

Intro to Digital Design

Combinational Logic

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Teaching Assistants:

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Wen Li

Introducing Your Course Staff

- ❖ Your Instructor: just call me Justin
 - CSE Associate Teaching Professor
 - From California (UC Berkeley and the Bay Area)
 - Raising a toddler takes up energy and dictates my schedule

- ❖ TAs:



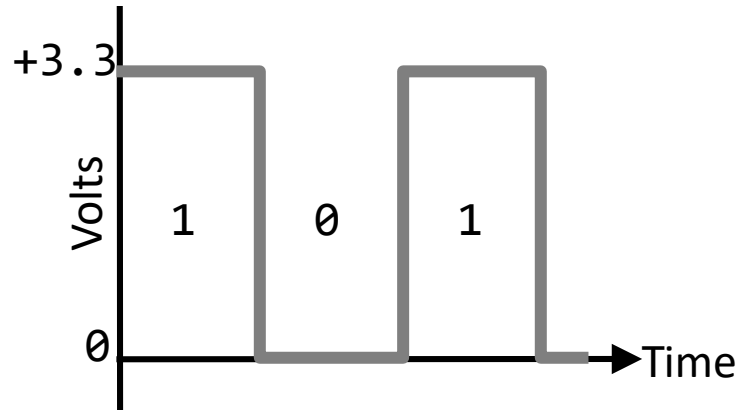
- Available in labs, office hours, and on Ed discussion
 - An invaluable source of information and help
- ❖ Get to know us – we are here to help you succeed!



Course Motivation

- ❖ Electronics an increasing part of our lives
 - Computers & phones
 - Vehicles (cars, planes)
 - Robots
 - Portable & household electronics
- ❖ An *introduction* to digital logic design
 - **Lecture:** How to think about hardware, basic higher-level circuit design techniques – preparation for EE/CSE469
 - **Lab:** Hands-on FPGA programming using Verilog – preparation for EE/CSE371

Digital vs. Analog



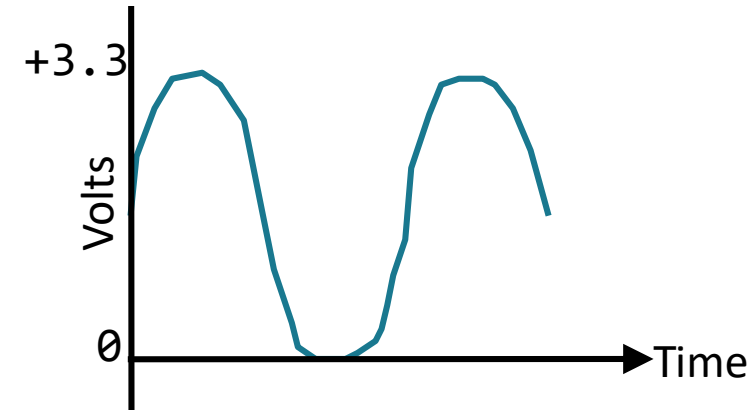
Digital:

Discrete set of possible values

Binary (2 values):

On, 3.3 V, high, TRUE, "1"

Off, 0 V, low, FALSE, "0"



Analog:

Values vary over a continuous range

Digital vs. Analog Systems

- ❖ Digital systems are more reliable and less error-prone
 - Slight errors can cascade in Analog system
 - Digital systems reject a significant amount of error; easy to cascade
- ❖ Computers use digital circuits internally
 - CPU, memory, I/O
- ❖ Interface circuits with “real world” often analog
 - Sensors & actuators

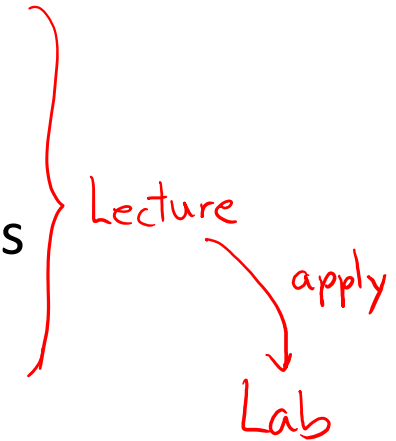
***This course is about logic design,
not system design (processor architecture),
and not circuit design (transistor level)***

Digital Design: What's It All About?

- ❖ Come up with an implementation using a set of primitives given a functional description and constraints
- ❖ Digital design is in some ways more art than a science
 - The creative spirit is in combining primitive elements and other components in new ways to achieve a desired function
- ❖ However, unlike art, we have objective measures of a design (*i.e.*, constraints):
 - Performance
 - Power
 - Cost

Digital Design: What's It All About?

