m_{ap_4}, rdb3 \rangle = Copy(m_{ap_3})$

Now we need to merge  $Diff(xsd2, m_{ap_4})$  into  $rdb3$

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## Change Propagation (cont'd)

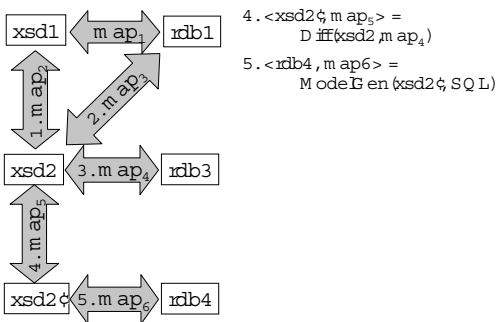


4.  $\langle xsd2c, m_{ap_5} \rangle = Diff(xsd2, m_{ap_4})$

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## Change Propagation (cont'd)



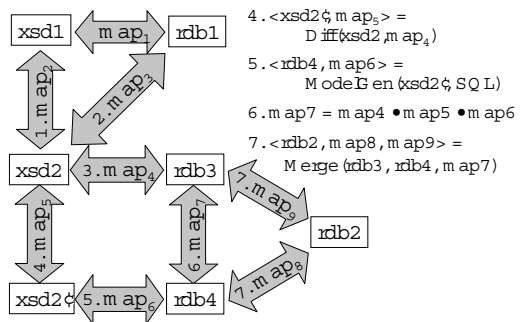
4.  $\langle xsd2c, m_{ap_5} \rangle = Diff(xsd2, m_{ap_4})$

5.  $\langle rdb4, m_{ap_6} \rangle = Merge(xsd2c, SQL)$

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## Change Propagation (cont'd)



4.  $\langle xsd2c, m_{ap_5} \rangle = Diff(xsd2, m_{ap_4})$

5.  $\langle rdb4, m_{ap_6} \rangle = Merge(xsd2c, SQL)$

6.  $m_{ap_7} = m_{ap_4} \bullet m_{ap_5} \bullet m_{ap_6}$

7.  $\langle rdb2, m_{ap_8}, m_{ap_9} \rangle = Merge(rdb3, rdb4, m_{ap_7})$

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## Complete Script in Rondo

```

OperatorDefinition: PropagateChanges(s1, d1, s1_d1, s2, c, s2_c)
1. s1_s2 = Match(s1, s2);
2. (d1 & d1 & d1) = Delete(d1, Traverse(All(s1) - Domain(s1_s2), s1_d1));
3. (c & c & c) = Extract(c, Traverse(All(s2) - Range(s1_s2), s2_c));
4. c & d1 & c = c & c * Invert(s2_c) * Invert(s1_s2) * s1_d1 * Invert(d1 & d1);
5. (d2, c & d2, d1 & d2) = Merge(c & d1 & c & d1 & c);
6. s2_d2 = s2_c * Invert(c & c) * c & d2 +
   Invert(s1_s2) * s1_d1 * Invert(d1 & d1) * d1 & d2;
7. return (d2, s2_d2);

```

OperatorUse:

```
SQL XSD: PropagateChanges(s1, d1, s1_d1, s2, Mode(s2, XSD));
```

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## Model reintegration

### 1 Design pattern

Reconcile independent changes

All changes of each model

No "duplicate additions"

### 1 Simplified example

"Additions" = add model element

(also: drop constraints, reorg. model)

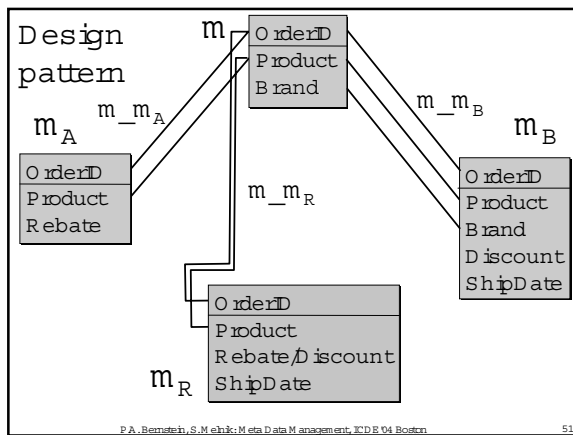
"Deletions" = delete model element

(also: add constraints, reorg. model)

