4.M iscellaneous ProjectIssues

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Lock Conversions

- This is step 3 of the course project
- Lock conversion upgrading an r-lock to a w -lock
 eg., T_i = read(x)... w rite(x)
- The main purpose of step 3 is to ensure you understand the lock manager code we supply.
- Deadlocks are an issue
 - if two txns converta lock concurrently, they'll deadlock (both get an r-lock on x before eithergets a w -lock)
 - To avoid the deadlock, a caller can get a w -lock first and down-grade to an r-lock if it doesn't need to write.

Lecture D ependencies

- Step 5 involves in plementing a workflow controller.
 It depends on
 - Transaction bracketing
 - Partitioning work between W C and transaction servers
 (= RM in the project)
 - Propagating transaction context in RM calls
 - Param eter-based routing (e.g. choose an RM based on flightnum ber)
 - I'll cover this on Feb 8.See also Chapter 2.
- Steps 6 10 involve two-phase commit.
 - I'll cover it by Feb 15. See also Chapter 9.

Process Structure

- For architectural reasons, RM and W C should be processes.
- Ordinarily, there's a TM pernetwork node. - To simulate this, a TM should be a process