Towards SHA-3

Christian Rechberger, KU Leuven

Fundamental questions in CS theory

Do oneway functions exist?

Do collision-intractable functions exist?

We don't know.

Do we care?

What we care about: computational properties

For cryptographic hash functions, it should be sufficiently hard to

- find preimages
- find collisions

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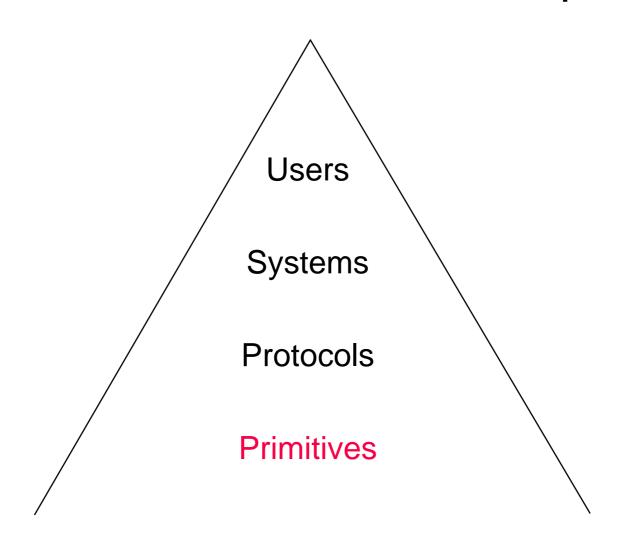
Secure? What properties?

Collision resistance Preimage resistance 2nd preimage resistance Near-collision resistance Pseudorandom generator Pseudorandom function Key derivation function Random oracle