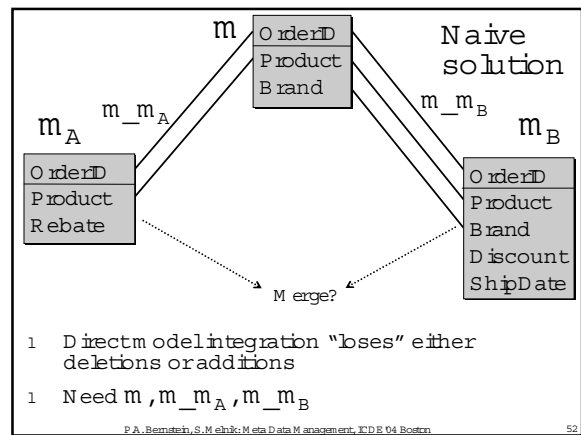
Mappings shown as lines between elements

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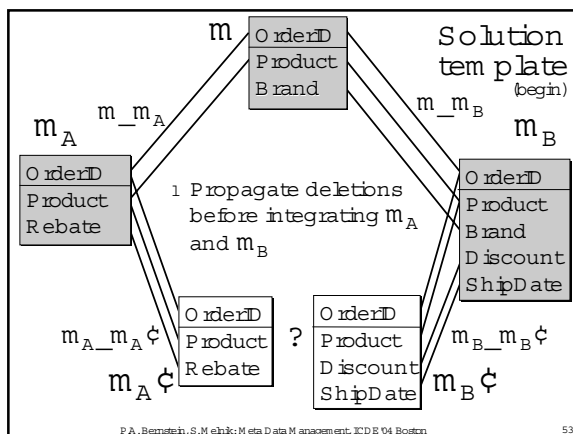
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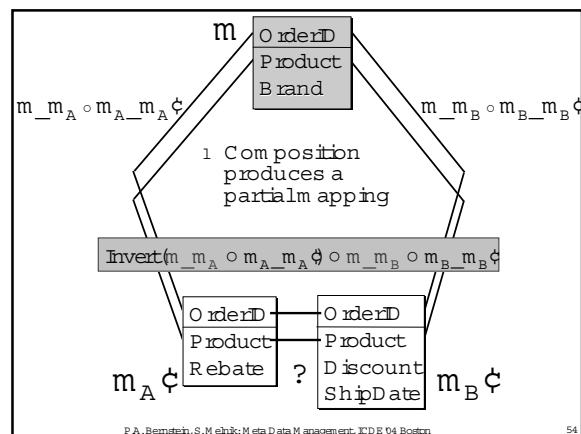
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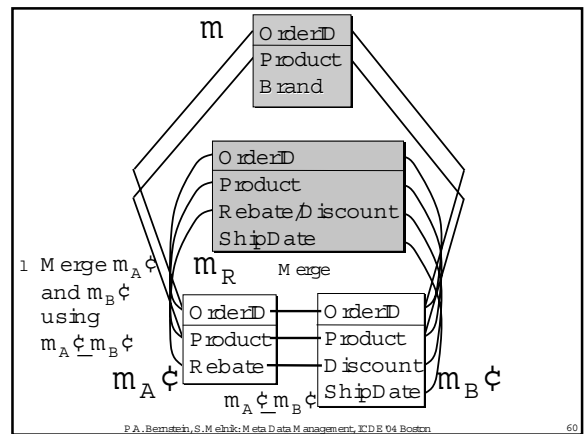
