

**Technology Overview** 

#### Modern SAT Solvers: Key Advances and Applications

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### Thanks...

For sharing their valuable knowledge, wisdom, slides, and more:

- Paul Beame, Henry Kautz (1998-2005, UW)
- Carla Gomes, Bart Selman (2005-2010, Cornell)
- Too many other colleagues to begin to mention here...

For exciting opportunities that lie ahead:

"Watson" and colleagues (2010-, IBM Watson)

For this wonderful chance to be back at UW!

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### Outline

- The challenge of automated reasoning
- Dramatic progress in SAT solvers in the last 20 years
- Applications of SAT
- Key advances in SAT reasoning technology
- Advances in our understanding
- Looking ahead: where is SAT research headed?
- Summary



### The Quest for Machine Reasoning

A cornerstone of Artificial Intelligence

Objective: Develop foundations and technology to enable effective, practical, large-scale automated reasoning

Machine Reasoning (1960-90s)

Computational complexity of reasoning *appeared* to severely limit real-world applications

**Current reasoning technology** 

Revisiting the challenge: Significant progress with new ideas / tools for dealing with complexity (scale-up), uncertainty, and multiagent reasoning

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### **General Automated Reasoning**





### **Challenges in Automated Reasoning**



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### Exponential Complexity Growth: The Challenge of Complex Domains





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### Progress in the Last 20 Years

Focus: Combinatorial Search Spaces specifically, the Boolean satisfiability problem, SAT

Significant progress since the 1990's. How much?

Problem size: We went from 100 variables, 200 constraints (early 90's) to 1,000,000+ variables and 5,000,000+ constraints in 20 years

Search space: from 10^30 to 10^300,000. [Aside: "one can encode quite a bit in 1M variables."]

- Is this just Moore's Law? It helped, but not much...
  - 2x faster computers does *not* mean can solve 2x larger instances
  - search difficulty does not scale linearly with problem size!
- Tools: 50+ competitive SAT solvers available

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# How Large are the Problems?

#### A bounded model checking (BMC) problem:

From "SATLIB":

http://www.satlib.org/benchm.html

SAT-encoded bounded model checking instances (contributed by Ofer Shtrichman)

In Bounded Model Checking (BMC) [BCCZ99], a rather newly introduced problem in formal methods, the task is to check whether a given model M (typically a hardware design) satisfies a temporal property P in all paths with length less or equal to some bound k. The BMC problem can be efficiently reduced to a propositional satisfiability problem, and in fact if the property is in the form of an invariant (Invariants are the most common type of properties, and many other temporal properties can be reduced to their form. It has the form of 'it is always true that ... '.), it has a structure which is similar to many Al planning problems.

#### Source: IBM

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### SAT Encoding

(automatically generated from problem specification)

p cnf 51639 368352	
-170	
-160	i.e., ((not $x_1$ ) or $x_7$ )
-150	$((not x_1) or x_2)$
<u> </u>	
<b>-130</b>	610.
<b>-120</b>	v v v oto ovo ouv Pooloon voriabloo
<u>-1 -8 0</u>	$x_1, x_2, x_3$ , etc. are our boolean variables (to be set to True or Feleo)
— <b>9</b> 15 0	(to be set to true of raise)
— <b>9</b> 14 0	
— <b>9</b> 13 0	Chauld y, he est to Felee 20
<b>-9 -12 0</b>	Should X <sub>1</sub> be set to False??
<b>-9 11 0</b>	
<b>-9 10 0</b>	
<b>-9 -16 0</b>	
-17 23 0	
-17 22 0	



### 10 Pages Later:





### 4,000 Pages Later:

. . .