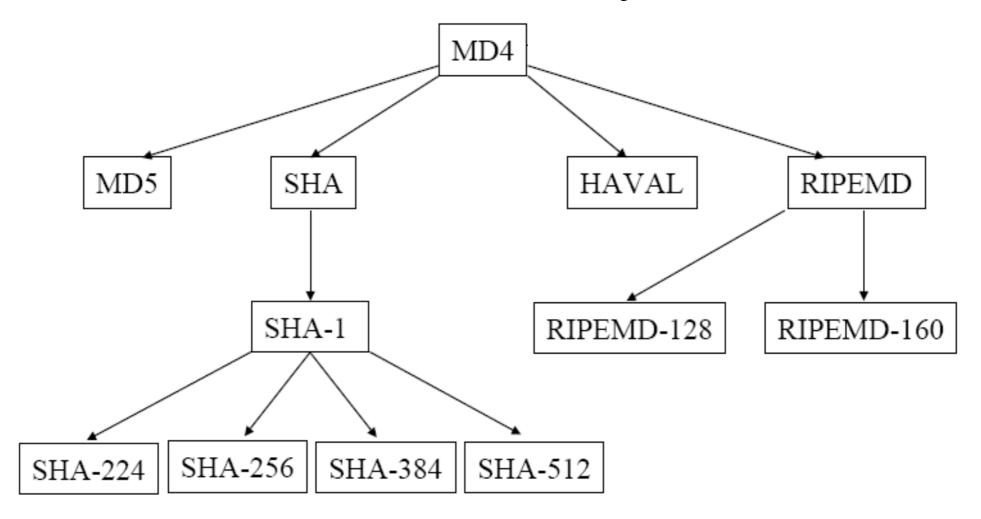


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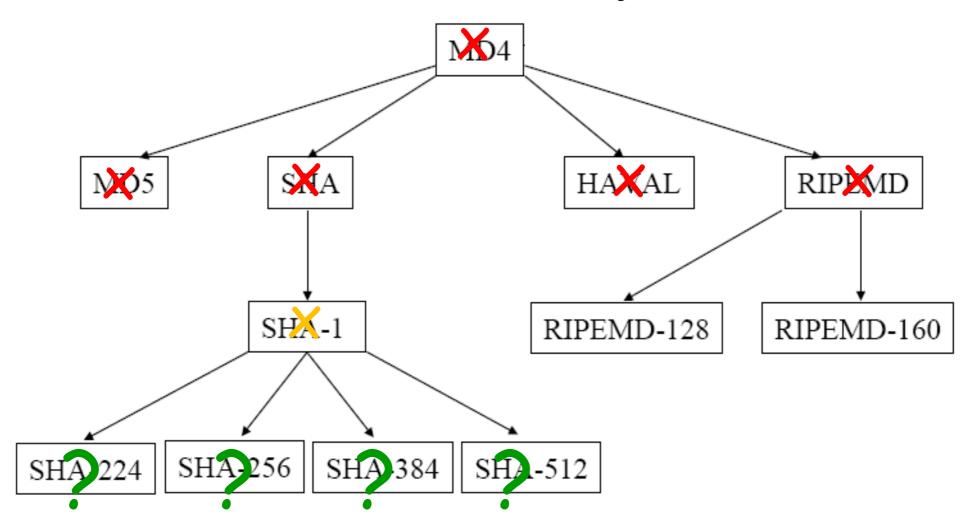
Hash functions as a fundamental primitive



MD4 family



MD4 family



Collisions for reduced SHA-1

40 rounds: Biham, Chen, 2005

58 rounds: Wang, Yu, Yin, 2005

64 rounds: De Cannière, R., 2006

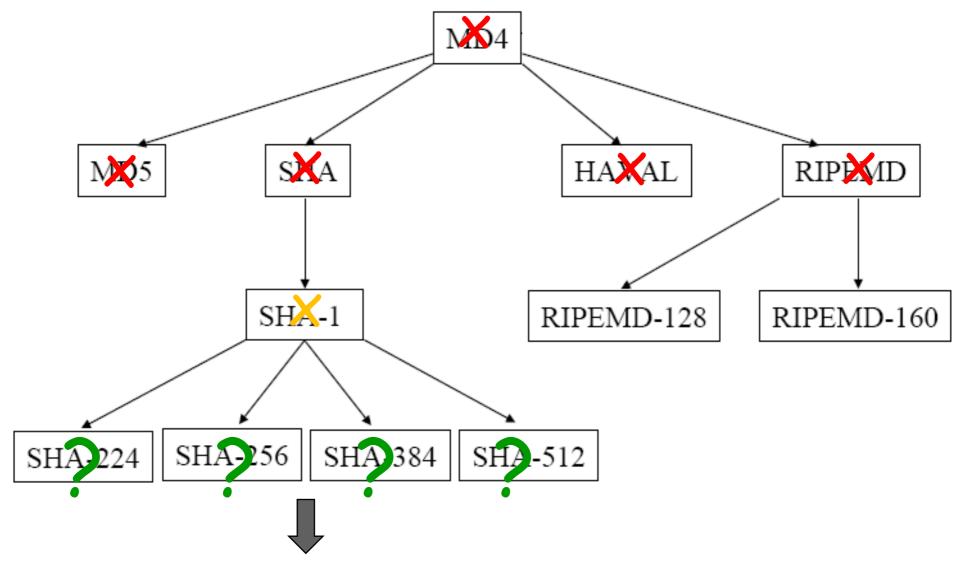
70 rounds: De Cannière, Mendel, R., 2007

Full 80 rounds?

What are the problems

- Too fast?
- Designers too optimistic
- New powerful variants of differential cryptanalysis

Road towards SHA-3



SHA-3 (selected in an open competition)

Design challenges for SHA-3

Faster than SHA-2 on many platforms

More secure than SHA-2, confidence

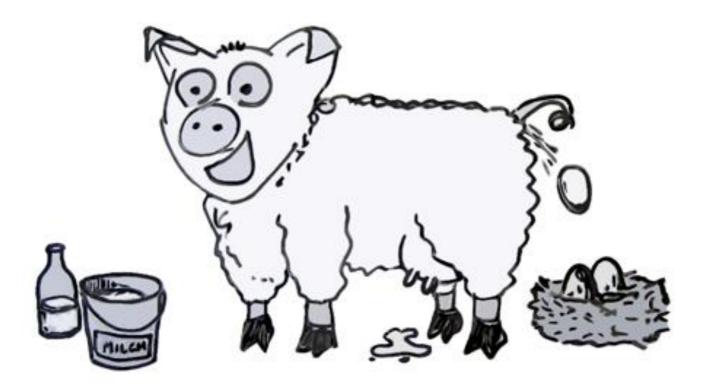
All the properties that you could think of now and in the years to come

