Collaborative Verification of Information Flow for a High-Assurance App Store

Michael D. Ernst, **René Just**, Suzanne Millstein, Werner Dietl*, Stuart Pernsteiner, Franziska Roesner, Karl Koscher, Paulo Barros, Ravi Bhoraskar, Seungyeop Han, Paul Vines, and Edward X. Wu

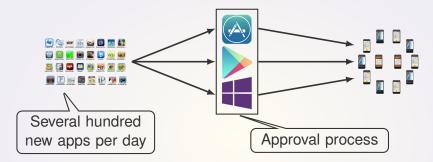


University of Washington *University of Waterloo

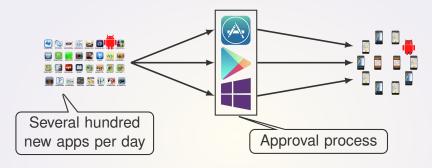
November 6, 2014



Current commercial app stores

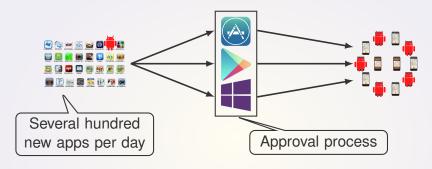


Current commercial app stores



Problem: Every major app store has approved malware!

Current commercial app stores



Problem: Every major app store has approved malware!

Best-effort solution: Malware removed when encountered

High-assurance app stores

Needed in multiple domains

- Government app stores (e.g., DoD)
- Corporate app stores (e.g., financial sector)
- App stores for medical apps

Require stronger guarantees

Verified absence of (certain types of) malware

Verification is costly

- Effort is solely on app store side
- Analyst needs to understand/reverse-engineer the app

High-assurance app stores

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Our solution: Collaboratively verify absence of malware

Our focus: Information-flow malware

App



Sudoku

Permissions

App



Permissions

App



Permissions

Sudoku



Read location Internet

App



Permissions

Sudoku



Read location Internet

Example: Information-flow malware

App



Permissions
Read location



Information flow



Read location Internet

Location → Internet

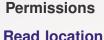
Introduction

Example: Information-flow malware

App



Sudoku





Internet, Kanada

Information flow







Example: Information-flow malware

App



Permissions

Read location



Information flow



Read location Internet





Read location Internet

Example: Information-flow malware

App



Permissions

Read location Internet <a>A



Information flow



Read location Internet





Read location Internet

Example: Information-flow malware

App



Permissions



Information flow









Read location Internet

Location → **Internet**

Introduction

Example: Information-flow malware

App



Permissions Read location



Information flow







Read location Internet

Internet





App



Permissions

Read location



Internet, Kanada

Information flow



Internet





Read location Internet



Location → BadGuy.com

App



Permissions

Read location Internet <u>—</u>



Information flow



Camera



Read location Internet







Location →





App



Sudoku



Camera



Read location Internet

Permissions



Read location

Internet

Information flow

Prevent malware using an information flow type-system



Location → Internet.





Location → BadGuy.com

Approach Evaluation Conclusion

Approach: Overview

Collaborative verification model

Leverage but don't trust the developer

Information Flow Type-checker (IFT)

- Finer-grained permission model for Android
- False positives and declassifications
- Implicit information flow

Evaluation

- Effectiveness: Effective for real malware in real apps
- Usability: Low annotation and auditing burden

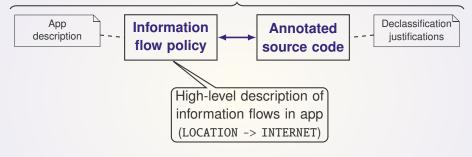
Collaborative verification model

Developer provides



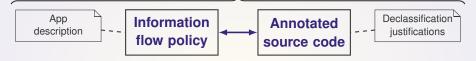
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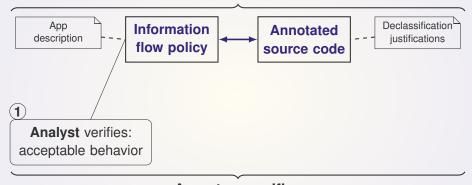
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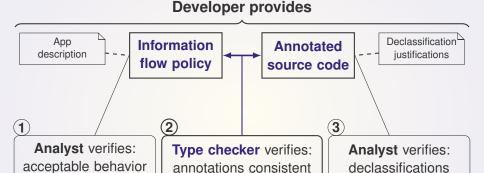
Collaborative verification model

Developer provides Declassification aaA Information **Annotated** description justifications flow policy source code **Analyst** verifies: Type checker verifies: acceptable behavior annotations consistent

Collaborative verification model

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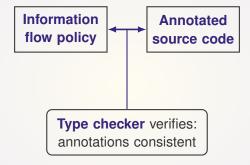
Collaborative verification model