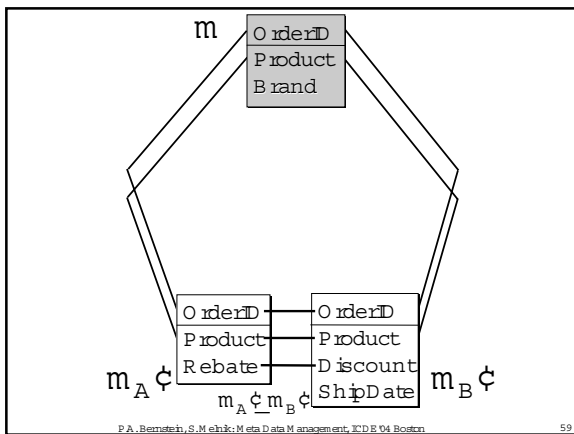
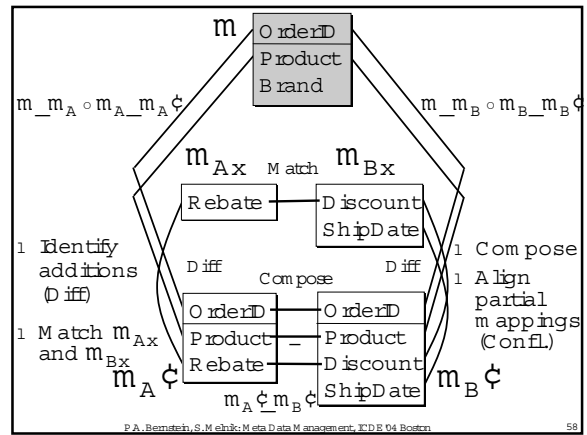
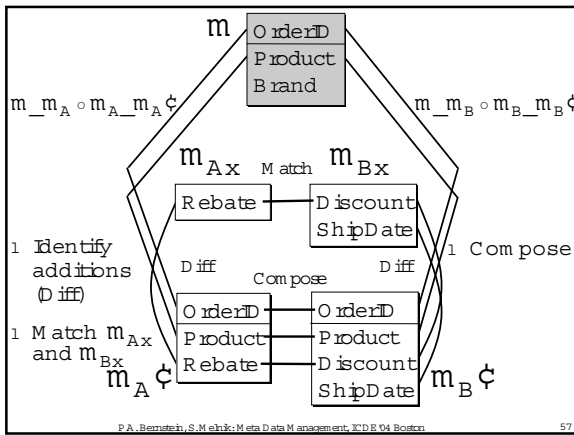
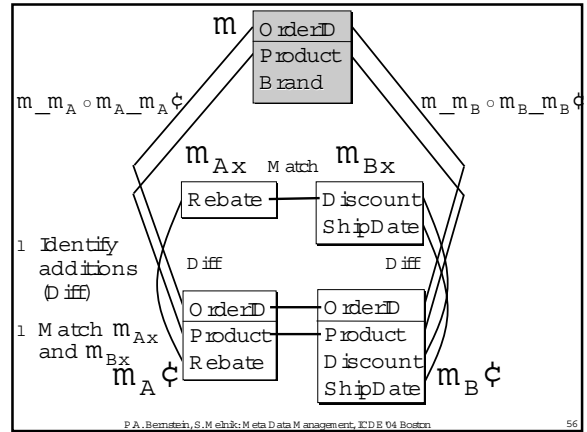
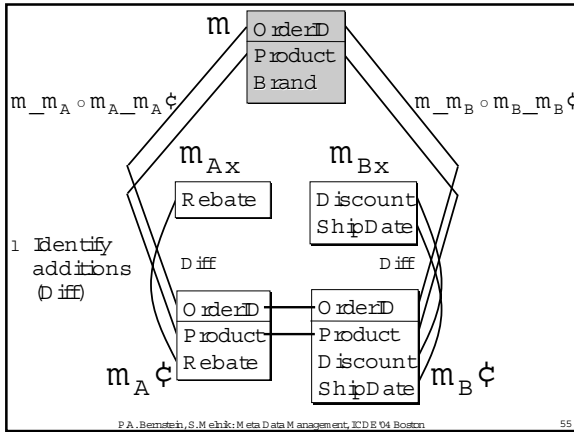
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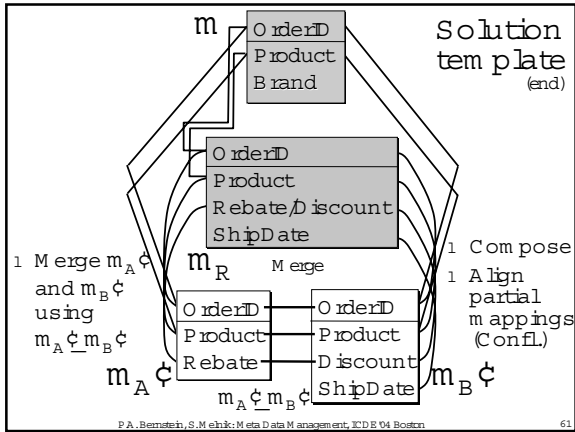


