

# Computers Are Not That Great!

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# It Matters How Fast Programs Run

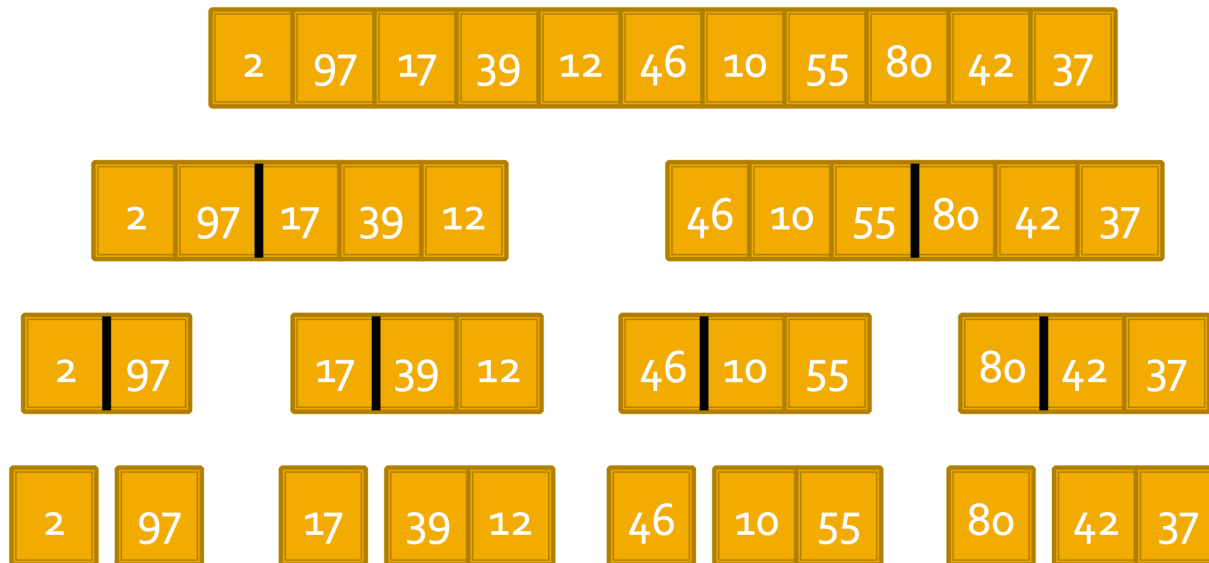
- Computers are amazingly fast ... but that's because we usually ask them to do really easy stuff; they can do billions of instructions per second (gips?) ...
- So, what's a "really easy" computation? *cn*
  - Looking up in a dictionary or address book how the letters you've typed might be completed
  - Recovering a losslessly compressed file
  - Looking in a file for a specific letter string
  - ...

# A Little More Work To Do

- How long would it take the Census Bureau to alphabetize the US population by first name?
  - Recall Exchange Sort & Bubble Sort – both are “ $n^2$  algorithms,” meaning they take  $cn^2$  seconds for some amount of time  $c$  – usually the overhead to process one item; lets estimate  $c \sim 0.5 \mu s$
  - If the US population is  $n = 310,000,000 = 3.1 \times 10^8$
  - $n^2 = 3.1 \times 10^8 \times 3.1 \times 10^8 = 9.6 \times 10^{16}$
  - $0.5 \mu s = .000\ 000\ 5s = 5.0 \times 10^{-7}s$
  - $cn^2 = 5.0 \times 10^{-7}s \times 9.6 \times 10^{16} = 48 \times 10^9s = 1521 \text{ years}$