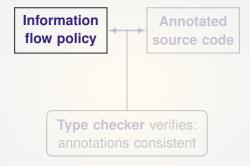
App store verifies

Developer and analyst do tasks that are easy for them

Verification of information flow



Verification of information flow



Information flow policy

High-level description of permitted information flows

READ_SMS -> INTERNET

READ_CLIPBOARD -> DISPLAY

USER_INPUT -> CALL_PHONE

ACCESS_FINE_LOCATION -> INTERNET(maps.google.com)

Information flow policy

High-level description of permitted information flows

	Source	flows to	Sink
•	READ_SMS	->	INTERNET
	READ_CLIPBOARD	->	DISPLAY
	USER_INPUT	->	CALL_PHONE
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Sources and Sinks

► Default Android permissions (145)

Not sufficient to model information flow!

Information flow policy

High-level description of permitted information flows

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Sources and Sinks

- Default Android permissions (145)
- Additional sensitive resources (28)

Information flow policy

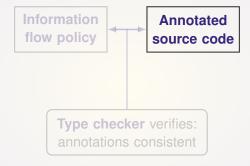
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Sources and Sinks

- Default Android permissions (145)
- Additional sensitive resources (28)
- Parameterized permissions

Verification of information flow



Information flow types: sources and sinks

@Source Where might a value come from? **@Sink** Where might a value flow to?

Information flow types: sources and sinks

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```
Android API
```

```
void sendToInternet(String message);
String readGPS();
```

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Information flow types: sources and sinks

@Source Where might a value come from? **@Sink** Where might a value flow to?

Android API

void sendToInternet(@Sink(INTERNET)String message);
String readGPS();

Information flow types: sources and sinks

@Source Where might a value come from? **@Sink** Where might a value flow to?

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Android API

void sendToInternet(@Sink(INTERNET)String message);

String readGPS();

(From Location)
```

Information flow types: sources and sinks

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Android API

void sendToInternet(@Sink(INTERNET)String message);
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void sendToInternet(@Sink(INTERNET)String message);
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```

```
App code
```

```
String loc = readGPS();
sendToInternet(loc);
```

uction Approach Evaluation Conclusion

Information flow types: sources and sinks

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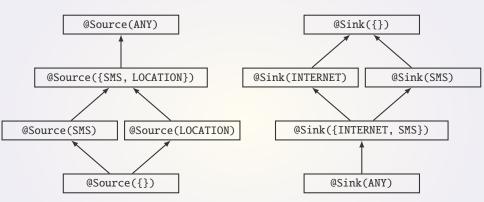
API annotations are pre-verified

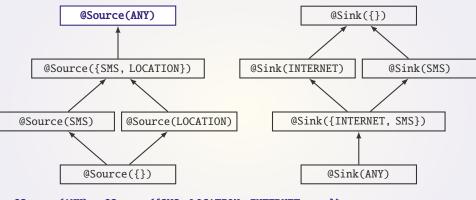
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void sendToInternet(@Sink(INTERNET)String message);
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App code

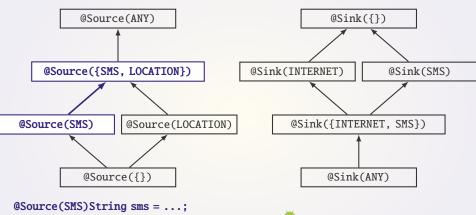
Developer annotations are not trusted

```
@Source(LOCATION)@Sink(INTERNET)String loc = readGPS();
sendToInternet(loc):
```

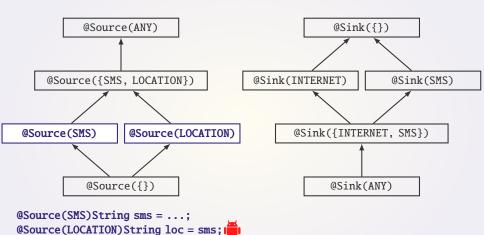


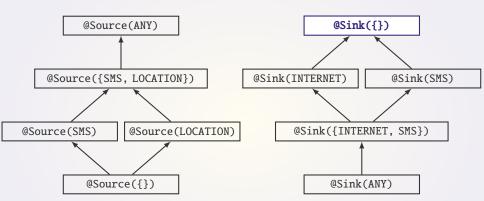


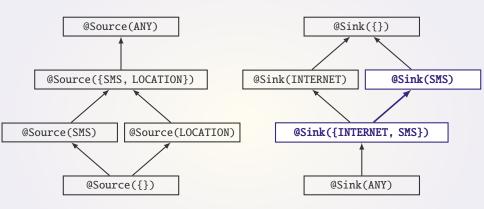
@Source(ANY) = @Source({SMS, LOCATION, INTERNET, ...})



@Source({SMS, LOCATION})String smsLoc = sms;