- ❖ How do we learn how to do this?
 - Learn about the primitives and how to use them
 - Learn about design representations
 - Learn formal methods and tools to manipulate representations
 - Look at design examples
 - Use trial and error – CAD tools and prototyping (practice!)



Lecture Outline

- ❖ **Course Logistics**
- ❖ Combinational Logic Review
- ❖ Combinational Logic in the Lab

Bookmarks

- ❖ Website: <https://courses.cs.washington.edu/courses/cse369/24sp/>
 - Schedule (lecture slides, lab specs), weekly calendar, other useful documents
- ❖ Ed Discussion: <https://edstem.org/us/courses/56771/>
 - Announcements made here
 - Ask and answer questions – staff will monitor and contribute
- ❖ Gradescope: <https://www.gradescope.com/courses/746339/>
 - Lab submissions, Quiz grades, regrade requests
- ❖ Canvas: <https://canvas.uw.edu/courses/1718545/>
 - Grade book, Zoom links, lecture recordings

Grading

- ❖ Labs (66%)
 - 6 regular labs – 1 week each
 - Labs 3-4: 60 points each, Labs 1&2, 5-7: 100 points each
 - 1 “final project” – 2 weeks
 - Lab 8 Check-In: 10 points, Lab 8: 150 points
- ❖ 3 Quizzes (no final exam)
 - Quiz 1 (10%): 20 min in class on April 23
 - Quiz 2 (10%): 30 min in class on May 14
 - Quiz 3 (14%): 60 min in class on May 28
- ❖ This class uses a straight scale ($\geq 95\% \rightarrow 4.0$)
 - Extra credit points count the same as regular points

Labs

- ❖ Lab Hours: Wed & Thu 2:30-5:20 pm (CSE 003)
- ❖ Each student will get a lab kit for the quarter
 - Lab kit picked up from CSE 003 during labs/OHs this week
 - Install software on laptop (Windows or VM)
- ❖ Labs are combination of report + demo
 - Submit via Gradescope **Wednesdays before 2:30 pm**
 - 10-minute demos done in lab sections (sign-up process)
- ❖ Late penalties:
 - No lab report can be submitted more than two days late
 - 5 late day tokens to prevent penalties, 10%/day after that
 - No penalties on lab demos, but must be done by EOD Friday

Collaboration Policy

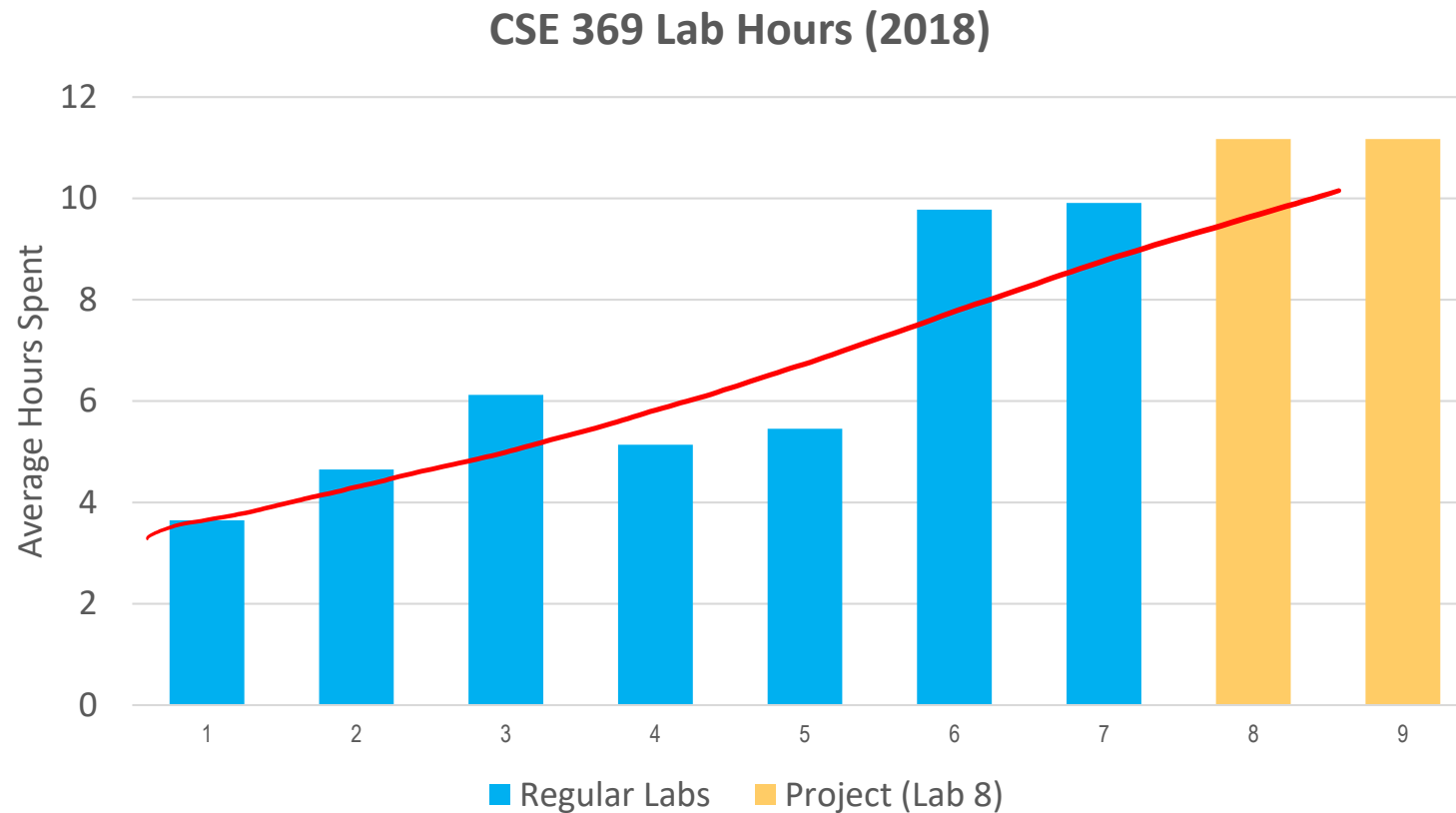
- ❖ Labs and project are to be completed *individually*
 - Goal is to give every student the hands-on experience
 - Violation of these rules is grounds for failing the class

