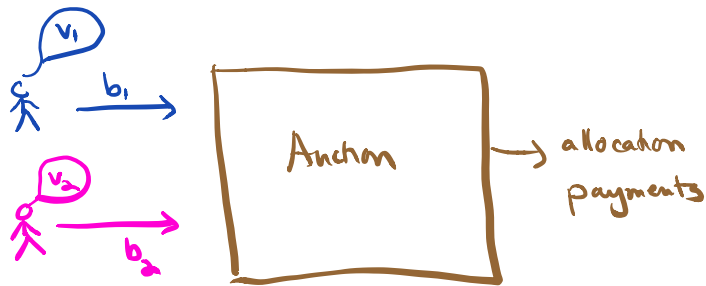


Plan for Today

- Continue with auctions
 - Sponsored Search
 - The VCG auction
 - The FCC Incentive Auction

Last time

v_1, v_2, \dots are private values



1st price auction: winner is highest bidder
winner pays her bid

2nd price auction: winner is highest bidder
winner pays 2nd highest bid

How to bid?

Model: each bidder bids to maximize (expected) utility, given beliefs about others

$$u_i = v_i \mathbb{1}_{\text{receives item}} - p_i \quad \text{utility of bidder } i$$

value payment

Vickrey auction

2nd price auction truthful
also called "strategy proof"
"dominant strategy incentive compatible"
DSIC

\equiv it is a dominant strategy to bid your value

also "individually rational" IR

\equiv bidder's utility guaranteed to be nonnegative

maximizes "welfare"
also called "surplus"

\equiv total happiness of all participants.
 \equiv sum of utilities

$$= \sum_{i=1}^n \underbrace{(v_i \mathbb{1}_{i \text{ wins item}} - p_i)}_{\text{utility of bidder } i} + \underbrace{\sum_{i=1}^n p_i}_{\text{utility of auctioneer}}$$

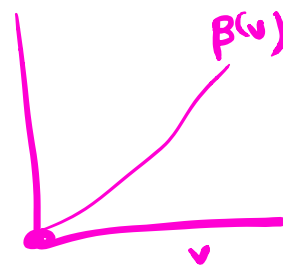
$$= \sum_{i=1}^n v_i \mathbb{1}_{i \text{ wins}}$$

1st price auction: no dominant strategies
 analyzed by assuming each player's
 value was drawn from distribution
 F

Shoved that if F is Uniform $[0, 100]$
 then $\beta(v) = \frac{v}{2}$ is a
Bayes-Nash equilibrium for $n=2$

$$E[\text{utility of bidding } \beta(v_i) \mid \text{other player bids } \beta(v_2)]$$

$$\geq E[\text{utility of bidding } w \mid \text{other player bids } \beta(v_2)]$$



Revenue Equivalence Thm

If $v_i \sim F \forall i$, then for any auction
 s.t. in equilibrium, item allocated to
bidder w/ highest value and payment of a
player with value 0 is 0
 has same expected revenue.

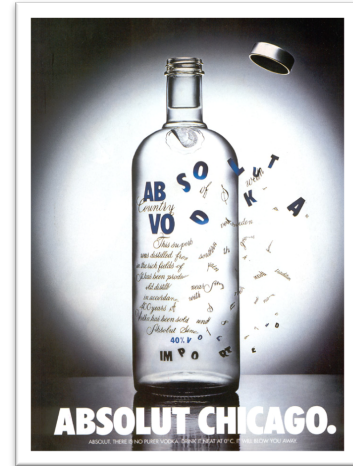
Advertising – how it used to be



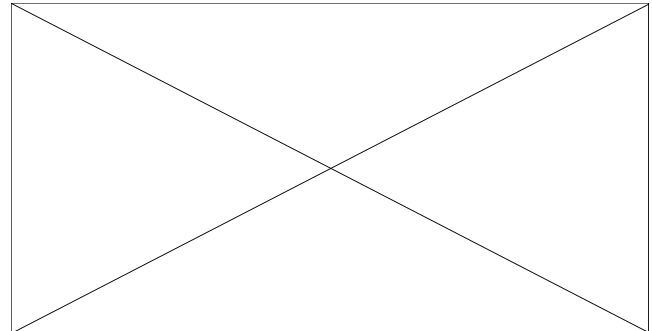
Posters
Billboards



newspapers



magazines
television



Pay-per-impression

Price depends on how many people your advertisement is shown to.