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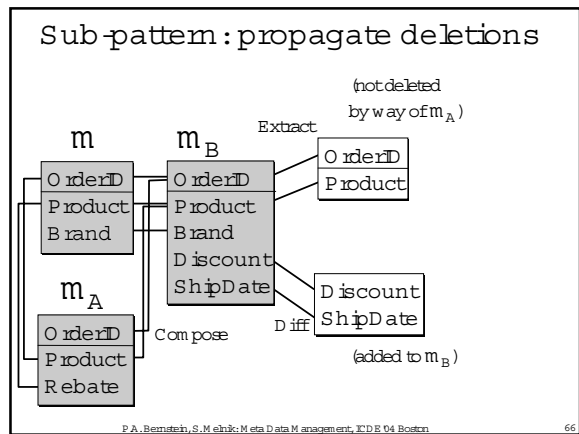
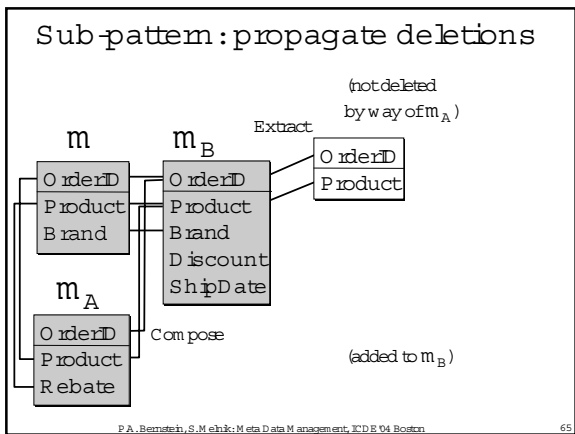
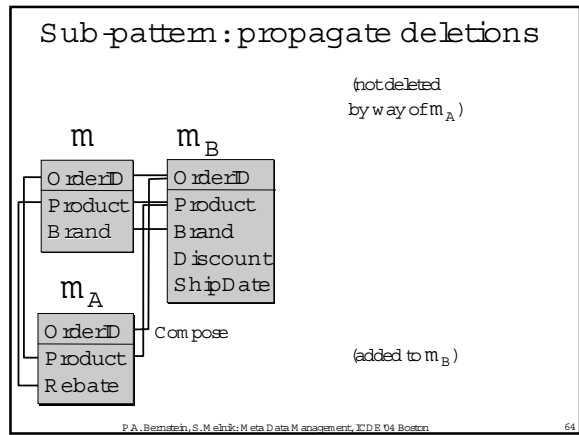
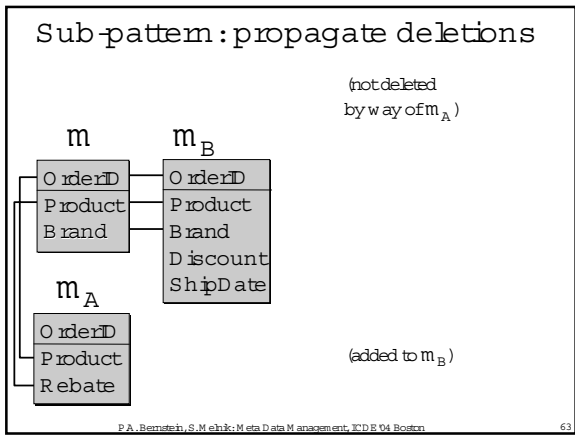
### Solution script

script R reintegrate( $m, m_A, m_m_A, m_B, m_m_B$ )