- ❖ **OK:**
 - Discussing lectures and/or readings, studying together
 - *High-level* discussion of general approaches
 - Help with debugging, tools peculiarities, etc.

- ❖ **Not OK:**
 - Developing a lab together
 - Giving away solutions or having someone else do your lab for you

Course Workload

- ❖ The workload (3 credits) ramps up significantly towards the end of the quarter:

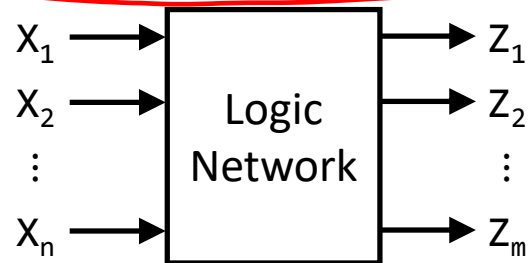


Lecture Outline

- ❖ Course Logistics
- ❖ **Combinational Logic Review**
- ❖ Combinational Logic in the Lab

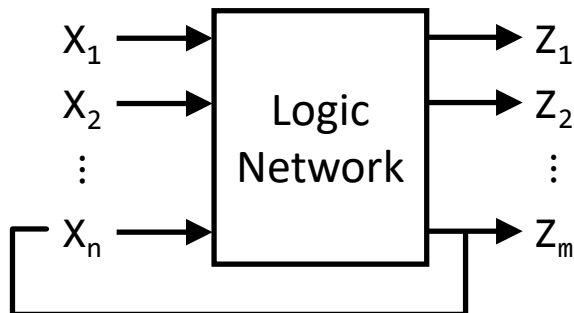
Combinational vs. Sequential Logic

❖ Combinational Logic (CL)



- Network of logic gates without feedback
- Outputs are functions only of inputs

❖ Sequential Logic (SL)



- The presence of feedback introduces the notion of "state"
- Circuits that can "remember" or store information

Representations of Combinational Logic

- 1 ❖ Text Description
 - 2 ❖ Circuit Description
 - ~~Transistors~~ Not covered in 369
 - Logic Gates
 - 3 ❖ Truth Table
 - 4 ❖ Boolean Expression
- ❖ All are equivalent!

Example: Simple Car Electronics

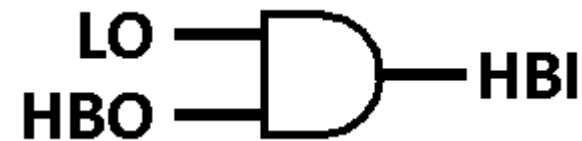
- ❖ Door Ajar (DriverDoorOpen, PassengerDoorOpen)

- $DA = DDO + PDO$



- ❖ High Beam Indicator (LightsOn, HighBeamOn)

- $HBI = LO \cdot HBO$



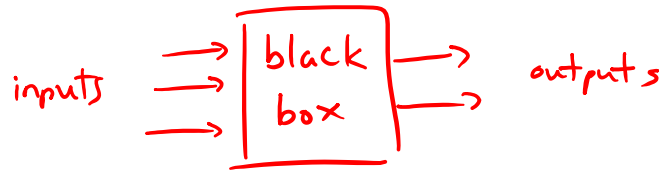
- ❖ Seat Belt Light (DriverBeltIn, PassengerBeltIn, Passenger)

