

# CSEP 561 – Bits and Links

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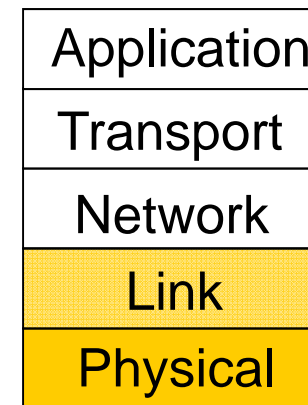
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# Topic

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- How do we send a message across a wire or wireless link?
- The physical/link layers:
  1. Different kinds of media
  2. Fundamental limits
  3. Encoding bits
  4. Model of a link

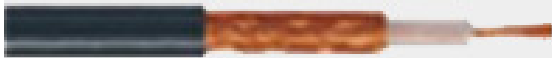


# 1. Wires

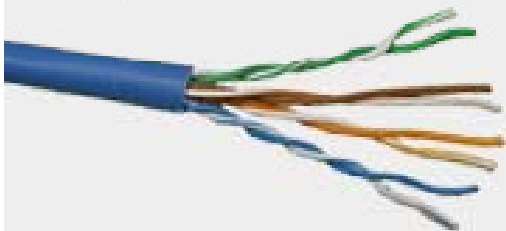
**10BASE5 - "Thicknet"**



**10BASE2 - "Thinnet"**



**10BASE-T**

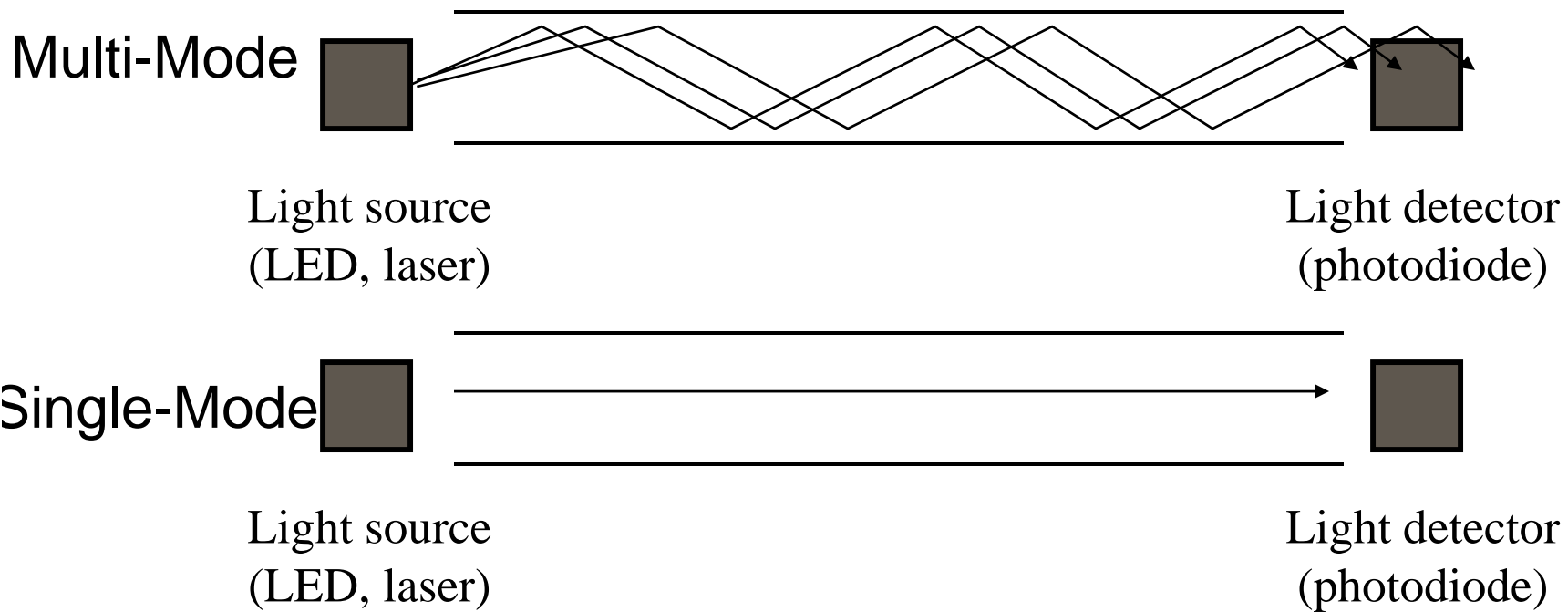
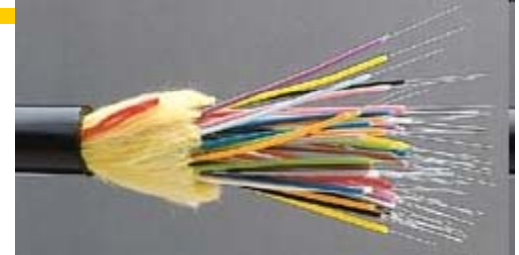


Now Cat 6, Cat 7 for GigE, four pairs

- Twisted pairs: twists reduce RF emission / crosstalk; also shielding can be added
- Coaxial cable: inner and outer ring conductor for superior noise immunity
- Many different specs/grades depending on application
- 100s of MHz for 100s of meters

