Diesel Highlights

Purely object-oriented language
- all data are instances of classes
- all operations & control structures
  via dynamically dispatched method calls

Multiply dispatched method calls

Closures, a.k.a. first-class lexically-nested functions

Static type system
- including fancy polymorphic types

Module system
- namespace management & encapsulation

Type-safe

Garbage-collected

Class declarations

To declare a class, use a `class` declaration, e.g.:

```
class shape;
class rectangle isa shape;
class rhombus isa shape;
class square isa rectangle, rhombus;
```

A class can have zero, one, or many superclasses
- multiple inheritance supported

A class *doesn't* declare any of its fields or methods;
these are separate top-level declarations

Can add new superclasses to existing classes externally,
e.g. in separate source files!

```
class printable;
extend class shape isa printable;
```

Each class defines a new type
- a subclass is a subtype

Field declarations

To declare the instance variables of a class, use `field` declarations

E.g.:

```
var field center(s:shape):point {new_origin()}
field width(r:rectangle):num;
field height(r:rectangle):num { r.width }
```

Fields are declared separately from classes

- the field is related to its "containing" class via the type of the field's argument
  - each object of that type (or subtype) stores a value for the field
- can add new fields to classes externally!

Must say `var` for assignable field

- immutable by default

A field can have a default initial value

- can be an expression, e.g. computing the field's initial value from the initial values of other fields

Function declarations

To declare a new top-level procedure, constructor, or method, use a `fun` declaration, e.g.:

```
fun new_point(x:num, y:num):point { ... }
fun new_origin():point { new_point(0,0) }
fun rect_area(r:rectangle):num {
  r.width * r.height }
fun move_rect(r:rectangle,
              new_center:point):void {
  r.center := new_center;
}
```

Functions are declared separately from classes

- receiver argument (if any) is explicit
- constructors have explicit names
- can add new functions to classes externally!

Can have different functions with same name but different numbers of arguments (a kind of static overloading)

A function body is a sequence of zero or more statements, followed by an optional result expression (`void` if absent)

Method declarations

Override an existing function for particular kinds of arguments using a `method` declaration

- method has same name and number of arguments as overridden function
- one or more formals' types declared using `@` instead of `:`
  - method applies only to run-time arguments whose dynamic class is an instance of the `@` type, called the **specializer**
- can override a method, too
  - more specific `@` types override less specific ones

E.g.:

```
fun resize(s:shape,dw:num,dh:num):shape {...}
method resize(r@rectangle,
              dw:num, dh:num):rectangle {...}
method resize(s@square,
              dw:num, dh:num):rectangle {...}
method resize(c@circle,
              dw:num, dh:num):shape {...}
```

Method body same syntax as function body

Abstract classes and functions

A class can be `abstract`

- can't have direct instances

E.g.:

```
abstract class shape;
```

A function declared for an abstract class need not have a body

- must be overridden by some method for every concrete subclass

E.g.:

```
fun resize(s:shape,dw:num,dh:num):shape;
-- must have resize methods for all concrete
-- subclasses of shape
```

Multiple dispatching

Can have multiple `@` specialized formals in a method
$\Rightarrow$ multiple dispatching

E.g.:

```
fun =(s1:shape, s2:shape):bool { false }
method =(r1@rectangle, r2@rectangle):bool {..}
method =(c1@circle, c2@circle):bool { ... }
```

All arguments treated uniformly

- any can be specialized, or not
- any number can be specialized
- specialization is always based on *dynamic* argument class, not *static* argument type

Method lookup rules

When invoke a function with some arguments
(a.k.a. send a message),
need to identify the right method to run

- consider a function with a body as a method with no `@`

Algorithm:

1. find set of *applicable* methods in invoked function
   - a method is applicable if, for each `@`*C* formal, the dynamic class of the corresponding argument is equal to or a subclass of *C*
   - if no applicable methods: report msg-not-understood error
2. select unique *most-specific* applicable method
   - a method is at least as specific as another if its specializers are uniformly at least as specific as the other's
   - if no uniquely most specific method: report msg-ambiguous error
3. run it

Static typechecking checks for these method lookup errors

Constraints on method types

Method argument and result types must conform to overriddee function/method's

- method's `@` formal types should be more specific than overridee's [*covariant*]
  - otherwise, wouldn't override!
  - safe, since tested dynamically via method lookup
- method's `:` formal types should be as general as (typically, the same as) the overridee's [*contravariant*]
- method's result type can be more specific than overridee's [*covariant*]