Design challenges for SHA-3

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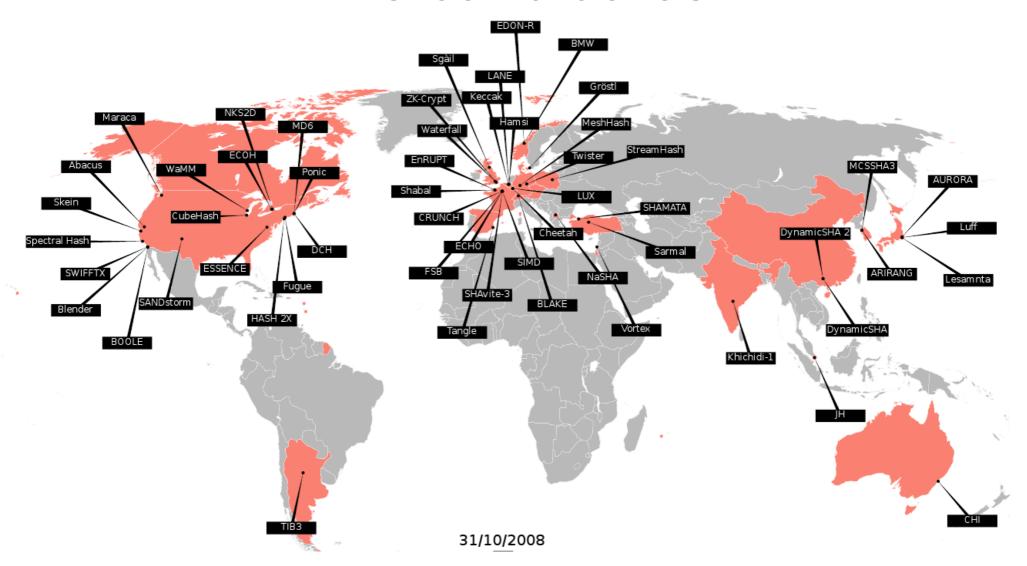
Outline

- Motivation
- SHA-3 competition
- Grøstl and the rebound attack
- SHA-3 candidates through the rebound lens
- Concluding discussions