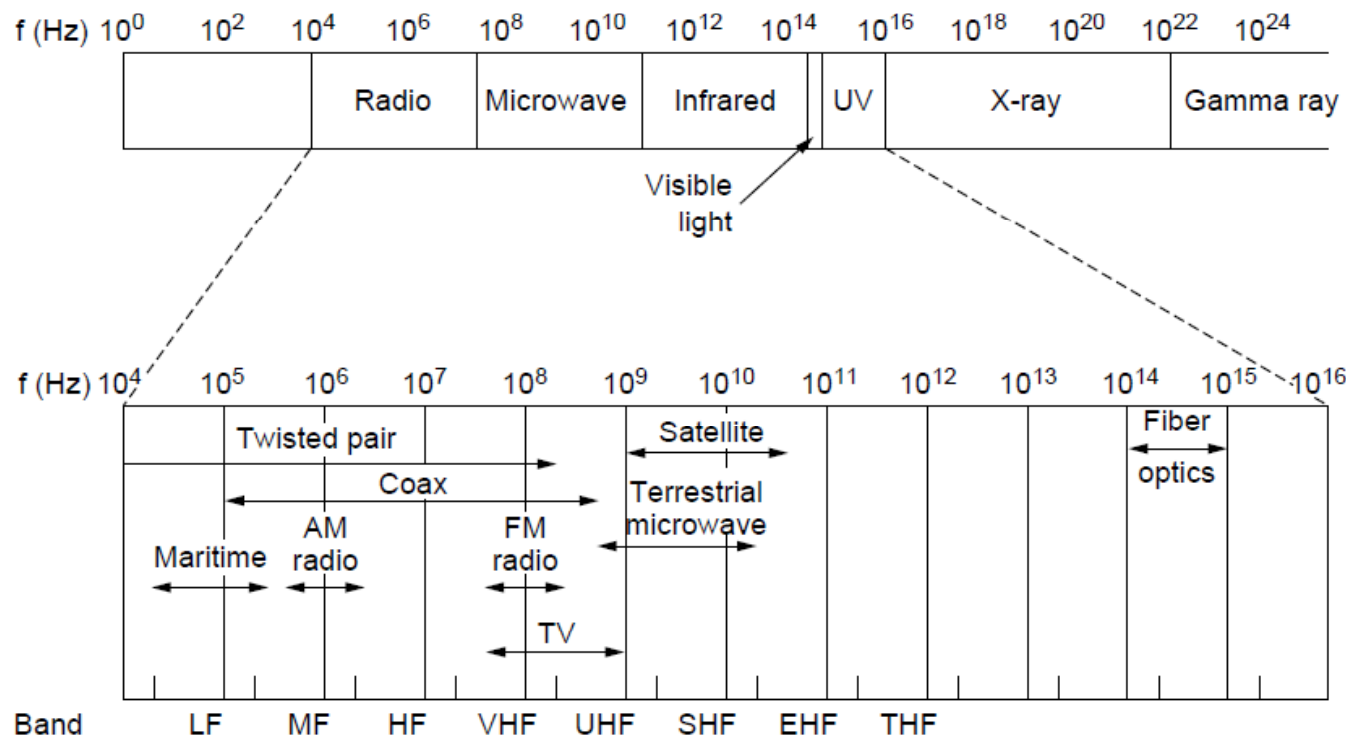
# Fiber Optic Cable

- Long, thin, pure strand of glass
  - light propagated with total internal reflection
  - enormous bandwidth available (terabits)



# Wireless

- Different frequencies have different properties
- Signals subject to atmospheric/environmental effects