(whether or not they look at it, or care about it)

“Half the money I spend on advertising is wasted; the trouble is, I don’t know which half”

Andrew Wanamaker, advertising pioneer

How is the price determined?

Complicated negotiations with high monthly premiums,

forms a barrier to entry for
small advertisers.

Modern Advertising

On the web,
Many different kinds of ads ...

Sponsored Search Ads



mesothelioma



All News Images Videos Books More Settings Tools



SafeSea

About 7,290,000 results (0.61 seconds)

Ad · www.nationalmesotheliomaclaims.com/

Mesothelioma Claim Center | Get Asbestos Cancer Payments

Learn how the \$30B Asbestos Trust Fund may pay for your asbestos-caused cancer.

Mesothelioma Claims

Asbestos Trust Fund Claims
This \$30 Billion Trust Can Help

Free Info Package

Covers what you need to know
about medical and financial options

Ad · www.mesotheliomabook.com/ (888) 637-6234

Family Hurt by Mesothelioma? | We're Here to Help You

We are the largest firm devoted only to Mesothelioma in the US, visit us today. We come to...

Ad · www.mesothelioma-survivor.com/mesothelioma/treatment (800) 336-0086

I Survived Mesothelioma Cancer - Learn How She Beat the Odds

Learn About Heather's Successful Surgery from Preparation to the Procedure & Recovery.

Ad · www.navy-veterans-mesothelioma.org/Financial/Benefits

Mesothelioma & Asbestos Risk | Navy Vets Asbestos Claims

Important info for Navy Vets. Learn About Mesothelioma Claims

Mesothelioma is a type of cancer that develops from the thin layer of tissue that covers many of the internal organs (known as the mesothelium). The most



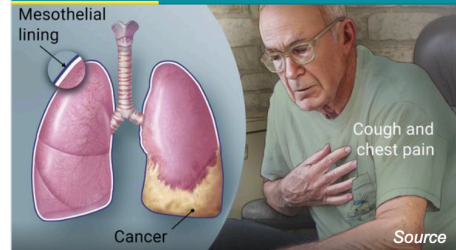
Mesothelioma



ABOUT

SYMPTOMS

TREATMENTS



A tumor of the tissue that lines the lungs, stomach, heart, and other organs.

Very rare

Fewer than 20,000 US cases per year

- Treatment can help, but this condition can't be cured
- Requires a medical diagnosis
- Lab tests or imaging always required

Price determined by auction – per keyword

The image shows a screenshot of a search engine results page for the query "computer science". The page includes a search bar, navigation links, and search results. Three blue rectangular boxes with white text are overlaid on the page to represent ad slots: "Slot One" and "Slot Two" are positioned over the top of the first search result, and "Ad A", "Ad B", and "Ad C" are positioned over the right side of the page. The search results include a Wikipedia entry for "Computer science", a gateway to "Computer Science" resources, and a department page from Carnegie Mellon University.

Slot One

Slot Two

Ad A

Ad B

Ad C

“Most people don’t realize that all that money comes in pennies at a time”
Hal Varian, Google Chief Economist

How are these ads different than the ads in the offline media?

In many cases: Pay-per-click

Model: per keyword auction.

Bidders: advertisers who have a standing bid on keyword.

Auction offer some # $k > 1$ of slots.

Slots not identical

CTR clickthrough rates

c_j : Prob of a click on the j^{th} slot. $j=1..k$

v_i : advertiser's value for a click.

q_i : CTR rate due to quality of ad. (all qualities are same)

overall: if put bidder's ad in slot j

$$\text{Prob of click} = c_j \cdot q_i$$

Value to bidder i if gets his ad placed
in slot j

exp $\underline{v_i \cdot c_j}$

exp. utility $v_i \cdot c_j - p_i \cdot c_j$

Generalized Second Price GSP

- allocate to slots in order of decreasing bid
- on a click to slot j , payment is next highest bid.

Model

Click-through rates

Value per click

$c_1 = 1$



value.
(bids)

7

PPC

6

utility

1

$c_2 = 0.4$



6

1

$(6-1) \cdot 0.4$

$c_3 = 0$

No slot



1

Slots

Advertisers

Example

$$c_1 = 1$$

$$c_2 = 0.99$$



in this example,
if bidder 1 bids below 9
then his utility will
be higher than it was
bidding truthfully.