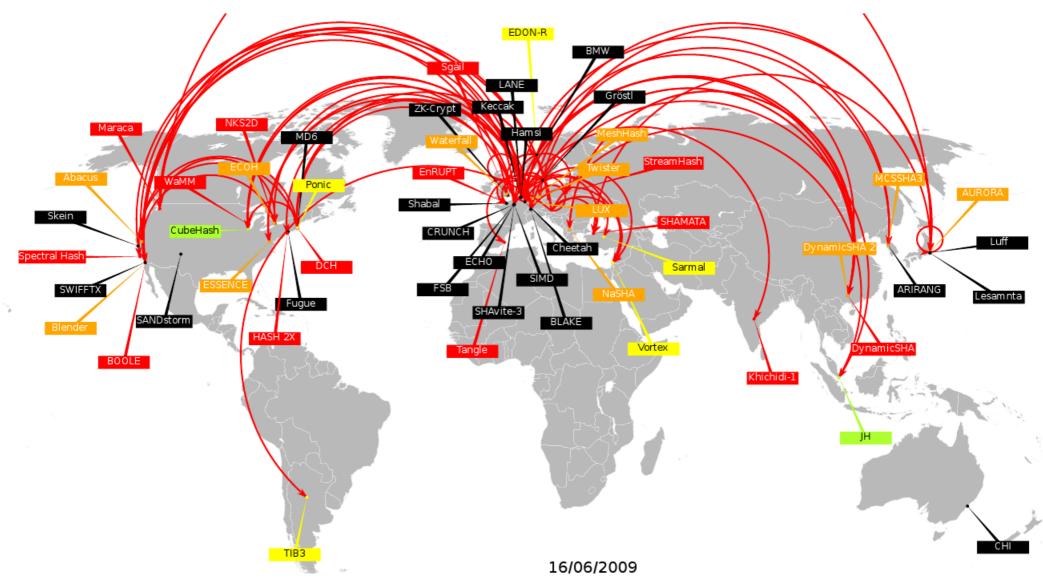
SHA-3 competition

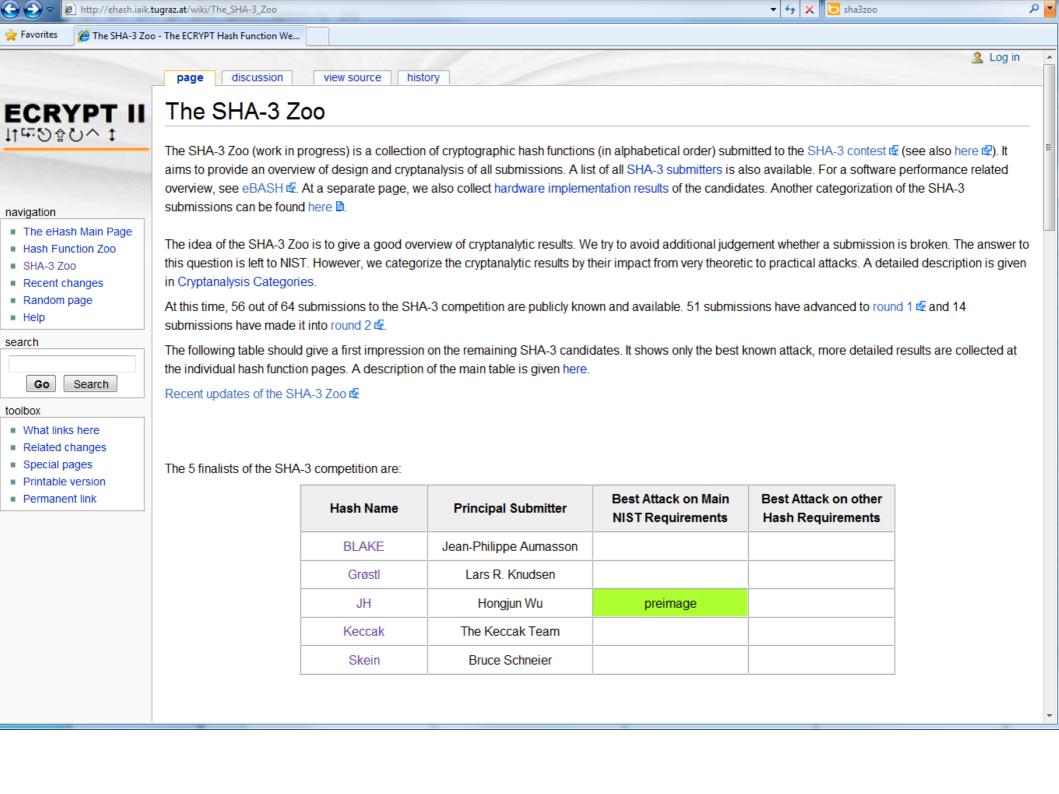
- 2006/2007: NIST drafts requirements and calls for submissions
- 10/2008: 64 submissions, >200 designers
- 12/2008: 51 round-1 candidates announced
- 07/2009: 14 round-2 candidates announced
- 12/2010: Five finalists announced
- Q2 2012: Final selection