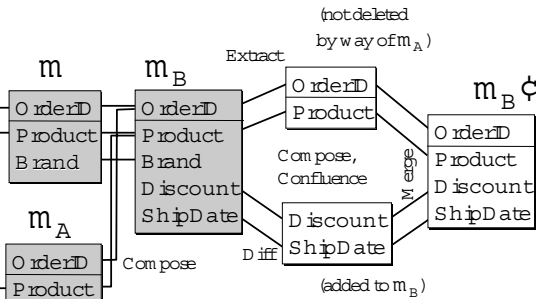
- $(m_A \zeta m_m_A \phi = \text{PropagateDeletions}(m_A, m_m_A, m_B, m_m_B));$
- $(m_B \zeta m_m_B \phi = \text{PropagateDeletions}(m_B, m_m_B, m_A, m_m_A));$
- $(m_A \times m_m_A \zeta m_m_A \phi) = \text{Diff}(m_A \zeta \text{Invert}(m_m_A \circ m_m_A \phi));$
- $(m_B \times m_m_B \zeta m_m_B \phi) = \text{Diff}(m_B \zeta \text{Invert}(m_m_B \circ m_m_B \phi));$
- $m_A \times m_m_B \zeta = \text{Match}(m_A \times m_m_B \zeta);$
- $m_A \zeta m_m_B \phi = (m_A \zeta m_A \times \circ m_A \times m_m_B \times \circ \text{Invert}(m_m_B \zeta m_m_B \phi)) \cup (\text{Invert}(m_m_A \circ m_m_A \phi \circ m_m_B \circ m_m_B \phi));$
- $(m_R, m_m_R, m_A \zeta m_m_B \phi) = \text{Merge}(m_A \zeta m_m_B \phi);$
- $m_m_R = (m_m_A \circ m_m_A \phi \circ \text{Invert}(m_m_A \phi)) \cup (m_m_B \circ m_m_B \phi \circ \text{Invert}(m_m_B \phi));$
- $\text{return}(m_R, m_m_R);$

Mapping reuse pattern

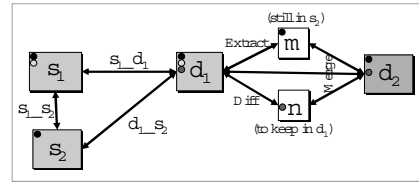
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### Sub-pattern: propagate deletions



### Solution script



```

script PropagateDeletions(s, d, s_d, s2, s_s2)
1. d_s = Invert(s_d) o s_s;
2. (m, d_m) = Extract(d, d_s);
3. (n, d_n) = D iff(d, Invert(s_d));
4. (d2, d2_m, d2_n) = Merge(m, n, Invert(d_m) o d_n);
5. d_d2 = (d_m o Invert(d_m)) ^ (d_n o Invert(d_n));
6. return (d2, d_d2);
    
```

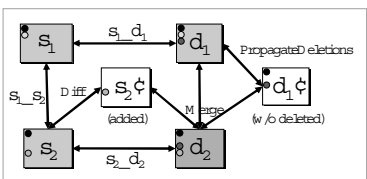
### Outline

- 1 Introduction
- 1 Meta data problem s
- 1 Design patterns
- 1 Solution templates
  - Change propagation
  - Model reintegration
  - ➔ Change propagation revisited
- 1 Research background
- 1 Wrap up

### Change propagation

- 1 Solution template
  - Propagate deletions
  - Include additions
  - Merge result

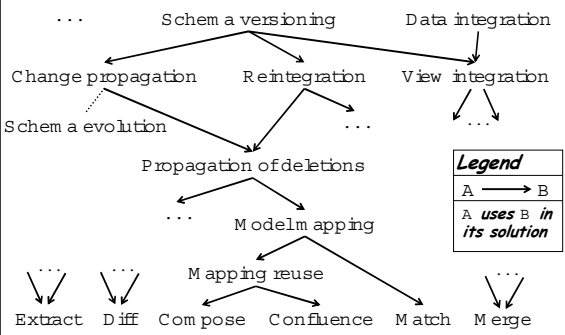
### Change propagation



```

script PropagateChanges(s, d, s_d, s2, s_s2)
1. (d_s, d_d) = PropagateDeletions(s, d, s_d, s2, s_s2);
2. (s_s, s_s') = D iff(s, Invert(s_s));
3. (d2, d2_s, d2_d) = Merge(s_s, d_s, Invert(s_s) o s_s', d_d);
4. d_d2 = (Invert(s_d) o s_s' o s_s' o Invert(d_s)) ^ (d_d o Invert(d_d));
5. s_d2 = (Invert(s_s) o s_d o d_d o Invert(d_d)) ^ (s_s' o Invert(d_s));
6. return (d2, d_d2, s_d2);
    
```