Truthful auction: VCG auction

chooses outcome to maximize social welfare.

$$v_1 > v_2 > \dots$$

social welfare of an allocation

$$= \sum_{i=1}^k v_i c_{s(i)}$$

slot i is assigned to

$$c_1 > c_2 > c_3 > \dots$$

Social welfare maximizing allocation

$$\sum_{i=1}^k v_i c_i$$

change each bidder the "externality" her
presence imposes on others

Bidder 1:

What would SW be w/o bidder 1 in auction

$$\sum_{i=2}^{k+1} v_i c_{i-1}$$

SW with bidder 1

$$\sum_{i=2}^k v_i c_i$$

VCG payment for bidder 1

charge bidder 1 less others incur because of her

$$\sum_{i=2}^{k+1} (c_{i-1} - c_i) v_i$$

$$c_{k+1} = 0$$

$$c_{i-1} > c_i$$

VCG payment for bidder j

$$\sum_{i=j+1}^{k+1} (c_{i-1} - c_i) v_i$$

VCG auction for sponsored search

Get a bid from each bidder

Relabel bids so $b_1 > b_2 > \dots > b_n$

for $i = 1..k$ assign bidder i to i 's slot.

for $i = 1..k$, for each click, charge bidder i

$$p_i = \frac{1}{c_i} \sum_{j=i+1}^{k+1} b_j (c_{j-1} - c_j)$$

$$c_i p_i = \left(\sum_{j=i+1}^{k+1} b_j (c_{j-1} - c_j) \right)$$

$$b_{i+1} \left(\frac{c_i - c_{i+1}}{c_i} \right) + b_{i+2} \left(\frac{c_{i+1} - c_{i+2}}{c_i} \right)$$

PPC

Then VCG is truthful, IR, and welfare maximizing

Proof of truthfulness

Fix other bids, let b_j denote j 'th highest bid except me.

slot k	$c_k b_k$
k-1	$(c_{k-1} - c_k) b_{k-1} + c_k b_k$
i	$\sum_{j=i}^k (c_j - c_{j+1}) b_j$ ←

my value.

v

Δ value as max from i to slot i-1

Δ price

$$(c_{i-1} - c_i) v$$

$$(c_{i-1} - c_i) b_{i-1}$$

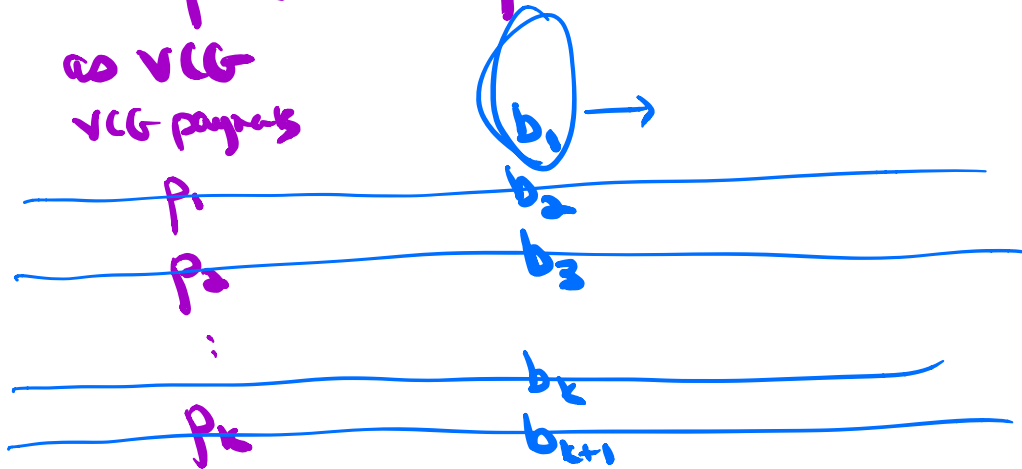
want to do this if $v > b_{i-1}$

$$b_i > v > b_{i+1}$$

↑

other bids
 $b_1 > b_2 > \dots$

If valuations & CTRs, the GSP auction has an equilibrium of same allocation & payments



no sweeping revenue comparison between
 GSP & VCG

Facebook

- what bidders can bid on.
various "events"
- placement, size & format of ads
- evaluation of quality of ads

$$v_i = c_j$$

more appropriate models

bidder has value for each event type
difficult for advertisers to figure out CTRs