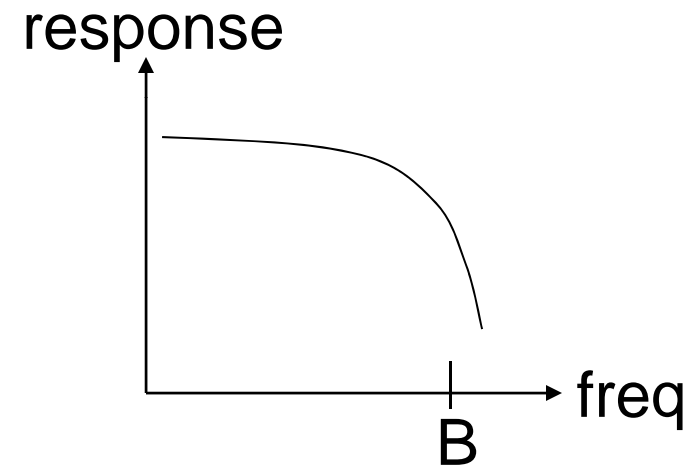




# Model of a wire

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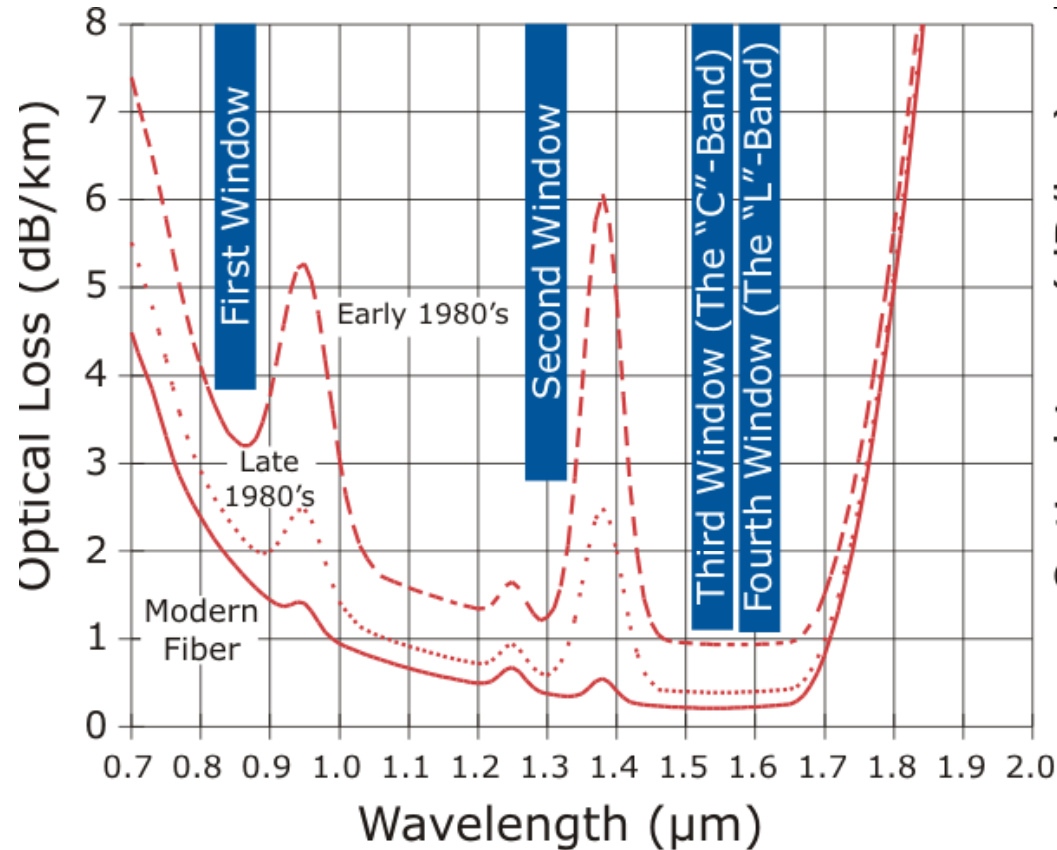
- Frequencies beyond cutoff highly attenuated
  - Bandwidth = passband (Hz)
- Signal also subject to:
  - Attenuation (repeaters)
  - Distortion (frequency and delay)
  - Noise (thermal, crosstalk, impulse)



EE: Bandwidth = width of frequency passband, measured in Hz  
CS: Bandwidth = information carrying capacity, measured in bits/sec

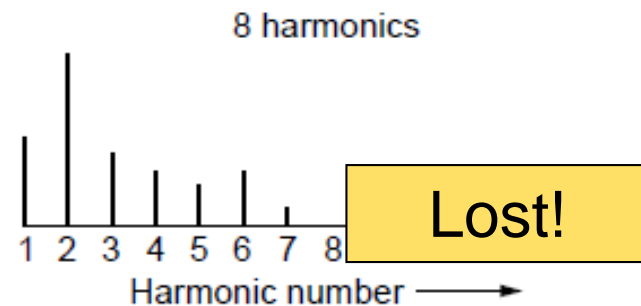
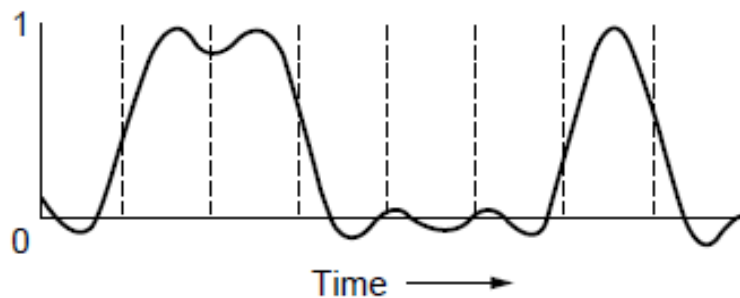
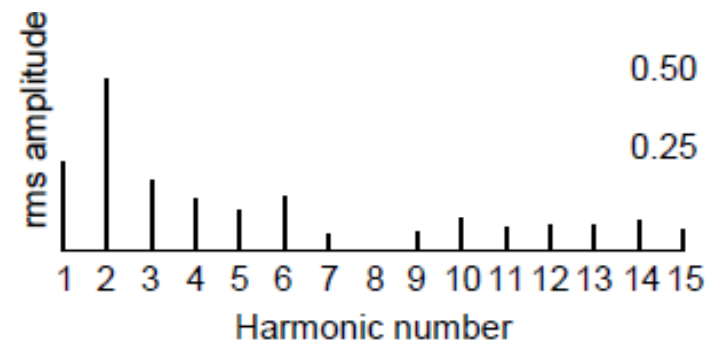
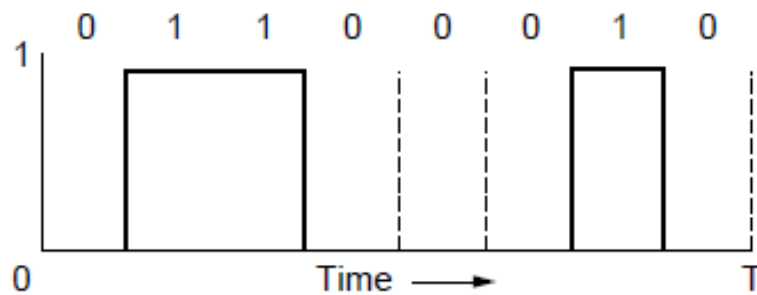
# Attenuation of optic fiber

