# Figures of Merit* 

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## 1 In the Spotlight

Watch out for the fellow who talks about putting things in order! Putting things in order always means getting other people under your control.

- Denis Diderot, 1796

In Theoretical Computer Science it is customary to alphabetize the names of coauthors on the title page of collaborative publications. For those of us whose name will never appear before the phrase et al., this can be cause for concern.

Given the custom of alphabetizing, and under the assumption that surnames are distributed uniformly over the alphabet, one would expect to see such collaborations as

$$
\begin{array}{|l|}
\hline \text { Beeri, Mendelzon, Sagiv, and Ullman [4]. }
\end{array}
$$

On the other hand, something seems amiss with the collaboration

$$
\begin{array}{|l|l|}
\hline \text { Vishkin and Wigderson [36] . }
\end{array}
$$

After all, if Uzi Vishkin chose a coauthor at random, what is the probability that that coauthor's surname would occur later in the alphabet than Vishkin? Armed with this example, we make the following definition.

Definition: In a collaboration of $X_{0}<X_{1}<\cdots<X_{k}$, the spotlight factor of $X_{0}$ is

$$
\star\left(X_{0}\right)=\left(1-. X_{0}\right)^{k} . *
$$

In words, the spotlight factor is the probability that $k$ coauthors chosen independently at random will all have surnames later in the alphabet than $X_{0}$; the lower the spotlight factor, the more impressive the achievement of the first author.

[^1]| $\star\left(X_{0}\right)$ | $X_{0}$ et al. |
| :--- | :--- |
| 0.1889 | Ogden, Riddle, \& Rounds [20] |
| 0.1719 | Vishkin \& Wigderson [36] |
| 0.1640 | Paul, Seiferas, \& Simon [24] |
| 0.12685 | Wong \& Yao [37] |
| 0.12680 | Wood \& Yap [38] |
| 0.1214 | Karmarkar, Karp, Lipton, Lovász, \& Luby [14] |
| 0.0919 | Ruzzo, Simon, \& Tompa [28] |
| 0.0851 | Schwartz, Sharir, \& Siegel [30] |
| 0.0664 | Paul, Pippenger, Szemerédi, \& Trotter [21] |
| 0.0255 | Santoro, Sidney J., Sidney S., \& Urrutia [29] |

Table 1: The Top Ten Spotlight Factors

Example 1 In the aforementioned collaboration of Vishkin and Wigderson [36],

$$
\begin{aligned}
\star(\text { Vishkin }) & =1-. \text { Vishkin } \\
& =1-\left(\frac{22}{27^{1}}+\frac{9}{27^{2}}+\frac{19}{27^{3}}+\cdots\right) \\
& \approx 0.1719
\end{aligned}
$$

In contrast, the less exotic collaboration of Beeri, Mendelzon, Sagiv, and Ullman [4] achieves a more modest spotlight factor:

$$
\begin{aligned}
\star(\text { Beer }) & =\left(1-. \text { Beer } i^{3}\right)^{3} \\
& =\left(1-\left(\frac{2}{27^{1}}+\frac{5}{27^{2}}+\frac{5}{27^{3}}+\cdots\right)\right)^{3} \\
& \approx 0.7756 \square
\end{aligned}
$$

Table 1 lists the top ten spotlight factors in the Theoretical Computer Science community. A two-author collaboration that exceeds the record of 0.0255 would have to be at least Zippel and Zuck, which would achieve $\star($ Zippel $) \approx 0.0238$. At the other extreme of the alphabet, Adleman would have to collaborate with 84 other coauthors in order to achieve a spotlight

| Paul, Pippenger, Szemerédi, \& Trotter [21] |
| :--- |
| Paul, Seiferas, \& Simon [24] |
| Paul \& Reischuk [23] |
| Paul, Prauß, \& Reischuk [22] |
| Paul \& Tarjan [25] |
| Paul, Tarjan, \& Celoni [26] |

Table 2: A Sampling of Wolfgang Paul's Collaborations
factor of 0.0246, and he would have to ensure that the list didn't include Aanderaa or Abelson.

Note that the only researcher who appears twice in the top ten of Table 1 is Wolfgang Paul. As corroboration of the statistical significance of the spotlight factor, a random sampling of Paul's publications is listed in Table 2.