### Finally, 15,000 Pages Later:

 $\begin{array}{c} -7\ 260\ 0\\ 7\ -260\ 0\\ 1072\ 1070\ 0\\ -15\ -14\ -13\ -12\ -11\ -10\ 0\\ -15\ -14\ -13\ -12\ -11\ 10\ 0\\ -15\ -14\ -13\ -12\ 11\ -10\ 0\\ -15\ -14\ -13\ -12\ 11\ 10\ 0\\ -7\ -6\ -5\ -4\ -3\ -2\ 0\\ -7\ -6\ -5\ -4\ -3\ 2\ 0\\ -7\ -6\ -5\ -4\ 3\ -2\ 0\\ -7\ -6\ -5\ -4\ 3\ 2\ 0\\ 185\ 0\end{array}$ 

Search space of truth assignments:  $2^{50000} \approx 3.160699437 \cdot 10^{15051}$ 

#### *Current SAT solvers solve this instance in just a few seconds!*



### SAT Solvers as Search Engines

- Systematic SAT solvers have become really efficient at searching fast
- E.g., on an IBM model checking instance from SAT Race 2006, with ~170k variables, 725k clauses, solvers such as MiniSat and RSat roughly:
  - Make 2000-5000 decisions/second
  - Deduce 600-1000 conflicts/second
  - Learn 600-1000 clauses/second (#clauses grows rapidly)
  - Restart

every 1-2 seconds (aggressive restarts)



### SAT Solver Progress Trend (1994 – 2001)

Solvers have continually improved over time

Instance	Posit' 94	Grasp' 96	Sato' 98	Chaff' 01
ssa2670-136	40.66s	1.20s	0.95s	0.02s
bf1355-638	1805.21s	0.11s	0.04s	0.01s
pret150_25	>3000s	0.21s	0.09s	0.01s
dubois100	>3000s	11.85s	0.08s	0.01s
aim200-2_0-no-1	>3000s	0.01s	< 0.01s	< 0.01s
2dlxbug005	>3000s	>3000s	>3000s	2.90s
c6288	>3000s	>3000s	>3000s	>3000s

Source: MarquesSilva-02



### Forces Driving Faster, Better SAT Solvers

#### From academically interesting to practically relevant

- "Real" benchmarks, with real interest in solving them
- Regular SAT Solver Competitions (Germany-89, Dimacs-93, China-96, SAT-02, SAT-03, ..., SAT-07, SAT-09, SAT-2011)
  - "Industrial-instances-only" SAT Races (2008, 2010)
  - A tremendous resource! E.g., SAT Competition 2006 (Seattle):
  - 35+ solvers submitted, downloadable, mostly open source
  - 500+ industrial benchmarks, 1000+ other benchmarks

#### 50,000+ benchmark instances available on the Internet

This constant improvement in SAT solvers is the key to the success of, e.g., SAT-based planning and verification



### Coming up: SAT Competition 2011!

#### www.satcompetition.org

- To be organized along-side the SAT-2011 Conference, Ann Arbor, MI
- Exciting event for the community

### SAT Competition 2011

A competitive event of the SAT 2011 Conference

June 19th - June 22nd 2011, Ann Arbor, MI, USA

- An annual 'check': *how much did we improve this year?*
- My task as one of the judges: ensure fairness and competitive spirit go hand-in-hand!
- If you have a great idea to improve SAT solvers, start working on it!

[submissions due March 2]



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### For details on the state of the art in SAT...

### Handbook of Satisfiability

 Everything from historical perspectives, theoretical foundations, practical solvers, applications, extensions, ...

[27 chapters contributed by numerous experts]





## Applications of SAT (1)

- Model Checking (hardware/software verification)
  - Bounded model checking (BMC)
  - Enhancement techniques: Abstraction and refinement
  - Major groups at Intel, IBM, Microsoft, and universities such as CMU, Cornell, and Princeton

### SAT has become a dominant back-end technology

- Key questions:
  - Safety property: can system ever reach state S?
  - Liveness property: will system <u>always</u> reach T after it gets to S?
- Main idea:
  - Encode "reachability" in a *finite* transition graph
  - CNF encoding: in simple terms, "similar" to planning... next slide!





## Applications of SAT (2)

### Classical Planning

- Parallel step optimal plans
- E.g., SATPLAN-06 fastest in this category at IPC-06 planning competition





#### Main idea:

- Create planning graph, unfolded up to T steps
- Encode legal transitions as CNF
  - Variables: action vars, fact vars for each time step t
  - Clauses encode action preconditions, effects, mutual exclusion, etc.
    e.g., (actionA\_(t+1) ⇒ preA1\_t AND preA2\_t AND preA3\_t), (NOT actionA\_t OR NOT actionB\_t), ...