- Enormous bandwidth in each window





## 2. Effect of limited bandwidth



- Less bandwidth permits less rapid signal transitions

# Signals over a wire

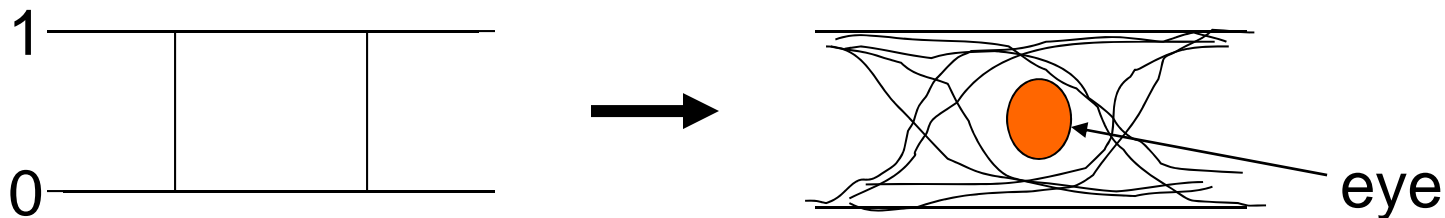
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- If we send a waveform, what do we get out?
  - Signal loss
  - Delay
  - Frequency-dependent attenuation
  - Noise at the receiver
- What do we want to get out?
  - Fidelity versus interpretability

# Nyquist Limit (~1924)

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- For a noiseless channel with bandwidth B
- Symbols will be distorted, and sending too fast leads to Inter-symbol Interference (ISI)

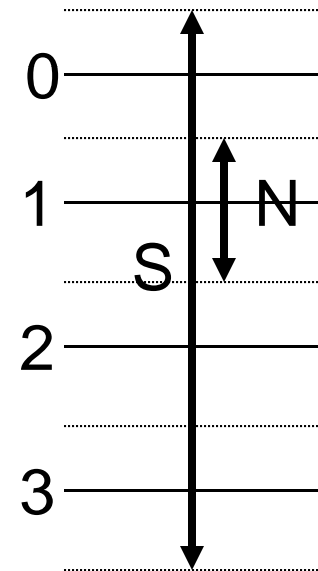


- The maximum rate at which it is possible to send:  
 $R = 2B$  symbols/sec  
e.g., 3KHz  $\rightarrow$  6Ksym/sec

# Taking Noise into Account

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- Noise limits how many signal levels we can safely distinguish between
  - $S = \text{max signal amp.}$ ,  $N = \text{max noise amp.}$
- The number of bits per symbol depends on the number of signal levels
  - E.g, 4 levels implies 2 bits / symbol



# Shannon limit (1948)

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$$\text{Capacity} = \text{Bandwidth} \times \log_2 \left( 1 + \frac{\text{Signal}}{\text{Noise}} \right)$$

↑                      ↑                      ↑

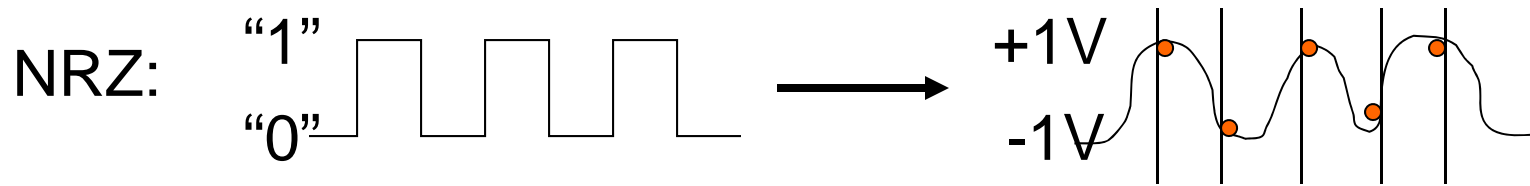
Bits/sec              How fast signal              How many signal  
                                 can change                      levels can be seen

- Shannon, “A Mathematical Theory of Communication,” 1948.
- SNR or Signal-to-Noise ratio defined as deciBels on a log scale
  - $\text{SNR} = 10 \log_{10} (\text{Signal} / \text{Noise})$ , e.g., 30dB = 1000 times