- $SBL = \overline{DBI} + (P \cdot \overline{PBI})$



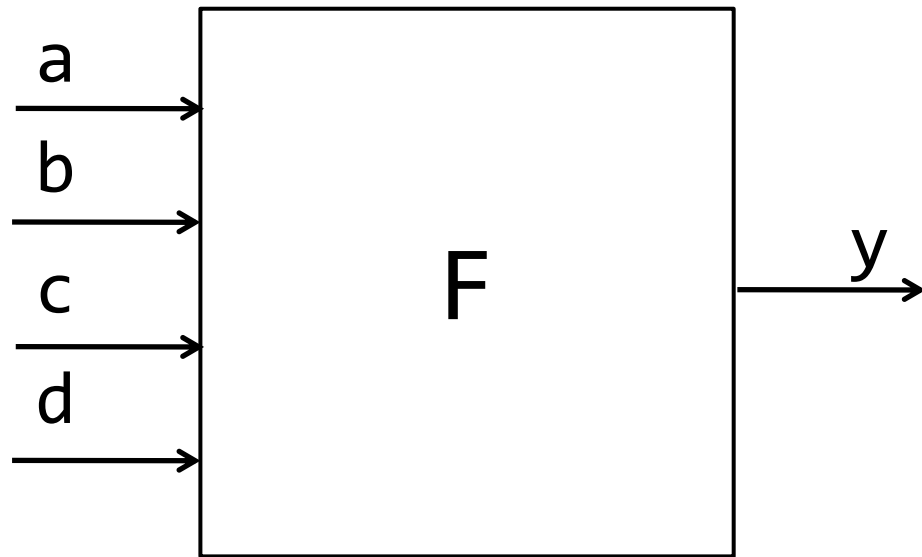
Truth Tables

- ❖ Table that relates the inputs to a combinational logic (CL) circuit to its output
 - Output *only* depends on current inputs
 - Use abstraction of 0/1 instead of high/low voltage
 - Shows output for every possible combination of inputs (“black box” approach)



- ❖ How big is the table?
 - 0 or 1 for each of N inputs 2^N rows
 - Each output is a separate function of inputs, so don't need to add rows for additional outputs

CL General Form



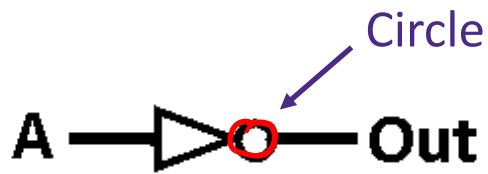
If we have N inputs, how many distinct functions F do we have?

2^N output "positions", each being 0/1
 so $2^{(2^N)}$ possible functions

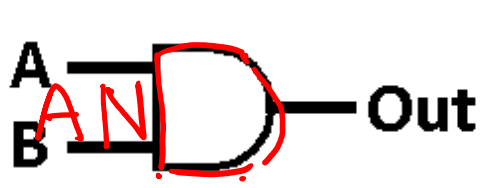
a	b	c	d	y
0	0	0	0	F(0,0,0,0)
0	0	0	1	F(0,0,0,1)
0	0	1	0	F(0,0,1,0)
0	0	1	1	F(0,0,1,1)
0	1	0	0	F(0,1,0,0)
0	1	0	1	F(0,1,0,1)
0	1	1	0	F(0,1,1,0)
0	1	1	1	F(0,1,1,1)
1	0	0	0	F(1,0,0,0)
1	0	0	1	F(1,0,0,1)
1	0	1	0	F(1,0,1,0)
1	0	1	1	F(1,0,1,1)
1	1	0	0	F(1,1,0,0)
1	1	0	1	F(1,1,0,1)
1	1	1	0	F(1,1,1,0)
1	1	1	1	F(1,1,1,1)

Handwritten notes: A red arrow points from the 'y' column header to the text '0/1'. A red bracket on the left side of the table spans all 16 rows and is labeled 2^N rows (16).