This section closes with a particularly reprehensible form of the spotlight phenomenon, namely, those professors who willfully choose their graduate student advisees with surnames later in the alphabet than theirs, hoping to cash in at the time of future collaborations. Some examples are given in Table 3.

There is one known instance in which a resourceful Ph.D. student named Yehuda outspotlighted his advisor Shimon Even. When it came time to publish the results of their collaboration, Even announced his inevitable intention of being first author. Yehuda responded by legally changing his name to Bar-Yehuda [3].

## 2 Out of the Spotlight

One of the most pernicious effects of haste is obscurity.

- Samuel Johnson, 1752

After completing our careful study of the spotlight factor, we were settling back to rest on our laurels when the collaboration

| Advisor | Advisee |
| :--- | :--- |
| Paul | Prauß |
|  | Reisch |
| Reischuk |  |
| Rollig |  |
| Schnitger |  |$|$| Shamir |
| :--- |
|  |
| Snir |
| Upfal |

Table 3: Reprehensible Form of Spotlight Factor

> Brassard \& Crépeau [8]
came to light. Something seemed biased about this collaboration, but in this case the fault cannot be Brassard's, who is likely to be first author in any collaboration. Closer inspection revealed Crépeau to be the culprit. After all, if Crépeau chose a coauthor at random, what is the probability that that coauthor's surname would upstage his by as little as Brassard's did? Armed with this example, we make the following definition.

Definition: In a collaboration of $X_{0}<X_{1}<\cdots<X_{k}$, the coefficient of obliviousness of $X_{i}$ is

$$
\dot{\iota}\left(X_{i}\right)=\left(. X_{i}-. X_{0}\right)^{i},
$$

for $1 \leq i \leq k$.
In words, the coefficient of obliviousness is the probability that $i$ coauthors chosen independently at random will all have surnames that precede $X_{i}$ as narrowly as does $X_{0}$; the lower the coefficient, the more oblivious $X_{i}$ is to fame.

Example 2 In the aforementioned collaboration of Brassard and Crépeau [8],

$$
\dot{¿}(\text { Crépeau })=. \text { Crépeau }-. \text { Brassard }
$$

$$
\begin{aligned}
& =\frac{1}{27^{1}}+\frac{0}{27^{2}}+\frac{4}{27^{3}}-\frac{3}{27^{4}}+\cdots \\
& \approx 0.0372
\end{aligned}
$$

In contrast, our standard example of Beeri, Mendelzon, Sagiv, and Ullman [4] achieves a more modest coefficient of obliviousness:

$$
\begin{aligned}
\dot{i}(\text { Ullman }) & =(. \text { Ullman }-. \text { Beer } i)^{3} \\
& =\left(\frac{19}{27^{1}}+\frac{7}{27^{2}}+\frac{7}{27^{3}}-\frac{5}{27^{4}}-\cdots\right)^{3} \\
& \approx 0.3635 \square
\end{aligned}
$$

Table 4 lists the top coefficients of obliviousness in the Theoretical Computer Science community.

## 3 The Fundamental Theorem

There's a sucker born every minute.

- attributed to Phineas T. Barnum

Comparing Tables 1 and 4, it is apparent that many of the collaborations that occur in the former also occur in the latter, and that in all cases the constants are much smaller in the latter. This leads us to the

Fundamental Theorem of Collaborative Sociology: If $X_{0}<X_{1}<\cdots<X_{k}$ agree to collaborate, $X_{k}$ 's obliviousness to fame exceeds $X_{0}$ 's appetite for fame.

Proof: $\left(. X_{k}-. X_{0}\right)^{k}<\left(1-. X_{0}\right)^{k}$

## 4 Monotone Erdös Number

These studies raise the question of a natural variation of the well known "Erdös number". Define a directed graph $G=(V, E)$, where $V$ is the set of all researchers, and $(u, v) \in E$ if and only if there is some publication in which $u$ appears earlier in the list of coauthors than $v$.