# One Other Approach

- Recall there was also the Merge Sort



Merge Sort requires  $cn \log_2 n$

$$cn \log_2 n = 5.0 \times 10^{-7} s \times 3.1 \times 10^8 \times 29$$

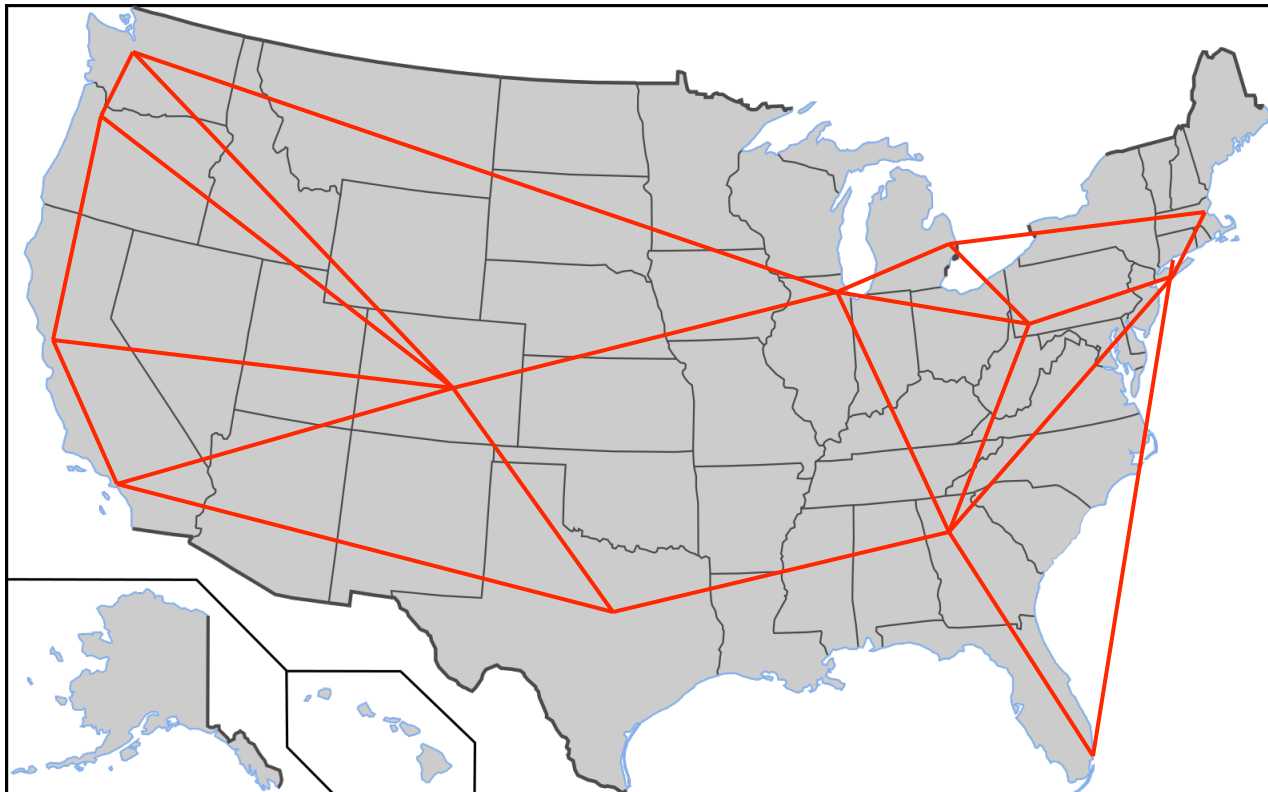
$$= 15.5 \times 10 s \times 29 = 4495 s = 1.24 \text{ hr}$$

# Summarizing Alphabetize Task

- The input size  $n = 310,000,000 = 3.1 \times 10^8$
- Exchange sort & Bubble sort require  $cn^2$  time
- For  $c=1/2$  microsecond,  $cn^2 = 1521$  years
- Merge sort requires  $cn \log_2 n$  time
- Because  $\log_2 n = 29$ ,  $cn \log_2 n = 1.24$  hours
  