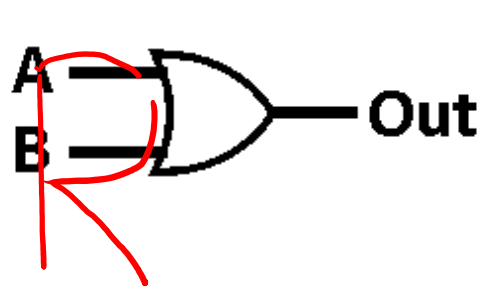
Logic Gate Names and Symbols

❖ **NOT** 

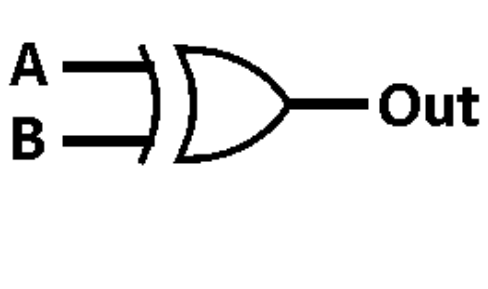
A	Out
0	1
1	0

❖ **AND** 

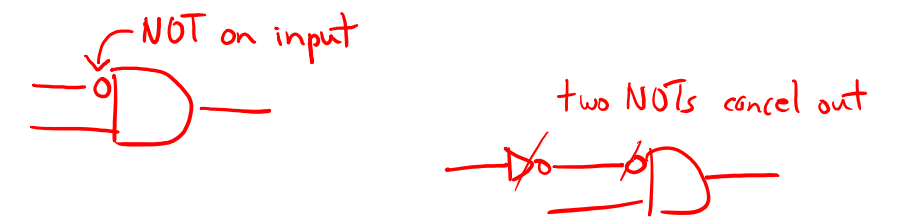
A	B	Out
0	0	0
0	1	0
1	0	0
1	1	1

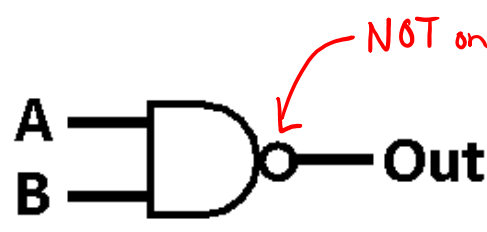
❖ **OR** 

A	B	Out
0	0	0
0	1	1
1	0	1
1	1	1

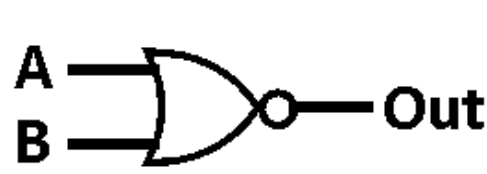
❖ **XOR** 

A	B	Out
0	0	0
0	1	1
1	0	1
1	1	0




❖ **NAND** 

A	B	Out
0	0	1
0	1	1
1	0	1
1	1	0

❖ **NOR** 

A	B	Out
0	0	1
0	1	0
1	0	0
1	1	0

❖ **XNOR** 

A	B	Out
0	0	1
0	1	0
1	0	0
1	1	1

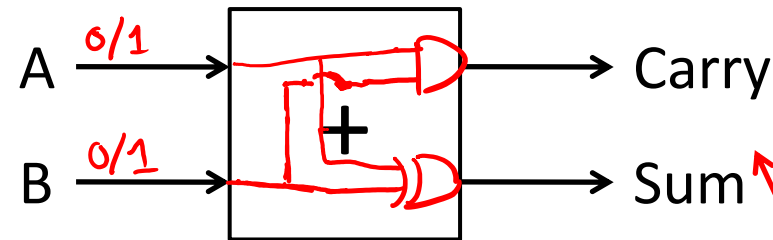
More Complicated Truth Tables

3-Input Majority

How many rows? $2^3 = 8$ rows

A	B	C	Out
0	0	0	0
0	0	1	0
0	1	0	0
0	<u>1</u>	<u>1</u>	→ 1
1	0	0	0
<u>1</u>	0	<u>1</u>	→ 1
<u>1</u>	<u>1</u>	0	→ 1
<u>1</u>	<u>1</u>	<u>1</u>	→ 1

1-bit Adder



A	B	Carry	Sum
0	0	0	0
0	1	0	1
1	0	0	1
1	1	1	0

2 separate functions (columns)

$A \cdot B$ $A \oplus B$

Boolean Algebra

- ❖ Represent inputs and outputs as variables
 - Each variable can only take on the value 0 or 1
- ⌋ ❖ Overbar is NOT: “logical complement”
 - If A is 0, then \bar{A} is 1 and vice-versa
- ∨ ❖ Plus (+) is 2-input OR: “logical sum”
- ∧ ❖ Product (\cdot) is 2-input AND: “logical product”
- ❖ All other gates and logical expressions can be built from combinations of these
 - *e.g.*, $A \text{ XOR } B = A \oplus B = \bar{A}B + \bar{B}A$

Truth Table to Boolean Expression

- ❖ Read off of table
 - For 1, write variable name
 - For 0, write complement of variable

a	b	c	row
0	0	0	1
0	1	1	2
1	0	1	3
1	1	0	4

❖ Sum of Products (SoP)

- Take rows with 1's in output column, sum products of inputs
- $C = \bar{A}B + \bar{B}A$

sets to 1 when input combination matches

We can show that these are equivalent!