## Applications of SAT (3)

### Combinatorial Design

- Complex mathematical structures with desirable properties
  - Latin squares
  - Quasi groups
  - Ramsey numbers
  - Steiner systems
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- Highly useful for
  - design of experiments
  - coding theory, cryptography
  - drug design and testing
  - crop rotation schedules

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## Applications of SAT (4)

- Has been applied to solve sub-problems in many domains!
  - Test pattern generation
  - Scheduling
  - Optimal Control
  - Protocol Design (e.g., for networks)
  - Multi-agent systems
  - E-Commerce (E-auctions and electronic trading agents), etc.



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### Key Advances in SAT Reasoning

- Learn from mistakes *during* search
  - Clause Learning: when DPLL backtracks, learn a concise reason: what went wrong
     → avoid similar 'mistakes' in the future!
  - Extremely powerful in practice



- Restart to avoid heavy-tailed behavior
  - SAT solvers known to sometimes "get stuck" and at other times be very fast... on similar problems, of similar size, etc.
  - Discovery: runtime distribution (across 'similar' instances or even on a single instance) is heavy-tailed → too many runs that can take way longer than the average!
  - Solution: restart computation differently, keep learned clauses!



### Key Advances in SAT Reasoning

- Machine learning to build algorithm portfolios
  - Observation: no single SAT solver is good on every family of instances
  - Features of a given instance can be used to predict, with reasonable accuracy, which solver will work well on it!
  - Solution: design a portfolio solver using ML techniques
    - Based on runtime prediction models
    - Recent work avoid complex models, use k-NN or clustering
- Automatic parameter tuning (generic and instance-specific)
  - SAT solvers are designed with many 'hardwired' parameters
  - Millions of parameter combinations impossible to explore all by hand!
  - Solution: use automatic parameter tuning tools based on local search, genetic algorithms, etc.



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## Advances in Our Understanding

- Phase transitions and hardness profiles
  - Sharp threshold for satisfiability
  - Easy-hard-easy behavior
  - Heavy-tailed runtime phenomenon
    - Inspiring rapid restarts and portfolio methods
- Theoretical Foundations
  - SAT solvers as a "proof system" within Resolution
  - CL adds exponential power to basic DPLL
  - "CL + lots of restarts" as powerful as Resolution!





## Advances in Our Understanding (2)

- Hidden Structure in "real-world" problems
  - Backdoors to satisfiability (and unsatisfiability)

constraint graph



- E.g., setting 16 variables (or even 5 vars) in this logistics planning problem with 843 vars, 7301 constraints simplifies it tremendously!
- Note: such hidden structure not captured by traditional "cutsets"
- Beautiful insights into solution space: clustering structure!
  - Connections with Statistical Physics / Survey Propagation





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### Looking Ahead: *Where is SAT Research Headed?*

- Direction A: getting more out of SAT solvers
  - Minimal/minimum **unsatisfiable cores**: very useful in practice!
  - MAXSAT, weighted MAXSAT
  - Circuit representations (rather than CNF)
- Direction B: tacking problems harder than SAT
  - Near-uniform **sampling** from the solution space
  - **Solution counting** (with relations to probabilistic inference)
    - #P-hard : challenging even to approximate with good confidence bounds
- Direction C: expanding the applicability of SAT technology
  - Pseudo-Boolean SAT (i.e., linear inequalities over Boolean vars)
  - SMT: Satifisiability Modulo Theories (e.g., linear arithmetic, bit-vector operations, uninterpreted functions)



### Summary

- SAT research has been on an exciting 20 year journey
  - From toy problems with 100-200 variables
  - To industrial applications with 1M variables, 5M constraints
- Powerful new techniques and insights, many extensions, new connections with other areas of AI and beyond, strong groups in universities and industry alike
- Handbook of SAT: a great resource!
- Questions/comments/more information: ashish.sabharwal@us.ibm.com