# 3. Encoding Bits with Signals

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- Generate analog waveform (e.g., voltage) from digital data at transmitter and sample to recover at receiver

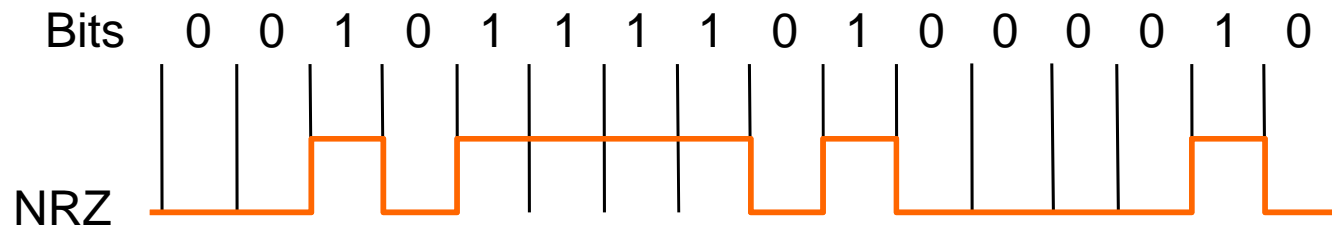


- We send/recover symbols that are mapped to bits
  - May have >2 different symbols, e.g., amplitudes
  - Thus distinguish symbol rate versus bit rate
- This is baseband transmission

# NRZ

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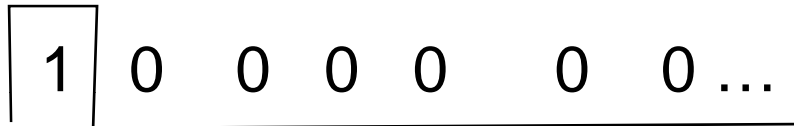
- Simplest encoding, NRZ (Non-return to zero)
  - Use high/low voltages, e.g., high = 1, low = 0



# Issue: Clock recovery

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- Um, how many zeros was that again?



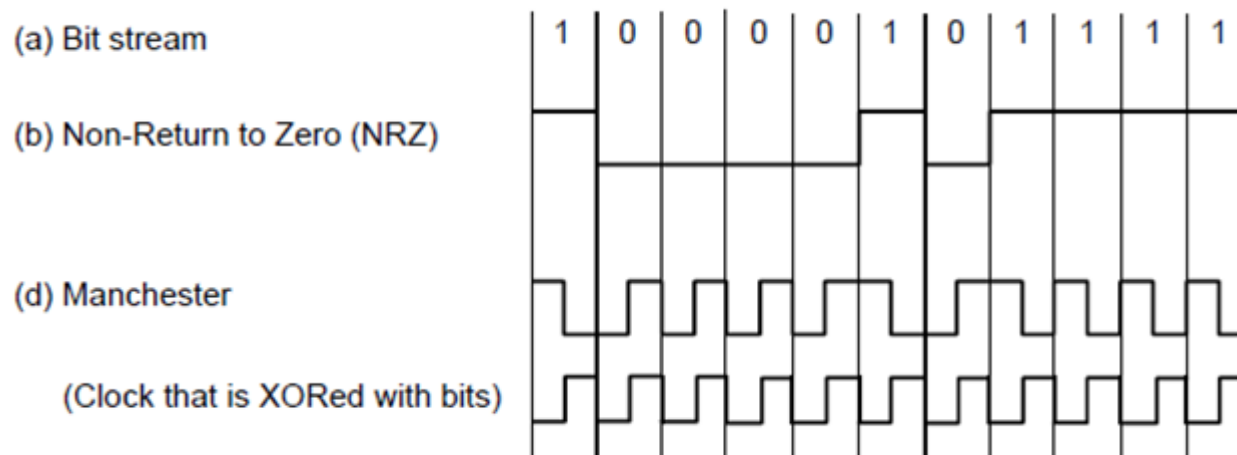
If sender and receiver have exact clocks no problem. But they don't!

- Any brilliant ideas?



# Embed clock in signal (Manchester)

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- Signal is XOR of data (NRZ) and clock (transition per bit)
  - Low-to-high is 0; high-to-low is 1
  - Signal rate is twice the bit rate
- Advantage: self-clocking, Disadvantage: BW inefficiency

# 4B/5B Codes

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- We want self-clocking transitions *and* efficiency ...
- Solution: map data bits (which may lack transitions) into code bits (which are guaranteed to have them)
- 4B/5B code:
  - 0000 → 11110, 0001 → 01001, ... 1111 → 11101
  - Never more than three consecutive 0s back-to-back
  - 80% efficiency, plus use “illegal” codes as markers
- Many more complex codes are available; some use multiple voltage level

# Scrambling

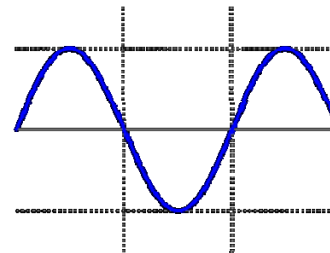
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- XOR data with known pseudo-random sequence
  - Can generate cheaply with linear feedback shift registers (LFSR)
  - Causes transitions with reasonable probability
  - Also tends to whiten data (better for RF)
- Reverse at receiver by XORing with same sequence