- Algorithms matter ... and smarter algorithms are better

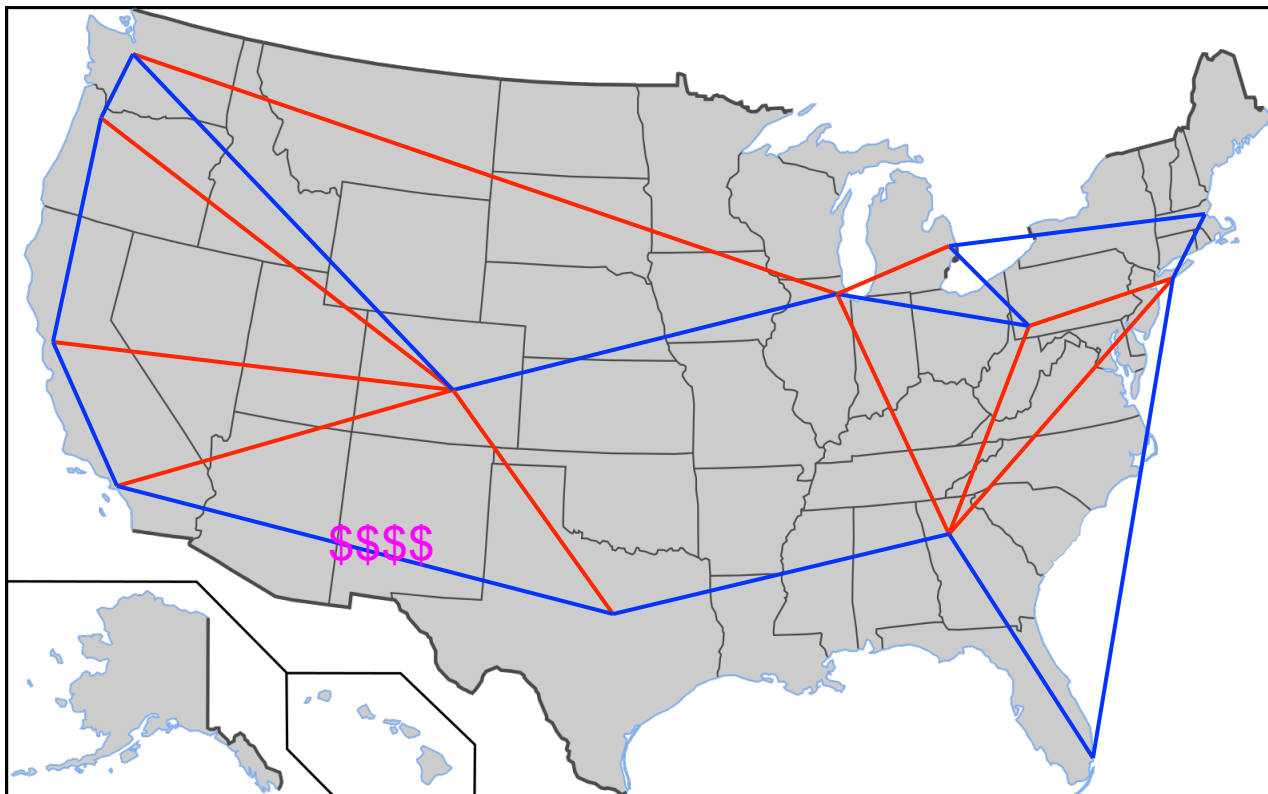
# Computations That Are Harder Still

- More data means more work, but sometimes it means a lot more work
- Traveling salesman problem:



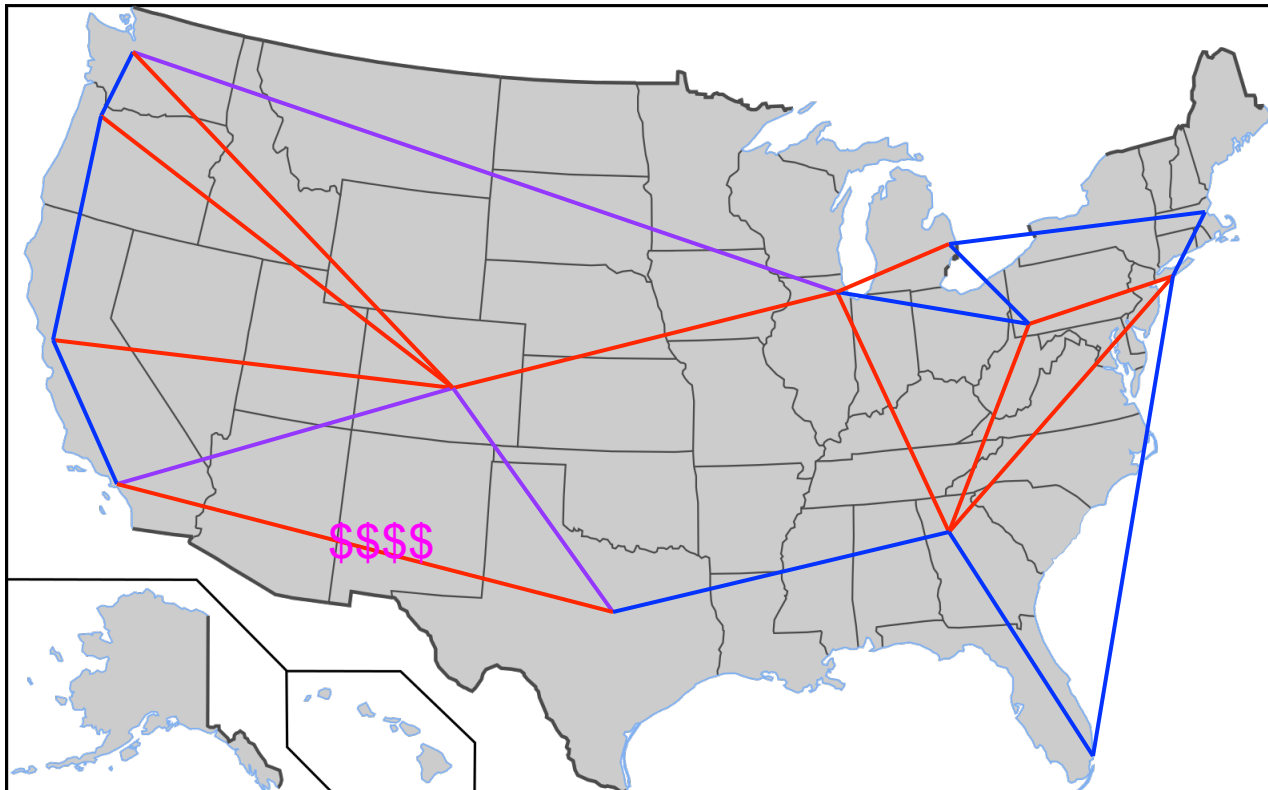
# TSP – Visit Each City Once

- Minimize the cost of the plane tickets
  - Finding a tour is reasonably easy
  - Finding the cheapest tour is NP-hard



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# NP-Hard & NP-Complete Problems

- NP, which stands for “nondeterministic polynomial time” (don’t learn that), is a class of problems with these features:
  - They are easy (like  $cn^2$ , perhaps) ways to solve if the computer can guess and is always right
  - They have no known easy (like  $cn^5$ , say) solutions, it seems, if the computer can’t guess, which it can’t
  - All known solutions effectively check all possible alternatives and pick the best
  - These are “normal” computations, like TSP
  - “Complete” means solve one and you’ve solved all

# In Computer Science Programs ...

- ... Are Data
- For Example: Processing is a program that accepts YOUR program as data and runs it ... so it “computes on” (processes) your program
- Except for really trivial languages (e.g. HTML) all programming languages are universal – CS people can write a program in that one language, say Processing, which can run programs in any other language – all programs
- This is the “Universality Principle”