E.g.:

```
fun resize(s:shape,dw:num,dh:num):shape;
method resize(r@rectangle,
              dw:num, dh:num):rectangle {...}
method resize(s@square,
              dw:num, dh:num):rectangle {...}
method resize(c@circle,
              dw:num, dh:num):shape {...}
```

Constraints ensure that if a call to a function typechecks, then no matter what method is invoked, its formal and result types will be compatible with the caller's expectations

Object creation

Create new instances of a class by evaluating `new` expressions

- can provide initial values for fields, or rely on fields' defaults, which are evaluated separately for each object

E.g.:

```
fun new_rectangle(w:num, h:num):rectangle {
  new rectangle {
    -- center gets default value
    width := w, height := h } }

fun new_square(w:num):square {
  new square {
    -- center gets default value
    -- height derived from width by default initializer
    width := w } }
```

Good programming style:
encapsulate all `new` expressions inside functions

Unlike traditional constructors, these functions:

+ can cache and return previously created objects
+ can create instances of different classes based on e.g. args
– cannot inherit field initialization code

Object declarations

Can declare one-of-a-kind objects using `object` declarations

- similar syntax to class declarations
- also can specify initial values for fields

E.g.:

```
object unit_square isa square { width := 1 };
```

Can inherit from and specialize on named objects just like classes

- cannot do that for anonymous objects created with `new`

Can reference named objects directly just like global variables

- cannot do that for classes

Expressions and statements

Constants, e.g.: `3`, `-4`, `5.6`, `"hi there\nbob"`, `'a'`

- all values are regular, first-class objects,
  e.g. `3` is an instance of `prim_int` class

Vector constructors, e.g.: `[]`, `[3+x, y*z, f(x)]`

- vectors are regular, first-class objects too

`new` expressions, e.g.: `new rectangle {...}`

Identifiers, e.g.: `x`, `joe_Bob_17`, `true`, `void`

- reference local var, formal, global var, or named object

Variable declaration statements, e.g.:

```
let w := y * z;
let var x:int := w + f(w);
```

- variables must be initialized at declaration
- assignable variables and globals should be given types

Assignment stmts, e.g.: `x := y * f(z);`

- cannot assign to formals or non-`var` locals/globals

Messages

Use standard procedure-call syntax to invoke a function with zero or more arguments:

```
start_prog()
center(r)
set_center(r, c)
draw(r, window, loc)
```

Infix & prefix syntax:

```
x + - y << z ** q!i
```

- any sequence of punctuation chars is a legal infix or prefix message name
- implement with normal functions & methods
- can specify precedence & associativity

Syntactic sugar supports standard "dot notation":

```
r.center   ⇔  center(r)
r.set_center(c)   ⇔  set_center(r, c)
r.draw(window, loc)   ⇔  draw(r, window, loc)
```

Syntactic sugar for `set_` messages to look like assignments:

```
r.center := c;   ⇔  set_center(r, c);
```


Accessing fields

Access fields solely by sending messages

- to read a field named `f` of object `o`, send `f` message to `o`, to invoke "get accessor" implicit method
  - syntactic sugar: `o.f`
- to update a (var) field named `f` of object `o` to new value `v`, send `set_f` message to `o` and `v`, to invoke "set accessor" implicit method
  - syntactic sugar: `o.f := v;`

Syntactic sugar makes accessing fields by messages syntactically "natural"

- can access methods as if they were fields, too

Allows fields to be reimplemented as methods & vice versa, and allows fields to be overridden with methods & vice versa, without rewriting callers

No explicit accessor methods or "properties" needed

Resends

In overriding method, can invoke overridden method, e.g.:

```
class visible_rectangle isa rectangle;
method resize(r@visible_rectangle,
              dw:num, dh:num):rectangle {...}
  let new_r := resend(r, dw, dh);
  r.undisplay;
  new_r.display;
  new_r }
```

Can use to resolve ambiguities, e.g.:

```
class square isa rectangle, rhombus;
fun area(s:shape):num;
method area(r@rectangle):num { ... }
method area(r@rhombus):num { ... }
method area(s@square):num {
  resend(s@rectangle) }
```

(Like Java's `super`)

Closures

First-class function objects

Used for:

- standard control structures (`if`, `while`, `&`, `|`, etc.)
- iterators (`do`, `find`, etc.)
- exception handling (`fetch`, `store`, etc.)