- high quality outcomes
- advertisers happy, bidding easy

VCG (fully general)

- n participants.
- finite set of outcomes of auction Ω ←
- each agent i has private valuation $v_i(w)$ $\forall w \in \Omega$
- $v_i(w) = (\text{value of some event}) \cdot \text{Pr}(\text{event occurs in outcome } w)$

VCG

choose outcome

$$\in \underset{w \in \Omega}{\text{argmax}} \sum_{i=1}^n v_i(w)$$

Let this outcome be (w^*)

$$p_i = \max_{w \in \mathcal{N}} \sum_{j \neq i} b_j(w) - \sum_{j \neq i} b_j(w^*)$$

$$u_i = v_i(w^*) - p_i$$

$$= v_i(w^*) + \sum_{j \neq i} b_j(w^*) - \max_{w \in \mathcal{N}} \sum_{j \neq i} b_j(w)$$

choose $\underline{b_i(w)}$ $\forall w \in \mathcal{N}$
to maximize

$$v_i(w^*) + \sum_{j \neq i} b_j(w^*) \quad (*)$$

$w^* \in$ argmax
 w $b_i(w) + \sum_{j \neq i} b_j(w)$

User interface for bidding.

- allow bidders ^{to bid} on events they
- allow bidders to specify budget

Computational requirements

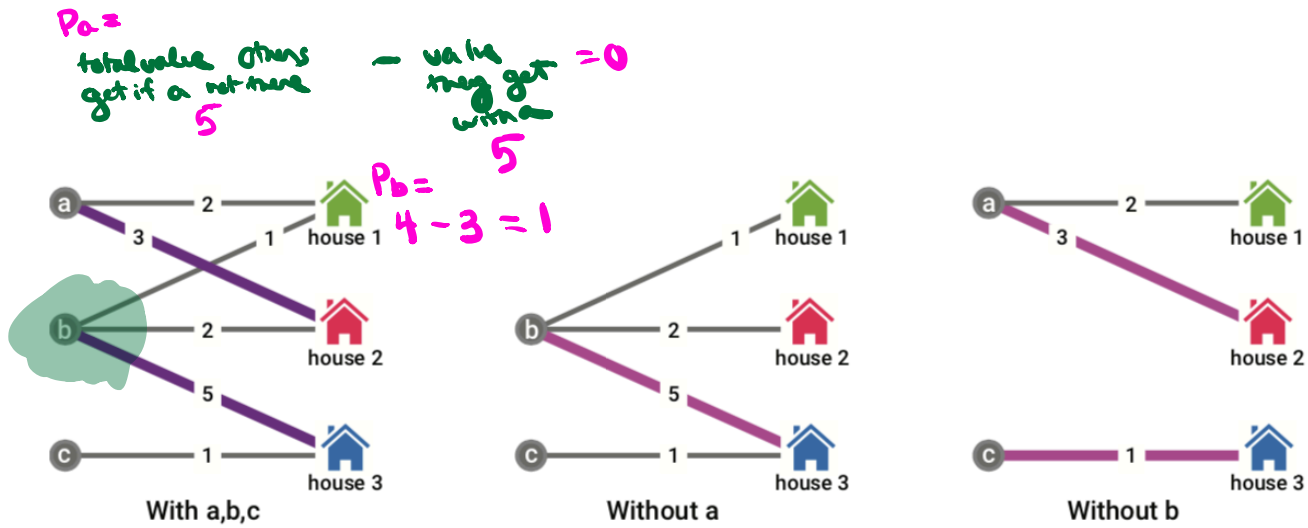


FIGURE 16.3. The label on the edge from i on the left to j on the right is the value v_{ij} that employee i has for a yearly lease of house j (say in thousands of dollars). The VCG mechanism allocates according to purple shaded edges. The payment of bidder a is 0 since in his absence house 3 is still allocated to bidder b . The payment of bidder b is 1 since in his absence the allocation is as follows: house 2 to bidder a and house 3 to bidder c , and therefore the externality he imposes is 1.