

## CSE 326 Selected Practice Problem Solutions

**2.6** As discussed in class, the answer to the first part is  $2^{2^{N-1}}$  and the answer to the second part is  $O(\log \log D)$ .

### 2.10

(a)  $O(N)$

(b)  $O(N^2)$

(c) The answer depends on how many digits past the decimal point are computed. Each digit costs  $O(N)$ .

### 3.22 Pseudocode:

```
Create stack
Read in first token
while (token is not "=")
    if (token is a number)
        push the token onto the stack
    else
        if (token is "+")
            pop a
            pop b
            push a + b
        if (token is "-")
            pop a
            pop b
            push a - b
        if (token is "*")
            pop a
            pop b
            push a * b
        if (token is "/")
            pop a
            pop b
            push a / b
read next token
```

4.1

(a)  $A$ .

(b)  $G, H, I, L, M,$  and  $K$ .

4.8

(a)  $- * * a b + c d e$ .

(b)  $(( a * b ) * ( c + d )) - e$ .

(c)  $a b * c d + * e -$ .

4.27 See Figures 1-4.



Figure 1: 4.27 After accessing 3

4.28 See Figure 5.



Figure 2: 4.27 After accessing 9

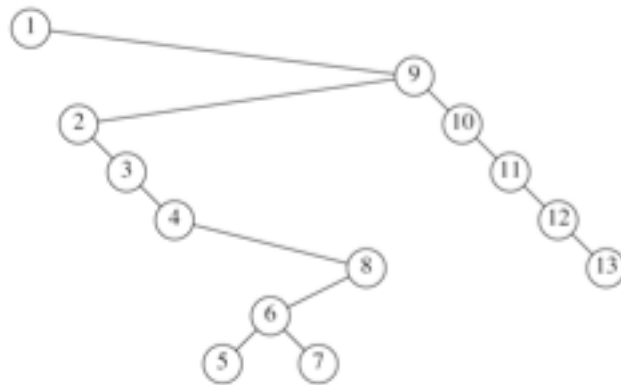


Figure 3: 4.27 After accessing 1

**6.2** See Figure 6.

**6.3** The result of three *deleteMins*, starting with both of the heaps in Exercise 6.2, is in Figure 7.

**6.30** Clearly the claim is true for  $k = 1$ . Suppose it is true for all values  $i = 1, 2, \dots, k$ . A  $B_{k+1}$  tree is formed by attaching a  $B_k$  tree to the root of a  $B_k$  tree. Thus by induction, it contains a  $B_0$  through  $B_{k-1}$  tree, as well as the newly attached  $B_k$  tree, proving the claim.

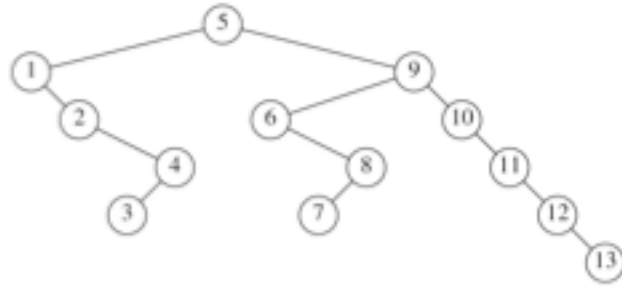


Figure 4: 4.27 After accessing 5



Figure 5: 4.28

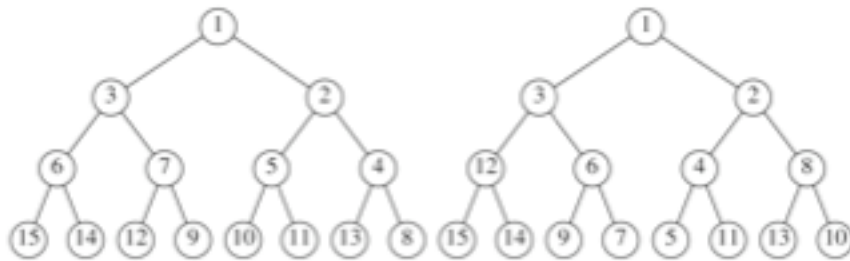


Figure 6: 6.2

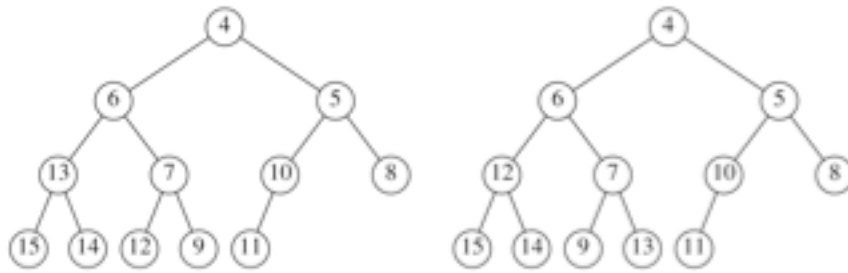


Figure 7: 6.3