Syntax

- **`&`**`(`*formals*`){` *zero or more stmts*`;` *result expr* `}`, e.g.:

  `&``(i:int,j:int){ let x := i + j; x*x }`

  - **`&`**`(int,int):int` is the type of this closure
- if no formals, can omit **`&`**`()`, e.g.: `{ print("hi"); }`

Examples of use:

```
if(i > j, { i }, { j })
[3,4,5].do(&(x:int){ x.print; })
table.fetch(key, { error("key is absent") })
```

Invoke closure by sending `eval` with right number of arguments

```
let cl := &(i:int){ i.print_line; };
...
eval(cl, 5);
```

Non-local returns

Can exit a function/method early via
a non-local return from a nested closure

```
{ ...; ^ result }
{ ...; ^ }
```

Example:

```
fun find_index(s:string,
               value:char,
               if_absent:&():int
              ):int {
  s.do_associations(&(i:int, v:char){
    if(v = value, { ^ i });
  });
  eval(if_absent) }

fun find_index(s:string, value:char):int {
  find_index(s, value,
             { error("not found") }) }
```


Parameterization

Can parameterize classes & functions

- functions can be implicitly parameterized using ` notation

Can provide upper bounds for parameter types

```
abstract class collection[T];
abstract class table[Key <= comparable[Key],
                     Value]
            isa collection[Value];
class array[Value] isa table[int,Value];

fun fetch(t:table[`Key,`Value], k:Key):Value;
fun find_key(
     t:table[`Key, `Value<=comparable[Value]],
     val:Value,
     if_absent:&():Key):Key {
  t.do_associations(&(k:Key,v:Value){
    if(v = val, { ^ k });
  });
  eval(if_absent) }
```

Explicit type parameters must be provided by client

Implicit formal type parameters inferred from argument types

Special types

`any`

- type of anything (akin to `Object` in Java)

`void`

- special object & type used for functions that don't return a useful result

`none`

- result type of functions that do not return normally, e.g. `error`, `loop`, `exit` argument closures

`dynamic`

- like `any`, but disables static checking
- the default type for formals & result, if explicit types omitted

`type1` **`&`** *`type2`*

- anything that is both a *`type1`* and a *`type2`*

`type1` **`|`** *`type2`*

- anything that is either a *`type1`* or a *`type2`*

Modules

Can wrap declarations in a `module` declaration, for encapsulation and namespace management

- mark named declarations as `public`, `protected` (the default), or `private` to control access outside the module
  - var fields have two names, with independent access control
- different modules can declare same names to mean different things

Can reference visible module contents using *`module`***`$`***`id`*

```
module Shapes {
  public abstract class shape;
  public get protected set
    var field center(s:shape):Points$point;
  public fun area(s:shape):num;
  fun shape_helper(s:shape):num { ... }
}

let s:Shapes$shape := ...;
let a:num := Shapes$area(s);
```

Module imports

Can `import` a module to give importing scope direct access to imported module's public names

E.g.:

```
module Shapes {
  import Points;
  public abstract class shape;
  public get protected set
    var field center(s:shape):point;
  public fun area(s:shape):num;
  fun shape_helper(s:shape):num { ... }
}
import Shapes;

let s:shape := ...;
let a:num := area(s);
```


Module extensions

Can declare that one module `extends` another module, to `import` other module and gain access to its `protected` things

```
module Rectangles;
  public extends Shapes;
  public class rectangle isa shape;
  public field width(r:rectangle):num;
  public field height(r:rectangle):num;
  public fun new_rectangle(w:num, h:num
                           ):rectangle {...}
  fun rect_area(r:rectangle):num { ... }
  method area(r@rectangle):num { r.rect_area }
end module Rectangles;
```

More on modules

Can write any of

```
module Name { ... }
module Name; ... end module Name;
module Name; ... end module;
module Name; ... <end of file>
```

interchangeably

Can declare a module within a module

- nested module declaration specifies its visibility

Can add new declarations to an existing module's body externally, e.g. in separate source files!

```
extend module Shapes {
  public fun =(s1:shape, s2:shape):bool {false}
}
extend module Rectangles {
  method =(r1@rectangle, r2@rectangle):bool...
}
```


Programs and files

A Diesel program is a file containing declarations and statements

- declarations visible throughout their scope
  - (mutual) recursion without forward declarations or header files
- statements executed in textual order
  - no `main` function necessary
  - access command-line arguments using `argv` object in standard library

To spread programs across multiple files, use `include` declarations

- an included file can include other files
- by default, Diesel programs implicitly include `prelude.diesel` to get standard library

E.g.

```
include "shapes.diesel";
```

Some standard control structures

```
if(test, { then });
if(test, { then }, { else }) -- returns a value
if_false(...);

test & { other_test } -- short-circuiting
test | { other_test } -- short-circuiting
not(test)

loop({ ... ^ ... });

while({ test }, { body });
while_false(...);
until({ body }, { test });
until_false(...);

exit(&(break:&():none){
  ... eval(break); ... });
exit_value(&(break:&(resultType):none){
  ... eval(break, result); ... });
loop_exit(...);
loop_exit_value(...);
loop_exit_continue(&(break,continue){...});
loop_exit_value_continue(&(brk,cont){...});
```