### First-cut taxonomy of patterns



## Outline

- 1 Introduction
- 1 Meta data problem s
- 1 Design patterns
- 1 Solution templates
- 1 Research background
- ➔ Wrap up

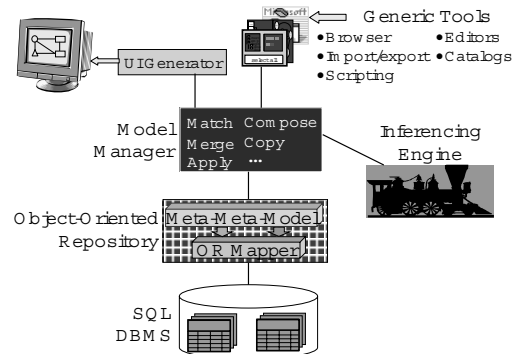
## The Commercial World

- 1 Books for IT professionals
  - A. Tanenbaum : Metadata Solutions, Addison-Wesley, 2001
  - D. Marco : Building and Managing the Meta Data Repository, Wiley, 2000
- 1 Standards-
  - UML, MOF, CWM (OMG)
  - XML, RDF, XML Schema, OWL (W3C)
- 1 Products and tools
  - Modeling: IBM RationalRose, Visio, CA AllFusion, Borland Together
  - General meta data managers: CA Advantage, Microsoft Meta Data Services, MetaIntegration
  - Meta data services in data warehousing ETL tools: Informatica, Ascential, ETL Data Advantage, ...

## The Research World

- 1 Model Management
  - A computational meta data framework based on models, mappings, and the operators described here (Match, Merge, Compose, ...)
- 1 Meta Data is a very active research area
  - Papers coming from many DB research groups
  - Some are problem-focused (e.g. data integration)
  - Some are operator-focused (e.g. Match, Merge)

## MM System Architecture



## Summary

- 1 Many DB problem s are easier to solve by manipulating meta data
- 1 Meta data problem s and solutions strongly resemble one another
- 1 Methodology: Use design patterns, solution templates, and operators to simplify development of meta data applications
- 1 There is much research to be done

## References

- 1 <http://research.microsoft.com/db/ModelMgt>
- 1 Overview
  - Bernstein, CDR 2003
  - Bernstein, Halvay, Pottinger, SIGMOD Record, Dec. 2000
- 1 Implementation
  - Mehk, Rahm, & Bernstein, SIGMOD 2003
  - and J. Web Semantics 1, 2003
- 1 Data Warehouse Exam ples
  - Bernstein & Rahm, ER 2000
- 1 Match Operation
  - Survey: Rahm & Bernstein, VLDB J., Dec. 2001
- 1 Merge Operation
  - Pottinger & Bernstein, VLDB 2003



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## The Match "Operator"

- 1 Schema matching (mapping discovery)  
Given two schemas, return correspondences that specify pairs of related elements
- 1 Semantic Mapping (query discovery)  
Given correspondences between two schemas, return an expression that translates instances of one schema into instances of the other.