❖ Product of Sums (PoS)

- Take rows with 0's in output column, product the sum of the complements of the inputs
- $C = (A + B) \cdot (\bar{A} + \bar{B})$

sets to 0 when input combination matched

Basic Boolean Identities

$$\diamond X + 0 = X$$

$$\diamond X + 1 = 1$$

$$\diamond X + X = X$$

$$\diamond X + \bar{X} = 1$$

$$\diamond \bar{\bar{X}} = X$$

$$\diamond X \cdot 1 = X$$

$$\diamond X \cdot 0 = 0$$

$$\diamond X \cdot X = X$$

$$\diamond X \cdot \bar{X} = 0$$

Basic Boolean Algebra Laws

❖ Commutative Law:

$$X + Y = Y + X$$

$$X \cdot Y = Y \cdot X$$

❖ Associative Law:

$$X + (Y + Z) = (X + Y) + Z$$

$$X \cdot (Y \cdot Z) = (X \cdot Y) \cdot Z$$

❖ Distributive Law:

$$X \cdot (Y + Z) = X \cdot Y + X \cdot Z$$

$$X + YZ = (X + Y) \cdot (X + Z)$$

Advanced Laws (Absorption)

$$\diamond X + XY = X$$

$$\diamond XY + X\bar{Y} = X$$

$$\diamond X + \bar{X}Y = X + Y$$

$$\diamond X(X + Y) = X$$

$$\diamond (X + Y)(X + \bar{Y}) = X$$

$$\diamond X(\bar{X} + Y) = XY$$

$$\begin{aligned} X + \bar{X}Y &= X \cdot 1 + \bar{X}Y = X \cdot (1 + Y) + \bar{X}Y \\ &= X + XY + \bar{X}Y \\ &= X + Y \end{aligned}$$
$$Y = \bar{X}Y + XY = Y(\bar{X} + X)$$

Practice Problem

❖ Boolean Function: $F = \bar{X}YZ + XZ$

Truth Table:

X	Y	Z	F
0	0	0	0
0	0	1	0
0	1	0	0
→ 0	1	1	1
1	0	0	0
⇒ 1	0	1	1
1	1	0	0
⇒ 1	1	1	1

Simplification:

$$= \bar{X}YZ + X\bar{Y}Z + XYZ$$

$$= \bar{X}YZ + XZ$$

$$= (\bar{X}Y + X)Z$$

$$= (X + Y)Z$$

$$= XZ + YZ$$

2 gates (1 OR, 1 AND)

Which of these is "simpler"?

3 gates (2 AND, 1 OR)

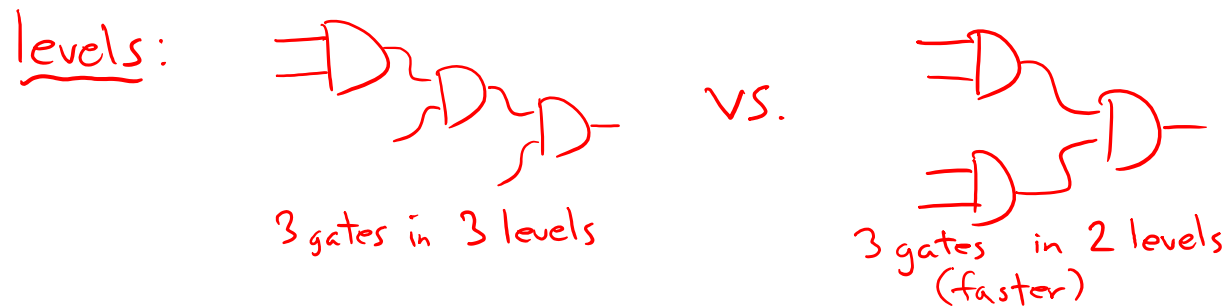
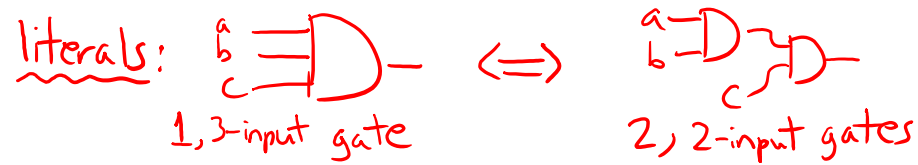
Technology Break

Lecture Outline

- ❖ Course Logistics
- ❖ Combinational Logic Review
-  **Combinational Logic in the Lab**

Why Is This Useful?

- ❖ Logic minimization: reduce complexity at gate level
 - Allows us to build smaller and faster hardware
 - Care about both # of gates, # of literals (gate inputs), # of gate levels, and types of logic gates



types: speed of NOT gate vs. AND gate?