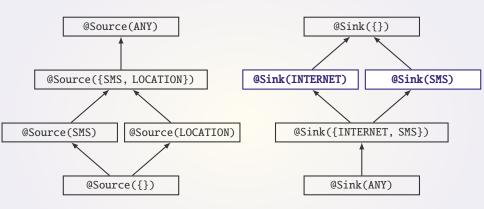




@Sink({INTERNET, SMS})String toInetSms;

@Sink(SMS)String toSms = toInetSms;



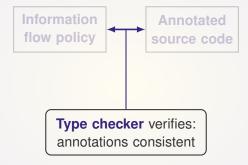


@Sink(SMS)String toSms;

@Sink(INTERNET)String toInet = toSms;



Verification of information flow



oduction Approach Evaluation Conclusion

Information Flow Type-checker (IFT): Overview

Guarantees of type-checking

- Annotations are consistent with code (type correctness)
- 2. Annotations are consistent with flow policy

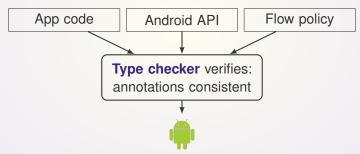
Type checker verifies: annotations consistent

oduction Approach Evaluation Conclusion

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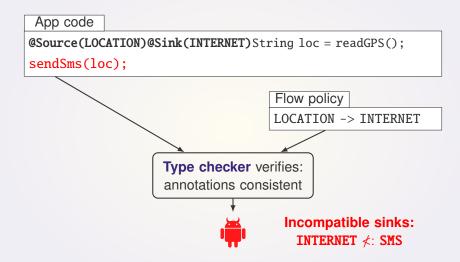


No undisclosed information flows in app

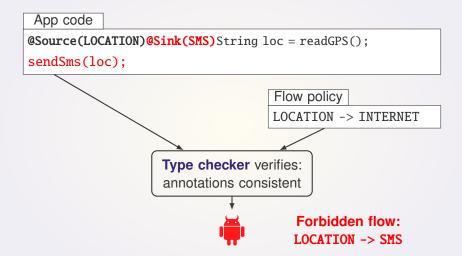
App code @Source(LOCATION)@Sink(INTERNET)String loc = readGPS(); sendToInternet(loc); Flow policy LOCATION -> INTERNET Type checker verifies: annotations consistent

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                                     Flow policy
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App code
@Source(LOCATION)@Sink(SMS)String loc = readGPS();
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                                     Flow policy
                                    LOCATION -> INTERNET
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```



False positives and declassifications

```
App code
@Source({LOCATION, SMS})String [] array;
array[0] = readGPS();
array[1] = readSMS();
@Source(LOCATION)String loc = array[0];
```

False positives and declassifications

```
App code
@Source({LOCATION, SMS})String [] array;
array[0] = readGPS();
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@Source(LOCATION)
@Source(LOCATION)String loc = array[0];
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False positives and declassifications

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@Source(LOCATION)String loc = array[0];

@Source(LOCATION, SMS)
```

False positives and declassifications

```
App code |
@Source({LOCATION, SMS})String [] array;
array[0] = readGPS();
array[1] = readSMS();
@SuppressWarnings("flow") // Safe: returns location data
@Source(LOCATION)String loc = array[0];
```

Declassifications

- Developer can suppress false-positive warnings
- App store employee verifies each declassification

oduction Approach Evaluation Conclusio

Reducing false positives

Flow sensitivity

Type refinement with intra-procedural data flow analysis

```
App code

@Source({LOCATION, SMS})String value;
if (...) {
  value = readSMS();
    ...  value: @Source(SMS)
}
...  value: @Source({LOCATION, SMS})
```

oduction Approach Evaluation Conclusion

Reducing false positives

Flow sensitivity

Type refinement with intra-procedural data flow analysis

Context sensitivity

► Polymorphism (e.g., String operations, I/O streams, etc.)

```
App code
```

```
@Source({LOCATION, SMS})String value = ...;
String substring = value.substring(0,8);
```

Returns @Source({LOCATION, SMS})

roduction Approach Evaluation Conclusion

Reducing false positives

Flow sensitivity

Type refinement with intra-procedural data flow analysis

Context sensitivity

Polymorphism (e.g., String operations, I/O streams, etc.)