The candidates

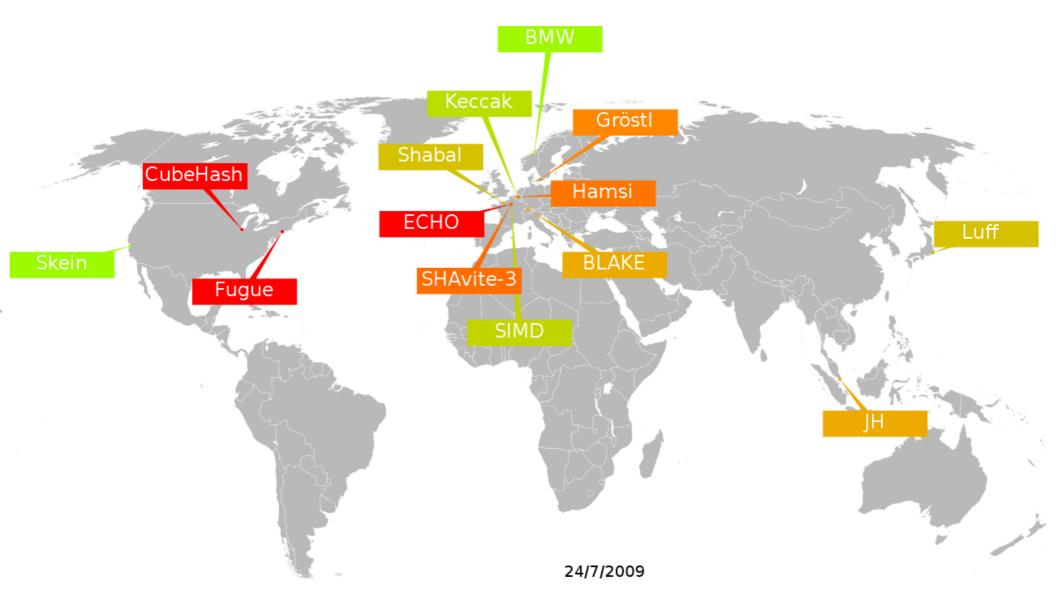


Preliminary cryptanalysis



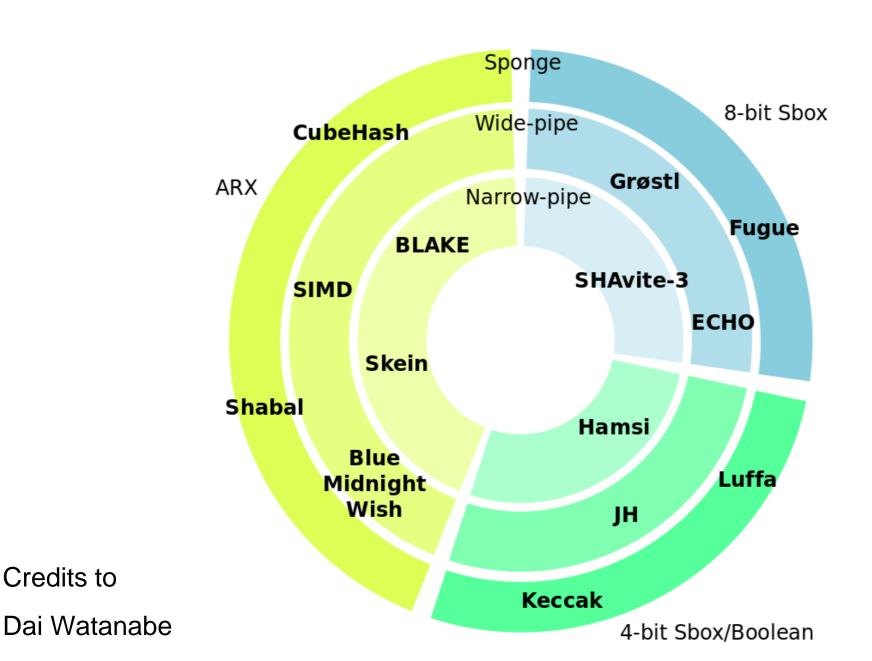


Round-2 candidates



How to categorize them?

How to categorize them?



Credits to

How to compare them?

- Security
- Performance/Implementation costs
 - Software (code size, speed, ...)
 - Hardware (lowest gate count, highest throughput, power consumption characteristics, ...)
 - Side-Channel countermeasures
- Confidence?