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## Schema Matching Problem

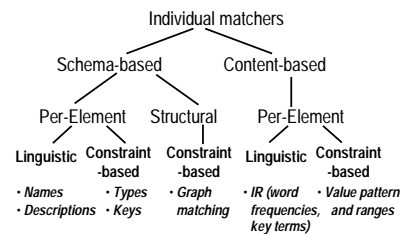
- 1 Input  
Schemas  $S_1$  and  $S_2$   
Possibly data instances for  $S_1$  and  $S_2$   
Background knowledge - thesauri, validated matches, standard schemas, constraints (keys, data types), ontologies, NL glossaries, etc.
- 1 Output  
Correspondences between elements of  $S_1$  and  $S_2$

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## Schema Matching Approaches

- 1 Many good ideas  
Rahm & Bernstein, VLDB J, Dec '01
- 1 But none are robust **fi** combine ideas

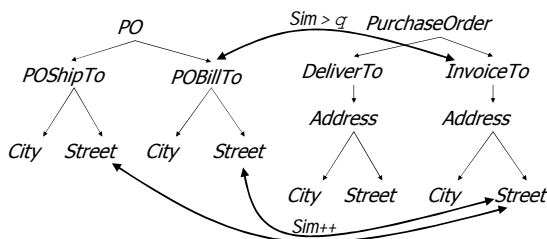


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## The Cupid Algorithm

- 1 Computes linguistic similarity of element pairs
- 1 Computes structural similarity of element pairs
- 1 Generates a mapping



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## Matching Anatomy Ontologies

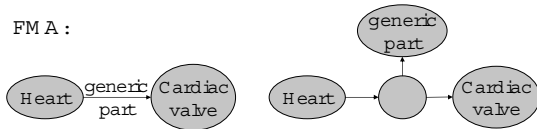
- 1 Match two human anatomy ontologies  
FMA - Univ. of Washington  
Galen CRM - Univ. of Manchester (UK)  
By Peter Mork (Univ. of Washington)  
Both models are big
- 1 Ultimate goal was finding differences
- 1 Like most match algorithms, ours calculates a similarity score for the  $m \cdot n$  pairs of elements

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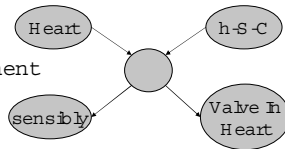
## Aligning Representations

FMA :



CRM :

Heart sensibly  
hasStructuralComponent  
ValveInHeart



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## Anatomy Matching Algorithm

### 1. Lexical Match

- Normalize string, UMLS dictionary lookup, convert to concept-ID from thesaurus

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## Anatomy Matching Algorithm

### 1. Lexical Match

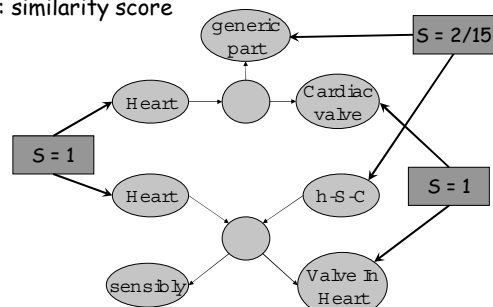
- Normalize string, UMLS dictionary lookup, convert to concept-ID from thesaurus
- String comparison **fi** 306 matches
- Adding spaces, ignoring case **fi** 1834 matches
- Lexical tools **fi** 3503 matches

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## Anatomy Matching Example

S: similarity score



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## Anatomy Matching Algorithm

### 1. Lexical Match

- Normalize string, UMLS dictionary lookup, convert to concept-ID from thesaurus

### 2. Structure Match

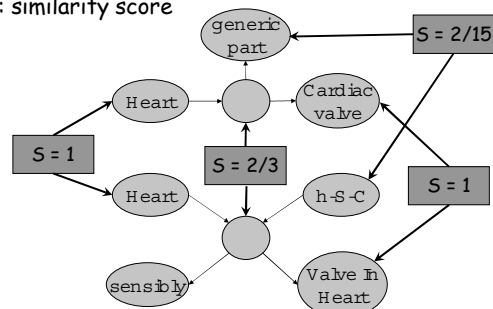
- Similarity (refined nodes)  
= Average (neighbors)
- Back-propagate to neighbors

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## Anatomy Matching Example

S: similarity score



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## Anatomy Matching Algorithm

1. Lexical Match
    - Normalize string, UMLS dictionary lookup, convert to concept-ID from thesaurus
  2. Structure Match
    - Similarity (reified nodes) = Average (neighbors)
    - Back-propagate to neighbors
- Adds 64 matches (to previous 3503)
  - Implies 875 reified relationship matches