Indirect control flow

- Constant value propagation
- Reflection analysis
- Intent analysis

Implicit information flow

App code

```
@Source(USER_INPUT)long creditCard = getCard();
long i=0;
while (true) {
  if (++i == creditCard) {
    sendToInternet(i);
  }
}
```

Implicit information flow

```
App code

@Source(USER_INPUT)long creditCard = getCard();
long i=0;
while (true) {
   if (++i == creditCard) {
        sendToInternet(i);
   }
}
Card number implicitly leaked
```

Classic approach (Denning and Denning, CACM'77)

- Taint all computations in dynamic scope
- Over-tainting may lead to taint explosion

Implicit information flow

```
App code

@Source(USER_INPUT)long creditCard = getCard();
long i=0;
while (true) {
  if (++i == creditCard) {
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  }
}
USER_INPUT -> CONDITIONAL
```

Our approach: Prune irrelevant conditions

- Add additional sink CONDITIONAL
- Type-checker warning for conditions with sensitive source

Implicit information flow

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}
USER_INPUT -> CONDITIONAL
```

Our approach: Prune irrelevant conditions

- Add additional sink CONDITIONAL
- Type-checker warning for conditions with sensitive source

Analyst must manually verify

- Analyst is aware of context
- ► No need to analyze dynamic scope for irrelevant conditions (e.g., null checks, malicious conditions, or trigger)

Evaluation: Overview

Are our permission model and type system effective?

- Adversarial Red Team challenge
- Evaluation of effectiveness for real malware

Is our approach effective and efficient in a timeconstrained set up?

- Control team study
- Comparison of effectiveness and efficiency to control team

Is our verification model applicable for real-world apps?

- Usability study with annotators and auditors
- Evaluation of annotation and auditing burden

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Apps are not pre-annotated

Adversarial Red Team challenge

Setup

- 5 independent Red Teams
- ► 72 Android apps (47 malicious with information-flow malware)
- 8,000 LOC and 12 permissions per app on average

Adversarial Red Team challenge

Setup

- 5 independent Red Teams
- ▶ 72 Android apps (47 malicious with information-flow malware)
- 8,000 LOC and 12 permissions per app on average

Results for 47 malicious apps



- ▶ 96% overall detection rate 4% require modeling of information flow paths (LOCATION -> ENCRYPT -> INTERNET)
- ► 60% of apps require our finer-grained sources and sinks

Control team study

Setup

- Control team using dynamic and static analysis tools
- ► 18 Android apps (13 malicious)
- 7,000 LOC and 16 permissions per app on average

Control team study

Setup

- Control team using dynamic and static analysis tools
- ▶ 18 Android apps (13 malicious)
- 7,000 LOC and 16 permissions per app on average

Results



Setup

- 2 groups acting as annotators and auditors
- ▶ 11 Android apps (1 malicious)
- 900 LOC and 12 permissions per app on average

Setup

- 2 groups acting as annotators and auditors
- 11 Android apps (1 malicious)
- 900 LOC and 12 permissions per app on average

Annotation burden

- 96% of type annotations are inferred
- Annotations required: 6 per 100 lines of code
- Annotation time: 16 minutes per 100 lines of code

Most time spent on reverse engineering

Declassifications

- ▶ 50% of apps had no declassifications
- On average 3 declassification per 1,000 lines of code

IFT's features effectively reduce false positives

Declassifications

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- On average 3 declassification per 1,000 lines of code

IFT's features effectively reduce false positives

Auditing burden

- Overall review time: 3 minutes per 100 lines of code
- 35% of time: review the flow policy
- ▶ 65% of time: review declassifications & conditionals

Only 23% of conditionals needed to be reviewed

Related work: Information flow

Jif (Myers, POPL'99)

- A security-typed language (incompatible Java extension)
- Supports dynamic checks and focuses on expressiveness

FlowDroid (Arzt et al., PLDI'14), SuSi (Rasthofer et al., NDSS'14)

- FlowDroid propagates sources and sinks found by SuSi
- SuSi classifies Android API methods using machine learning

IFT makes static verification of Android apps practical

- Finer-grained sources and sinks at type level
- Compiler plug-in using standard Java type annotations

Related work: Collaborative verification model

Verifying browser extensions

- ▶ IBEX (Guha et al., S&P'11)
 - Verification of Fine (ML dialect) against complex policies
- Lerner et al., ESORICS'13
 - Verification of private browsing using annotated JavaScript

IFT verifies information flow in Android apps using a high-level flow policy

Automated policy verification

- Crowd-sourcing (Agarwal & Hall, MobiSys'13)
- Natural language processing (Pandita et al., USENIX'13)
- Clustering (Gorla et al., ICSE'14)

Could aid manual verification of flow policies

Conclusions

Collaborative verification model

- Low overall verification effort for developer and app store analyst
- ► IFT combined with other analyses

Developer provides App description Information flow policy Annotated source code Annotated source code Type checker verifies: acceptable behavior Analyst verifies: annotations consistent Analyst verifies: declassifications

App store verifies

Information Flow Type-checker (IFT)

- Context and flow-sensitive type system
- Fine-grained model for sources and sinks
- High-level information flow policy

App code Android API Flow policy Type checker verifies: annotations consistent

Evaluation

- Detected 96% information-flow malware
- Low annotation and auditing burden
- Low false-positive rate



https://www.cs.washington.edu/sparta