Grøstl

Grøstl is inspired by

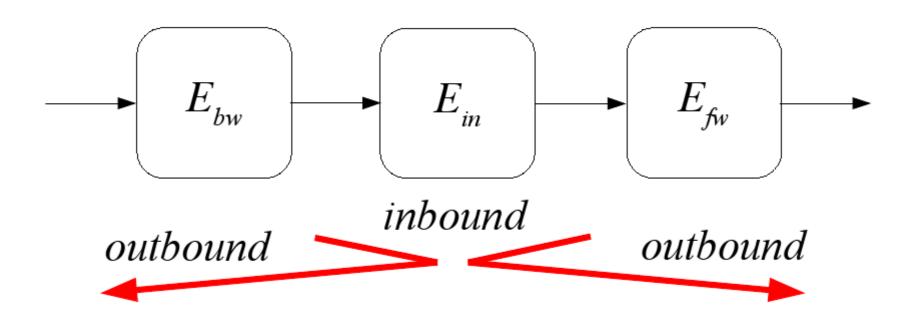
- Rijndael/AES (Daemen, Rijmen, 1997)
- SMASH (Knudsen, 2005)
- Grindahl (Knudsen, R., Thomsen, 2007)

Proofs against differential attacks

Proofs against generic shortcut attacks

Rebound attack

New variant of differential cryptanalysis, FSE 2009 Developed during the design of Grøstl



Origins of the rebound attack

Differential attack, Biham and Shamir, 1989

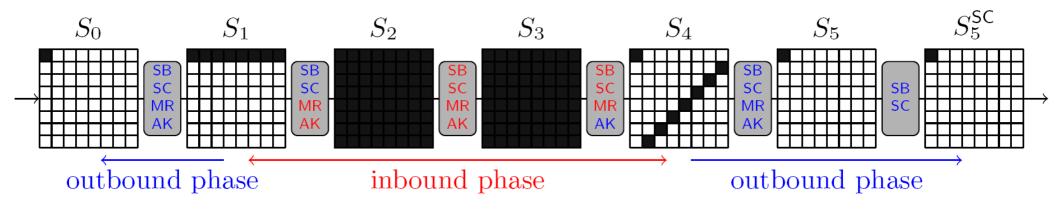
Inside-out approach, Dobbertin 1995, Wagner 1998

Truncated differential, Knudsen, 1994

Original Goal:

Get a good estimate of the security margin of Grøstl

Example of a rebound attack



Within a few months, others became a "victim":

- Twister (round-1 SHA-3 candidate)
- LANE (round-1 SHA-3 candidate)
- Whirlpool (ISO standard, unbroken since 2001)

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Further technical developments

The

Linear solving variant (SAC 2009)

Start-in-the-middle variant (SAC 2009)

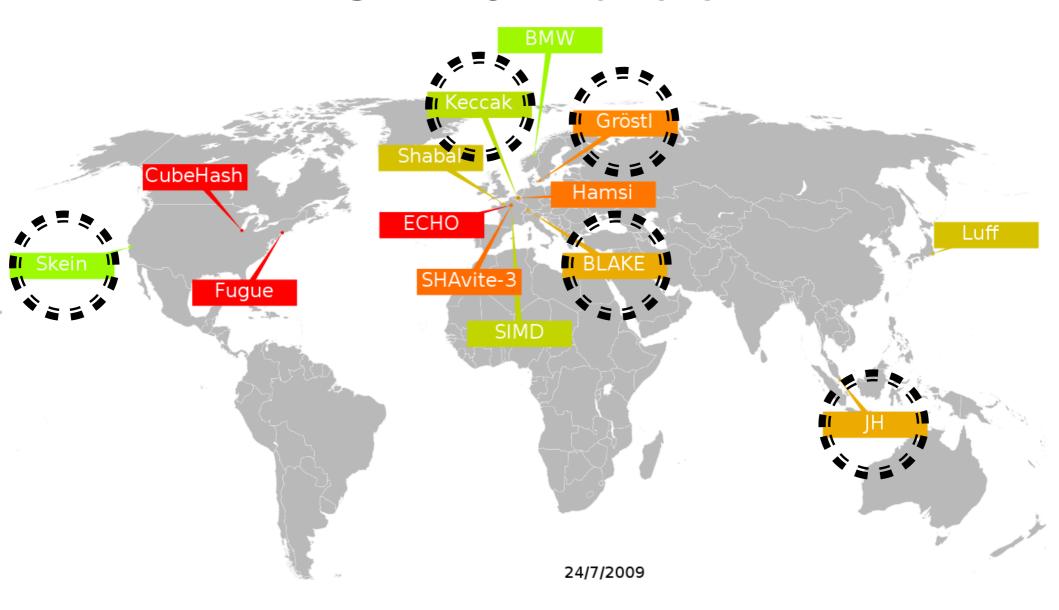
Super(S)box variant (Asiacrypt 2009 and FSE 2010)