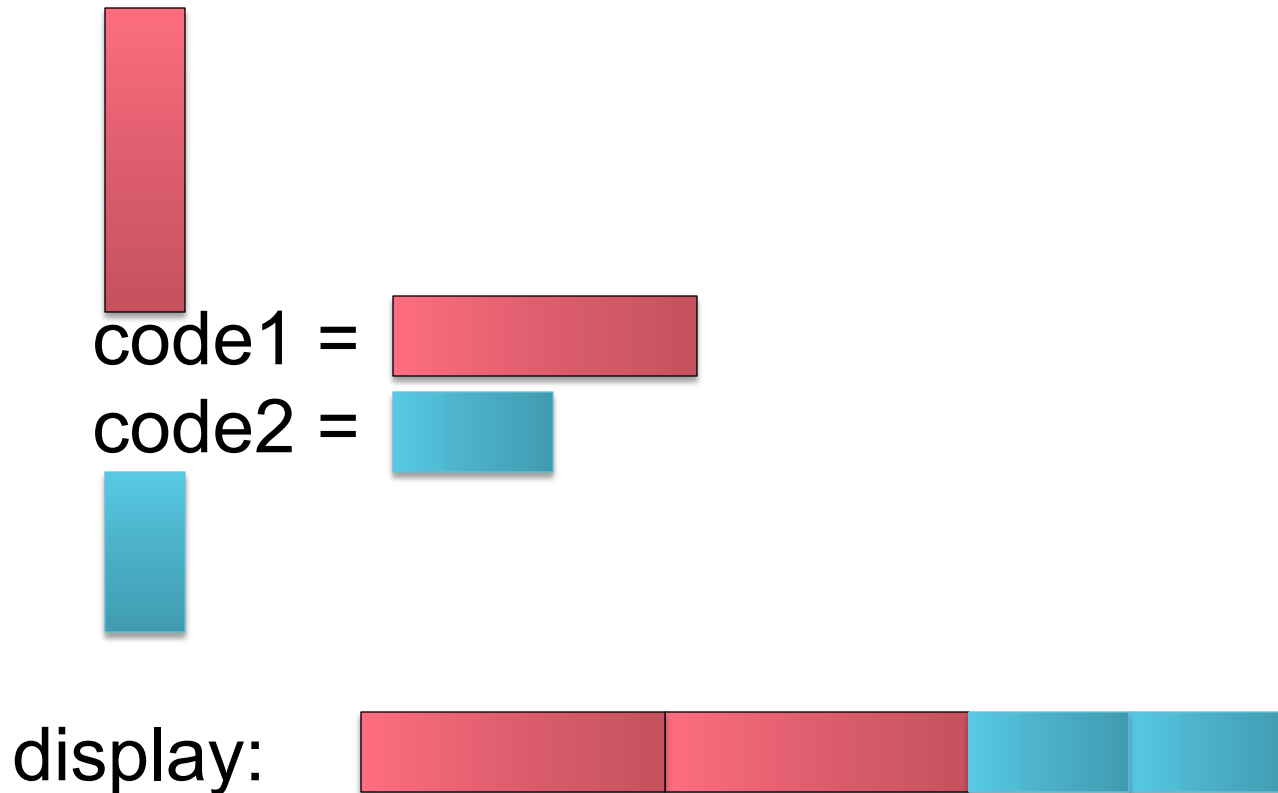
# A Program To Print Itself Out

```
String code1, code2;
void setup( ) {
  size(500, 400);
  background(255);
  noLoop( );
  fill(0);
}
void draw( ) {
  code1 = "\\String code1, code2; " +
    "void setup( ) { " +
    "size(500,400); " +
    "background(255); " +
    "noLoop( ); fill(0); } " +
    "void draw( ) { " +
    "code1 = ;\\"";
  code2 = "code1 = code1 + code1 + \\code2 = \" + code2 + code2; \" +
    "text(code1, 50, 50, 200, 400); }\\\" ";
  code1 = code1 + code1 + "code2 = \" + code2 + code2;
  text(code1, 50, 50, 200, 400);
}
```