Standard operations for all objects

Printing:

`print_string` -- *return printable string*

`print, print_line` -- *print out* `print_string`

Comparing:

`==, !==` -- *compare objects'* ***identities***

`=, !=` -- *compare* `comparable` *objects'* ***values***

Some standard classes and objects

```
bool
  true, false

num
  integer
    int         -- limited-precision integers
    big_int     -- arbitrary-precision integers
  float
    single_float
    double_float

character
  char        -- ascii
  unicode_char

pair, triple, quadruple, quintuple

mb[type]    -- type | absent
  absent

file        -- unix files
```


Standard collection classes and functions

```
abstract collection[T]
  length, is_empty, non_empty
  do, includes, find, pick_any
  copy

abstract unordered_collection[T]
  sets and bags

abstract ordered_collection[T]
  linked lists

abstract table[Key,Value]
  hash tables, association lists

abstract indexed[Value]
                isa table[int,Value],
                    ordered_collection[Value]
  arrays, vectors, strings

abstract sorted_collection[T <= ordered]
  binary trees, skiplists
```

Unordered collections

```
abstract unordered_collection[T]
                isa collection[T]
  add, add_all
  remove, remove_some, remove_any, remove_all

abstract bag[T] isa unordered_collection[T]

class list_bag[T] isa bag[T]
  new_list_bag[T]

abstract set[T] isa unordered_collection[T]
  union, intersection, difference
  is_disjoint, is_subset

class list_set[T] isa set[T]
  new_list_set[T]
class hash_set[T <= hashable] isa set[T]
  new_hash_set[T]
class bit_set[T] isa set[T]
  new_bit_set[T]
```


Ordered collections

```
abstract ordered_collection[T]
              isa collection[T]
  do (over 2-4 ordered collections in parallel)
  add_first, add_last, remove_first/_last
  || (concatenate)
  flatten (for collections of strings)

abstract list[T] isa ordered_collection[T]
  first, rest
  set_first, set_rest

class simple_list[T] isa list[T]
  cons
object nil[T] isa simple_list[T]
```

- cannot add in place to simple lists

```
class m_list[T] isa list[T]
  new_m_list[T]
```

Keyed tables

```
abstract table[Key,Value]
                 isa collection[Value]
  do_associations, includes_key, find_key
  fetch, !
  store, set_!, fetch_or_init
  remove_key, remove_some_keys, remove_all

class assoc_table[K,V] isa table[K,V]
  new_assoc_table[K,V]

class hash_table[K <= hashable,V]
                           isa table[K,V]
  new_hash_table[K,V]
```


Indexed collections

```
abstract indexed[T] isa ordered_collection[T],
                        table[int,T];
  first, second, ..., fourth, last
  set_first, ..., set_last
  includes_index, find_index
  pos, contains
  swap, sort
```

Fixed length (no `add`, `remove`):

```
class vector[T] isa indexed[T]
class i_vector[T] isa vector[T]
  new_i_vector[T](len, filler)
  new_i_vector_init[T](len, &(i){ value })
  new_i_vector_init_from[T](c, &(c_i){value})
class m_vector[T] isa vector[T]
  new_m_vector[_init[_from]][T](...)
```

Extensible:

```
class array[T] isa indexed[T]
  new_array[T]()
  new_array[_init[_from]][T](...)
```

`new_`*X*`_init_from` is like ML's `map`

Strings

```
abstract string isa indexed[char]
  to_lower_case, to_upper_case
  copy_from
  has_prefix, has_suffix
  remove_prefix, remove_suffix
  pad, pad_left, pad_right
  parse_as_int, parse_as_float
  print
```

Fixed length:

```
abstract vstring isa string
```

```
class i_vstring isa vstring
  new_i_vstring(len, filler)
  new_i_vstring_init(len, &(i){ value })
  new_i_vstring_init_from(c, &(c_i){value})
```

- `"..."` is an `i_vstring`

```
class m_vstring[T] isa vstring
  new_m_vstring[_init[_from]](...)
```


Other collections

```
class stack[T] isa m_list[T]
  push, pop, top
  new_stack[T]
```

```
class queue[T] isa m_list[T]
  enqueue, dequeue
  new_queue[T]
```

```
class histogram[T] isa hash_table[T,int]
  new_histogram[T]
  increment
```