Multiple-inbound phase variant (Asiacrypt 2009)

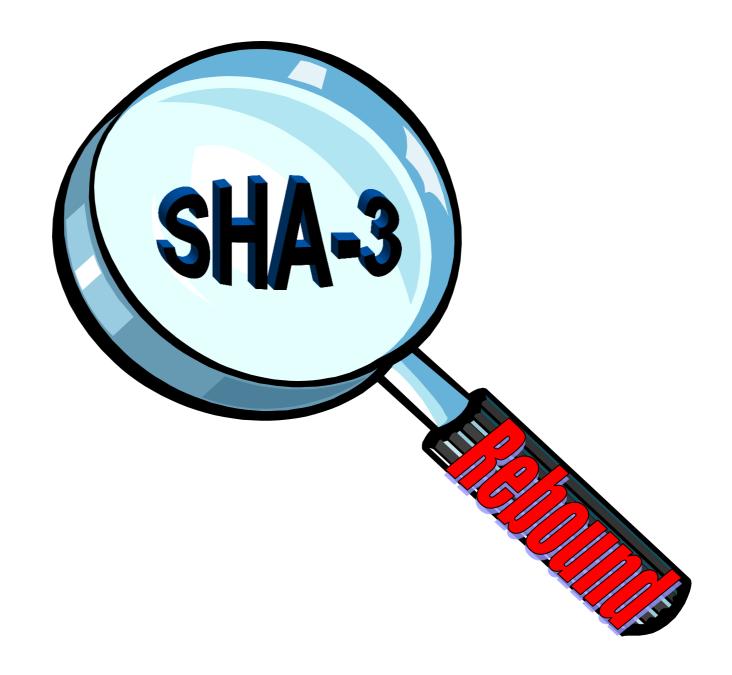
Rotational variant (Asiacrypt 2010)

...of the rebound attack

SHA-3 finalists







SHA-3 round-2 candidates through the rebound lens

4 or 8-bit S-box based

Others

Grøstl

ECHO

JH

Luffa

Shavite-3

Fugue

Hamsi

Skein

BMW

Blake

Cubehash

Keccak

SIMD

Shabal

SHA-3 round-2 candidates through the rebound lens

4 or 8-bit S-box based

Hamsi

Others

Shabal

Skein Grøstl - ノノノノ **ECHO BMW** // JH Blake // Luffa Cubehash Shavite-3 Keccak **Fugue** SIMD

Most recent case: Skein

- $R_{r,i}$
- Recent analysis by Khovratovich, Nikolic, R. in 2010
- Rebound idea for the first time applied to ARX construction
- Results in perspective:
 - 2009: Related-key differential attack: 34 rounds
 - 2010: Rotational attack: 42 rounds
 - New: Rebound rotational attack: 57 rounds

SHA-3 **finalists** through the rebound lens

4 or 8-bit S-box based

Others



SHA-3 finalists in numbers

Geography:

3 from Europe, 1 from Asia, 1 from America

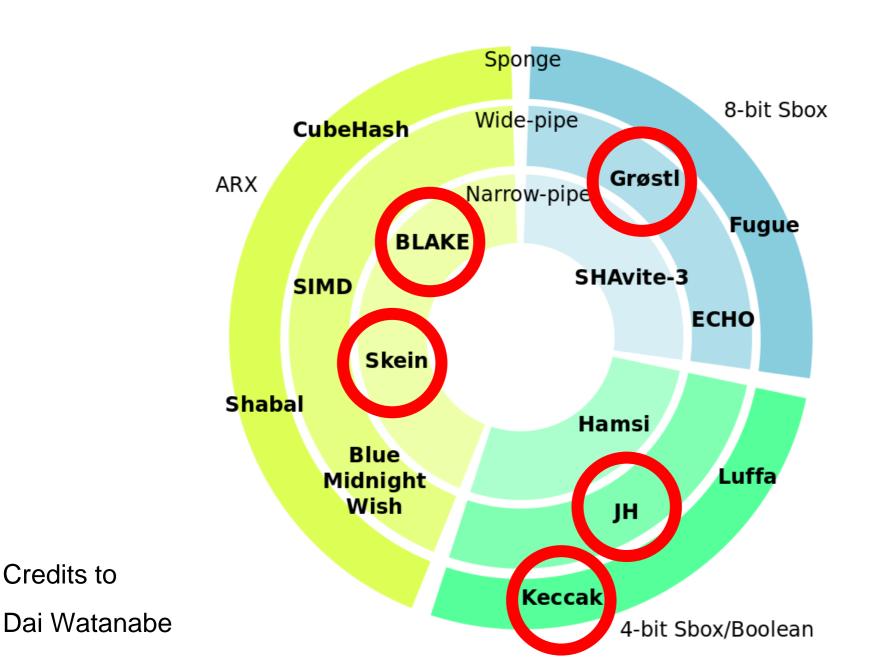
Tweaks:

all 5 got tweaked, 2 got tweaked twice

Team members also AES finalist: 3

Teams that designed a hash function before: 2

How to categorize them?



Credits to

SHA-3 finalists

Compression strategy:

Single Permutation: Blake (with finalization), JH, Keccak

Two Permutations: Grøstl

Large family of permutations (block cipher): Skein

Source of non-linearity:

64-bit: Skein

32/64-bit: Blake

8-bit: Grøstl

4/5-bit: JH

3-bit: Keccak

Conclusion (1/2) Assurance?

Very complicated attacks against MD5 and SHA-1

- (1) Differential trail with complicated carry interactions
- (2) Degrees of freedom utilization for speedup

Level of assurance provided by finalists against this class of attacks:

Blake, Skein: ARX, issues similar to SHA-1/SHA-2

Grøstl: both (1) and (2) done by rebound attacks

JH: (1) and (2) may be possible, open problem

Keccak: seems infeasible

Conclusion (2/2)

Building confidence in a new cryptographic primitive takes time

A lot remains to be done for a final SHA-3 selection by 2012

Upcoming: ECRYPT Hash Workshop 2011, May 19-20, Tallinn

The road ahead

- Application of new cryptanalytic techniques to other areas, examples
 - Internal fixed points:
 - Collision and preimage attack on GOST hash: 2008
 - Key recovery attack on GOST block cipher: 2011
 - Local collisions:
 - Collisions in SHA-0: 1998
 - Related-key attacks on AES: 2009
- New lightweight algorithms, where designers cut corners

Towards SHA-3

Q&A

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Backup slides

Addendum: Grøstl?



Call for input

Name	Country	
Gröstl	Austria	
Hash	USA	
Bubble and squeak	United Kingdom	
Rumbledethumps/Stovies	Scotland	
Colcannon	Ireland	
Bauernfrühstück	Germany	
Stamppot	Netherlands	
Pyttipanna	Finland, Norway, Sweden	
Biksemad	Denmark	
Roupa Velha	Portugal	
Bergerdil	Malaysia	
Ha'DIBaH 'ay'mey 'oQqar je	Qo'noS (Klingon)	