Fixing the tiny syntactic differences is easy

# Schematic of Self-Printing Pgm

- Divide the program into two halves --



# Running the Program ...

## ■ Output

```
"String code1, code2; void setup(
) { size(500,400);
background(255); noLoop( );
fill(0); } void draw( ) { code1 =
;""String code1, code2; void
setup( ) { size(500,400);
background(255); noLoop( );
fill(0); } void draw( ) { code1 =
;"code2 = code1 = code1 +
code1 + "code2 = " + code2 +
code2; text(code1, 50, 50, 200,
400);}" code1 = code1 + code1
+ "code2 = " + code2 + code2;
text(code1, 50, 50, 200, 400); }"
```

## Helpfully Formatted Output

```
"String code1, code2; void setup(
) { size(500,400);
background(255); noLoop( );
fill(0); } void draw( ) { code1 = ;"
```

```
"String code1, code2; void setup(
) { size(500,400);
background(255); noLoop( );
fill(0); } void draw( ) { code1 = ;"
```

```
code2 = " code1 = code1 +
code1 + "code2 = " + code2 +
code2; text(code1, 50, 50, 200,
400); }"
```

```
code1 = code1 + code1 + "code2
= " + code2 + code2;
text(code1, 50, 50, 200, 400); }"
```

# Adding Additional Code

- Notice that new code can be added, and the program can still print itself out

Put the new code here and here

```
void draw( ) {  
  code1 = "\"String code1, code2; \" +  
    "void setup( ) { \" +  
    "size(500,400); \" +  
    "background(255); \" +  
    "noLoop( ); fill(0); } \" +  
    "void draw( ) { \" +  
    "code1 = ;\";  
  code2 = "code1 = code1 + code1 + \"code2 = \" + code2 + code2; \" +  
    "text(code1, 50, 50, 200, 400); }\"  ";  
  code1 = code1 + code1 + "code2 = \" + code2 + code2;  
  text(code1, 50, 50, 200, 400);  
}
```

# Summarizing

- A self-printing program shows that programs can manipulate program text ...
- Examples of programs manipulating programs
  - The highlighter that “colors” your programs
  - The translator that converts Processing code into machine code so a computer can run it
  - The code that figures out what you did wrong when you forget a semicolon
  - A debugger can help you find errors in your pgm

# A Problem That Can't Be Solved

- Suppose we want to determine if a Processing program draws a red circle or not
- It seems possible, perhaps ...
  - Analyze the code to see if it displays any circles
  - Check if any of the circles it draws are red
  - Etc.
- Suppose **Boolean** check-pde(**String** code) is a Processing function that determines if a Processing program draws a red circle (**return true**) or does not draw a red circle (**return false**)



# Assuming check-pde() works ...

```
String code = "void trick( )... ";
void setup( ) {
    size(200,200); background(255); noLoop( );
}
void draw( ) {
    trick( ); //Guaranteed to get it wrong!
}
void trick( ) {
    if (check-pde(code)){ //does code draw red circle?
        fill(0,0,255); //check-pde says yes
    } else {
        fill(255,0,0); //check-pde says no
    }
    ellipse(100,100,10,10);
}
```

Analyze What Happens

# The Impact

- There are simple problems that computers cannot solve, b/c probs are not algorithmic ... no deterministic sequence of operations can find the answer; debugging is an example
- Alan Turing's insight in 1936



# Summary

- We considered how “hard” computations can be, where “hard” is measured as running time
- Linear time – thinking about how long the code runs
- Quadratic;  $N \log N$  – thinking about sorting
- NP Hard and the TSP
- Universal machine – yeah Turing!
- Undecidability