## Anatomy Matching Algorithm

1. Lexical Match
    - Normalize string, UMLS dictionary lookup, convert to concept-ID from thesaurus
  2. Structure Match
    - Similarity (reified nodes) = Average (neighbors)
    - Back-propagate to neighbors
  3. Align Super-classes
    - Super-class similarity = average similarity of children, grandchildren, great-grandchildren
- Adds 213 matches (to 3567)

## Some Lessons

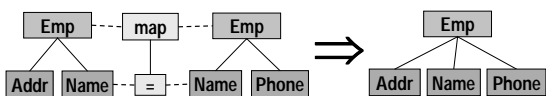
1. A common encoding of models is hard and involves compromises
  - Different styles of reifying relationships
  - CRM stores transitive relationships
1. Match needs to invent generalizations
  - In FMA, **arterial supply, venous drainage, nerve supply, lymphatic drainage**
  - In CRM, these all map to **isServedBy**
1. On big models, Match is expensive
  - Some steps required days to execute
  - Cross-product filled 80 GB (< 1GB input).

## Outline

- Introduction to Model Management
  - Using MM to solve meta data problems
  - Matching anatomy ontologies
1. Model merging
  1. Wrap-up

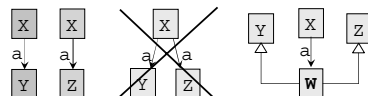
## Merge ( $M_1, M_2, map$ )

1. Return the union of models  $M_1$  and  $M_2$ 
  - Use map to guide the Merge
  - If elements  $x = y$  in map, then collapse them into one element



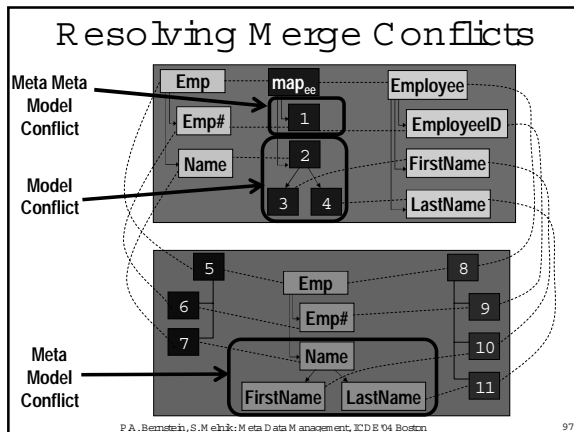
## Merge ( $M_1, M_2, map$ )

1. [Buneman, Davidson, Kosky, EDBT 92]
  - Meta-model has aggregation & generalization only
  - Union, and collapse objects having the same name
  - Fix-up step for inconsistencies created by merging



Successive fixups lead to different results **L**  
 Batch them at the end, to get a unique minimal result  
 Now enrich the meta-model (containment, complex mappings, ...) & merge semantics (conflicts, deletes)





- ### Contributions to Merge [Pottinger & Bernstein, VLDB 03]
- 1 Generic correctness criteria for Merge
  - 1 Use of first-class input mapping (not just correspondences)
  - 1 Taxonomy of conflicts & resolution strategies
  - 1 Characterize when Merge can be automatic
  - 1 A merge algorithm for an EER representation
  - 1 Experimental evaluation
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- ### An Approach to ModelGen [Atzeni & Torlone, EDBT 96]
- 1 Meta-models are made of patterns
 

Object has sub-object $\Delta$	Aggregation $\otimes$ has attributes $\square$	Aggregation $\otimes$ has key $\ddagger$
(a)	(b)	(c)
  - 1 Define pattern transformations as rules  
For XSDL to SQL,  $\Delta \Rightarrow \square + \ddagger$
  - 1 To translate  $M_s$  into meta-model  $(MM_t)$ ,  
Apply rules that replace patterns in  $M_s$  that are not in  $MM_t$  by patterns that are in  $MM_t$
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- ### ModelGen Research
- 1 More complete repertoire of patterns
  - 1 Make patterns more generic
  - 1 Integrate with rules engine (avoid cycles, control search)
  - 1 Implement it
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