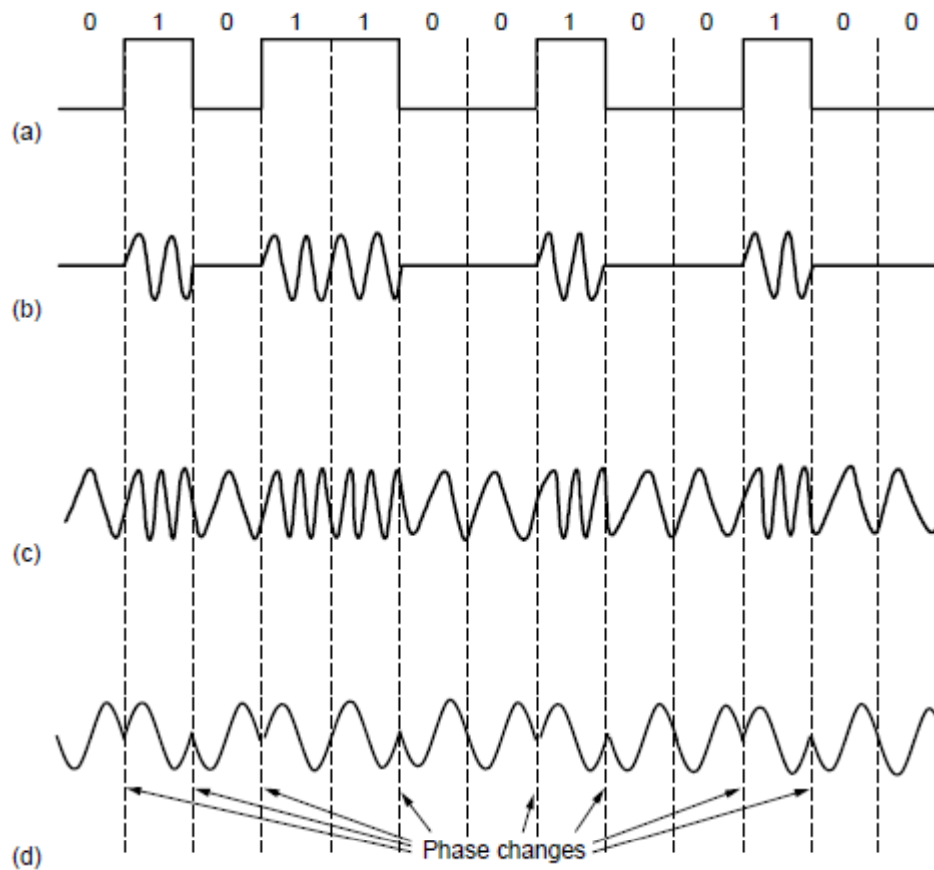
# Passband transmission

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- For wireless, fiber, need to encode signal by modulating carrier wave ... can't propagate at baseband
- Carrier frequency set by assigned bandwidth,
  - e.g., 2.45GHz WiFi
- Modulation: can change carrier
  - Amplitude
  - Phase/frequency