| $\dot{\ell}\left(X_{k}\right)$ | al. et $X_{k}$ |
| :--- | :--- |
| 0.0372 | Brassard \& Crépeau [8] |
| 0.0369 | Borodin \& Cook [6] |
| 0.0364 | Vishkin \& Wigderson [36] |
| 0.0334 | Ladner \& Lynch [16] |
| 0.0278 | Bhatt \& Cai [5] |
| 0.0185 | Alon \& Azar [2] |
| 0.0169 | Garey, Graham, \& Johnson [10] |
| 0.0151 | Ogden, Riddle, \& Rounds [20] |
| 0.0149 | Kung \& Leiserson [15] |
| $0.0148 *$ | Paul, Seiferas, \& Simon [24] |
| 0.0095 | Aggarwal \& Anderson [1] |
| 0.0086 | Shamir \& Snir [31] |
| 0.0071 | Solovay \& Strassen [32] |
| 0.0062 | Paul, Prauß, \& Reischuk [22] |
| 0.0050 | Paul, Pippenger, Szemerédi, \& Trotter [21] |
| 0.0043 | Ruzzo, Simon, \& Tompa [28] |
| $9.30 \times 10^{-4}$ | Santoro, Sidney J., Sidney S., \& Urrutia [29] |
| $2.23 \times 10^{-4}$ | Kahn, Klawe, \& Kleitman [13] |
| $2.00 \times 10^{-4}$ | Lenstra A., Lenstra H., \& Lovász [19] |
| $6.57 \times 10^{-5}$ | Schwartz, Sharir \& Siegel [30] |
| $2.24 \times 10^{-5}$ | Leighton \& Leiserson [18] |
| $1.64 \times 10^{-5}$ | Karmarkar, Karp, Lipton, Lovász, \& Luby [14] |
| $1.40 \times 10^{-6}$ | Brassard \& Bratley [7] |
| $3.58 \times 10^{-7}$ | Yao A. \& Yao F. [39] |
| $3.39 \times 10^{-7} \dagger$ | Goldreich, Goldwasser, \& Micali [11] |
| $8.56 \times 10^{-13} \ddagger$ | Lenstra A., Lenstra H., \& Lovász [19] |
| $4.17 \times 10^{-15 §}$ | Vazirani U. \& Vazirani V. [35] |
| $1.53 \times 10^{-15}$ | Plumstead B. \& Plumstead J. [27] |

*A popular conjecture at this time was that the limit of $i(X)$ was $\left(\frac{5 \ln \pi}{27}\right)^{e} \approx$ 0.01475 .
${ }^{\dagger}$ i(Goldwasser)
$\ddagger$ ¿(Lenstra H.)
${ }^{\S}$ A long-standing conjecture held that this was the quantum of obliviousness, based on the remarkable coincidence that it is Planck's constant when measured in electron volts - seconds. This conjecture was finally disproved by Joan Boyar [personal communication].

Table 4: The Top Coefficients of Obliviousness

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1. Erdös, Graham, & Szemerédi [9]
2. Graham, Lawler, Lenstra, & Rinnooy Kan [12]
3. Lawler, Tarjan, & Valdez [17]
4. Tarjan & Vishkin [33]
5. Vishkin & Wigderson [36]
```

Table 5: Monotone Erdös Numbers
Definition: The monotone Erdös number of $X$ is the length of a longest directed path in $G$ between Paul Erdös and $X .{ }^{\dagger}$

Table 5 illustrates that Wigderson's monotone Erdös number is 5. We know of no one in the Theoretical Computer Science community with a greater finite value.

## 5 Conclusions

Figures often beguile me, particularly when I have the arranging of them myself; in which case the remark attributed to Disraeli would often apply with justice and force: "There are three kinds of lies: lies, damned lies, and statistics."

- Mark Twain, 1924

This paper introduced the following new statistics:

1. Spotlight factor
2. Coefficient of obliviousness
3. Monotone Erdös number
4. Johnson number

Definition: The Johnson number $\mathcal{J}(X)$ is the number of statistics that $X$ has studied before David Johnson.

[^2]Theorem: $\mathcal{J}($ Tompa $)=4$.

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[^0]:    *A preliminary version of these results was presented at the $29^{\text {th }}$ Annual Symposium

[^1]:    *In this definition, the notation ". Vishkin", for example, is of course the radix 27 fraction, where $a=1, b=2, \ldots, z=26$, and blanks and punctuation represent 0 .

[^2]:    ${ }^{\dagger}$ The ordinary Erdös number is the length of a shortest path in the undirected version of $G$. In contrast, the challenge in this new definition is to find some $X$ with as great a finite value as possible.