# Modulation examples



(a) A binary signal

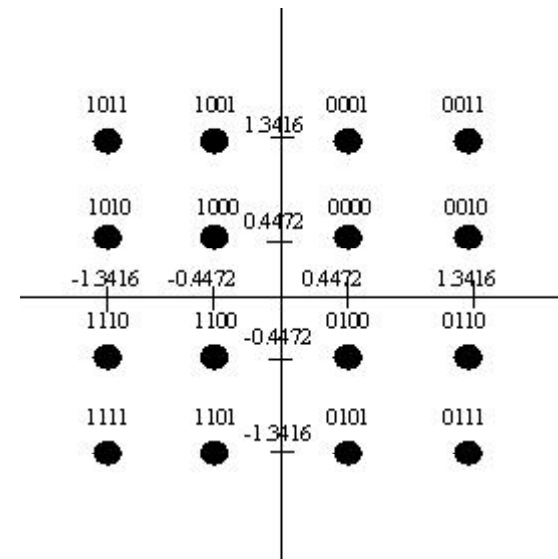
(b) Amplitude shift keying

(c) Frequency shift keying

(d) Phase shift keying

# Constellations

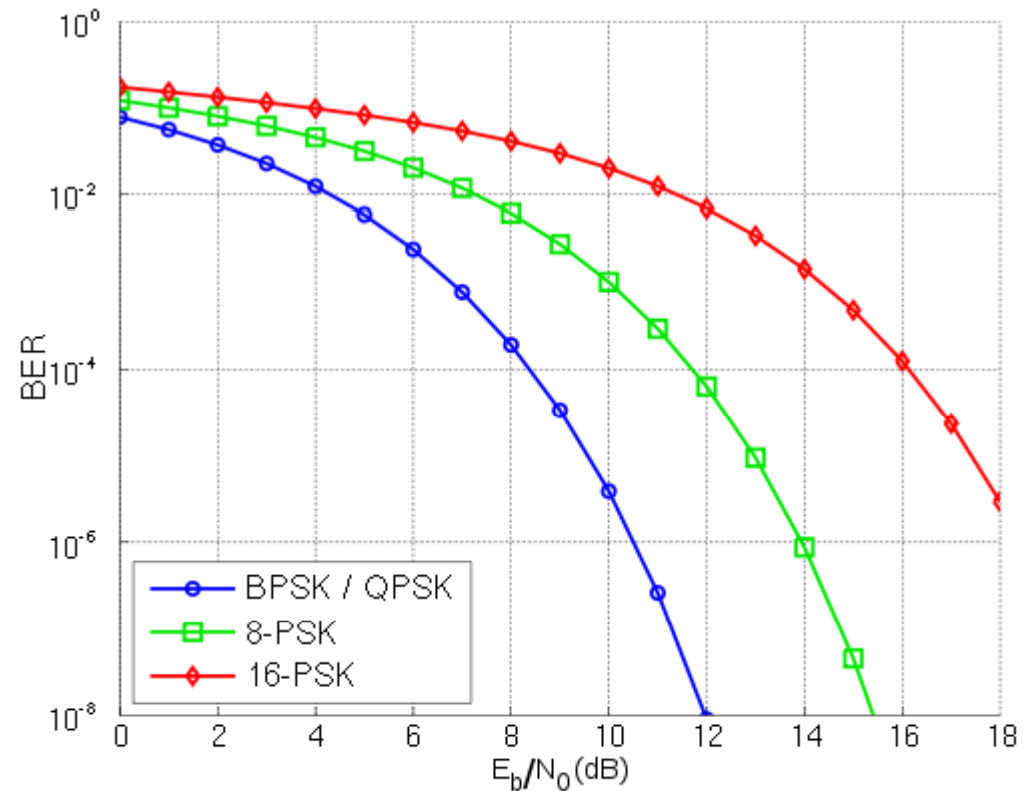
- Express modulation as a “constellation”
  - Points are amplitude/phase modulations for valid symbols
- Many names for schemes:
  - BPSK, QPSK ... QAM



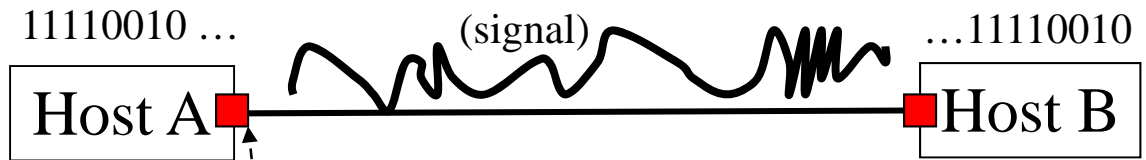
QAM 16 constellation  
in 3G (HSPDA)

# BER versus SNR

- Need higher SNR for more complex modulations to keep a low bit error rate

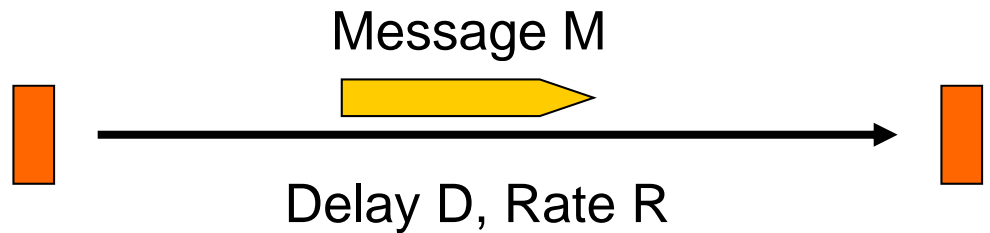


# 4. Abstract model of a link



What really happens

Network interface cards (NICs)  
(also called “network adaptors”)



Abstract link for our purposes



# Model

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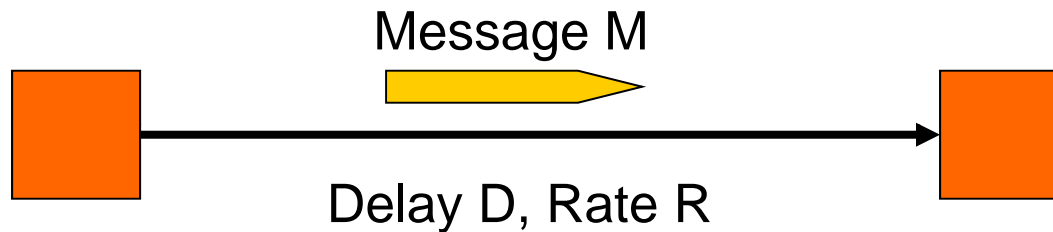


- Typically all we will need (but not so good for wireless!)
- Other parameters that are important:
  - Whether the media is broadcast or not
  - The kind and frequency of errors (bit error rate, BER)

# Message Latency

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- How long does it take to send a message?



- Two terms:
  - Propagation delay = distance / speed of signal in media
    - How quickly a message travels over the wire
    - $2/3c$  for copper wire
  - Transmission delay = message (bits) / rate (bps)
    - How quickly you can inject the message onto the wire
- Later we will see queuing delay ...

# One-way Latency

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Dialup with a modem:

- $D = 10\text{ms}$ ,  $R = 56\text{Kbps}$ ,  $M = 1024$  bytes
- Latency =  $10\text{ms} + (1024 \times 8) / (56 \times 1024)$  sec = 153ms!

Cross-country with T3 (45Mbps) line:

- $D = 50\text{ms}$ ,  $R = 45\text{Mbps}$ ,  $M = 1024$  bytes
- Latency =  $50\text{ms} + (1024 \times 8) / (45 \times 1024 \times 1024)$  sec = 50ms!

- Either a slow link or long wire makes for large latency

# Bandwidth-delay product: Messages occupy space on the wire

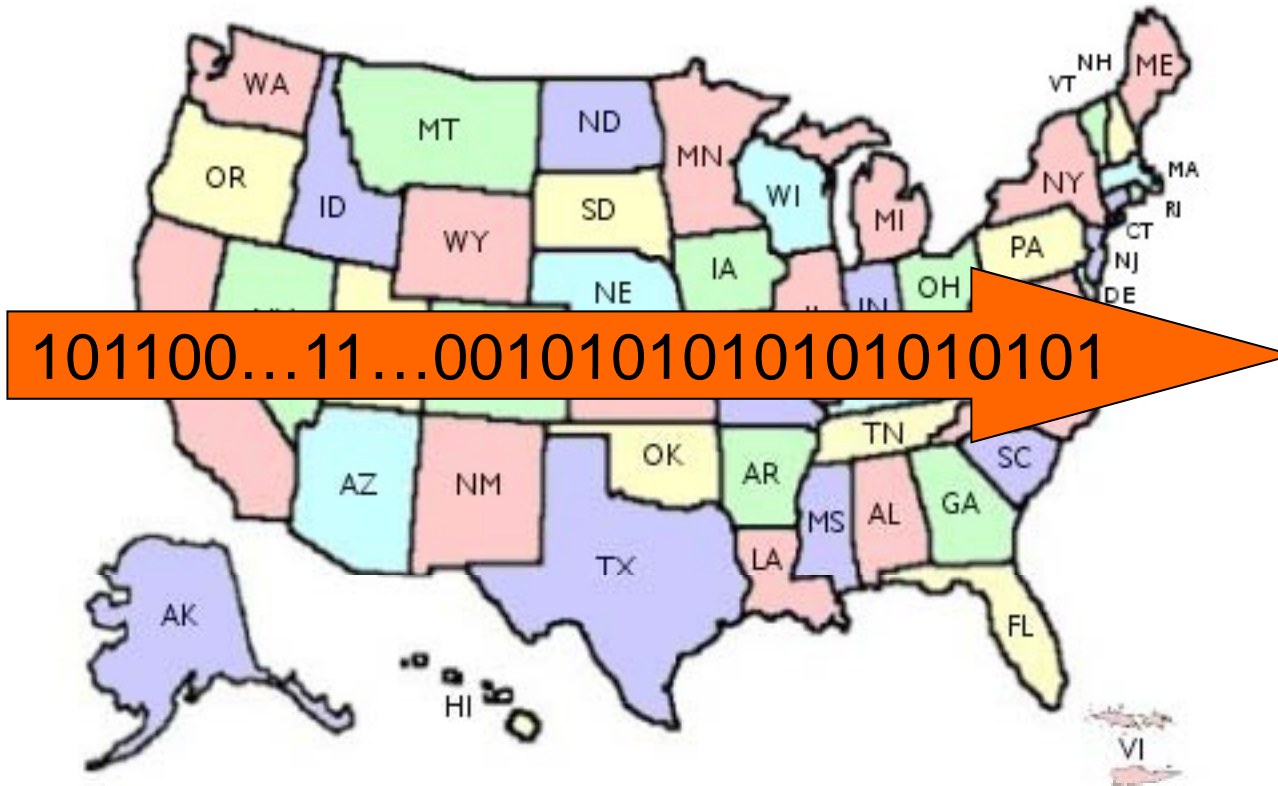
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- Consider a 1b/s network, suppose latency is 16 seconds.
  - How many bits can the network “store”?
  - This is the bandwidth-delay product
  - Measure of “data in flight.”
  - $1\text{b/s} * 16\text{s} = 16\text{b}$
- Tells us how much data can be sent before a receiver sees any of it.
  - Twice BD tells us how much data we could send before hearing back from the receiver something related to the first bit sent.

# BD Example

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$$\text{BD} = 50\text{ms} * 45\text{Mbps} = 2.25 * 10^6 = 280\text{KB}$$



# Wireless versus Wired links

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- Wireless complications:
  - Broadcast channel has interference effects
  - Link capacity varies lots over time, e.g., as endpoints move
  - Which “wireless links” are up even varies over time
- Endpoint moves → SNR changes due to RF effects → rate must go down if SNR falls to keep low BER; or rate wants to go up if SNR rises to use spectrum efficiently
- Wired is about engineering the right link properties
- Wireless is about adapting to the channel capacity