Why Is This Useful?

- ❖ Logic minimization: reduce complexity at gate level
 - Allows us to build smaller and faster hardware
 - Care about both # of gates, # of literals (gate inputs), # of gate levels, and types of logic gates
- ❖ Faster hardware?
 - Fewer inputs implies faster gates in some technologies
 - Fan-ins (# of gate inputs) are limited in some technologies
 - Fewer levels of gates implies reduced signal propagation **delays**
 - # of gates (or gate packages) influences manufacturing costs
- ★ Simpler Boolean expressions → smaller transistor networks → smaller circuit delays
→ faster hardware



Are Logic Gates Created Equal?

❖ No!

2-Input Gate Type	# of CMOS transistors
NOT	2
AND	6
OR	6
NAND	4
NOR	4
XOR	8
XNOR	8

← simplest, but not too useful

} useful, and simpler than alternatives

❖ Can recreate all other gates using only NAND or only NOR gates

- Called “universal” gates
- e.g., $A \text{ NAND } A = \bar{A}$, $B \text{ NOR } B = \bar{B}$
- DeMorgan’s Law helps us here!



x	y	NAND	
→ 0	0	1	0 → 1
0	1	1	
1	0	1	
→ 1	1	0	1 → 0

DeMorgan's Law

X	Y	\bar{X}	\bar{Y}	NOR $\overline{X+Y}$	$\bar{X} \cdot \bar{Y}$	NAND $\overline{X \cdot Y}$	$\bar{X} + \bar{Y}$
0	0	1	1	1	1	1	1
0	1	1	0	0	0	1	1
1	0	0	1	0	0	1	1
1	1	0	0	0	0	0	0

❖ $\overline{X + Y} = \bar{X} \cdot \bar{Y}$

❖ $\overline{X \cdot Y} = \bar{X} + \bar{Y}$

❖ In Boolean Algebra, converts between AND-OR and OR-AND expressions

▪ $Z = \bar{A}\bar{B}C + \bar{A}B\bar{C} + A\bar{B}\bar{C}$

▪ $\bar{Z} = (A + B + \bar{C}) \cdot (A + \bar{B} + \bar{C}) \cdot (\bar{A} + B + \bar{C})$

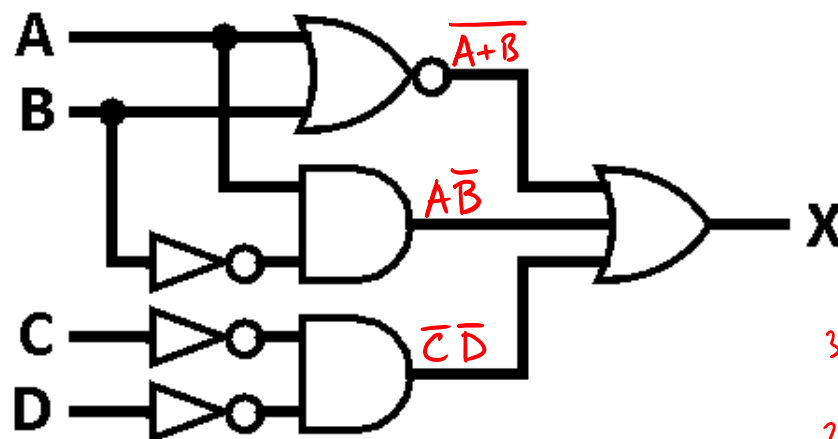
❖ At gate level, can convert from AND/OR to NAND/NOR gates

▪ "Flip" all input/output bubbles and "switch" gate



DeMorgan's Law Practice Problem

❖ Simplify the following diagram:



$$X = \overline{A+B} + A\bar{B} + \bar{C}D$$

DeMorgan's

$$X = \bar{A}\bar{B} + A\bar{B} + \bar{C}D$$

5 gates

$$X = \bar{B} + \bar{C}D$$

3-4 gates

$$X = \bar{B} + \overline{C+D}$$

DeMorgan's

$$X = \overline{B(C+D)}$$

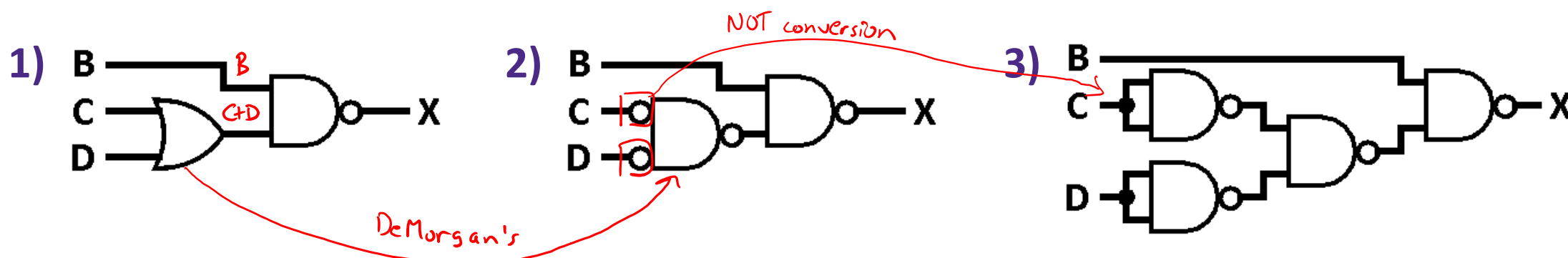
2-3 gates

let E=C+D, so

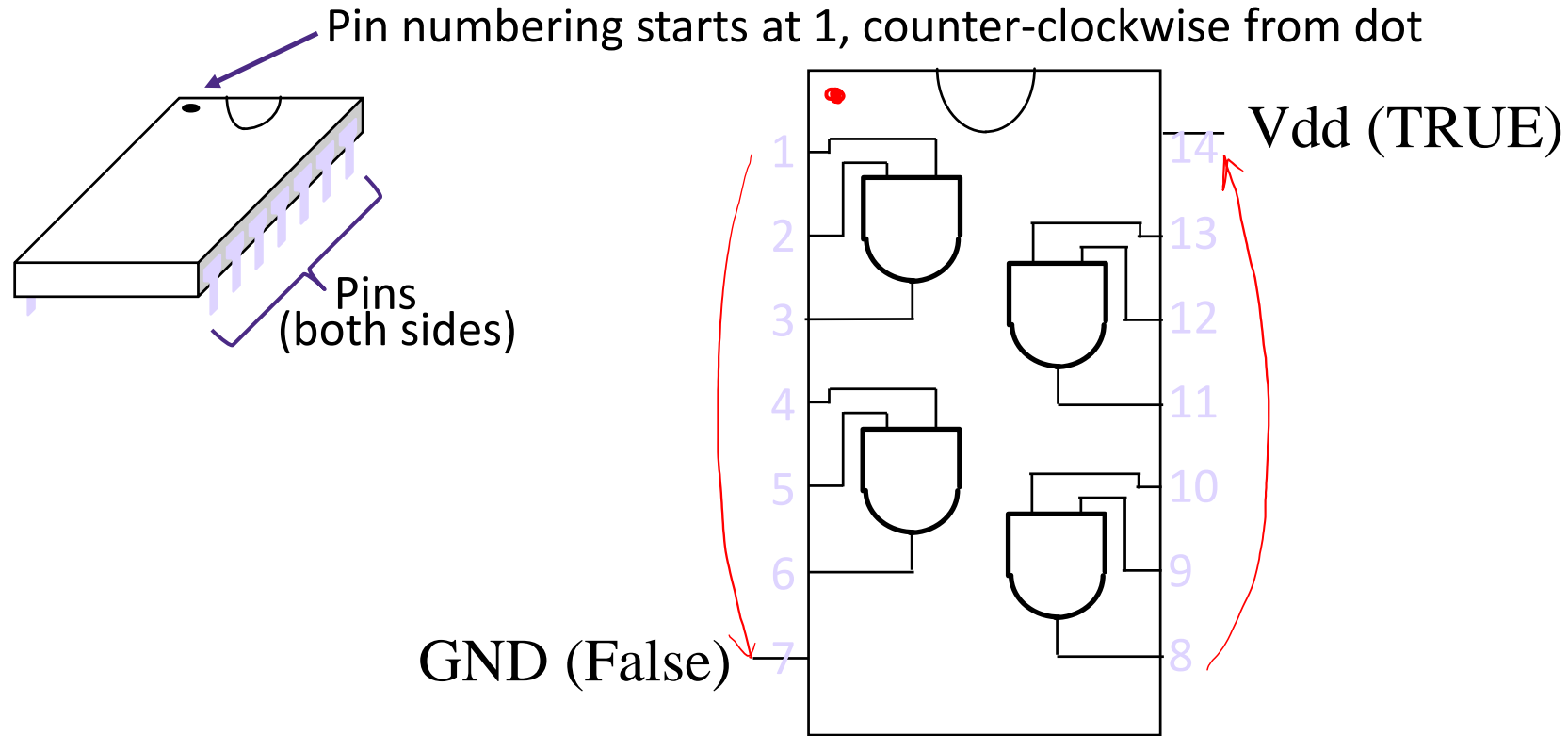
$$X = \bar{B} + \bar{E}$$

$$X = \overline{BE}$$

❖ Then implement with only NAND gates:



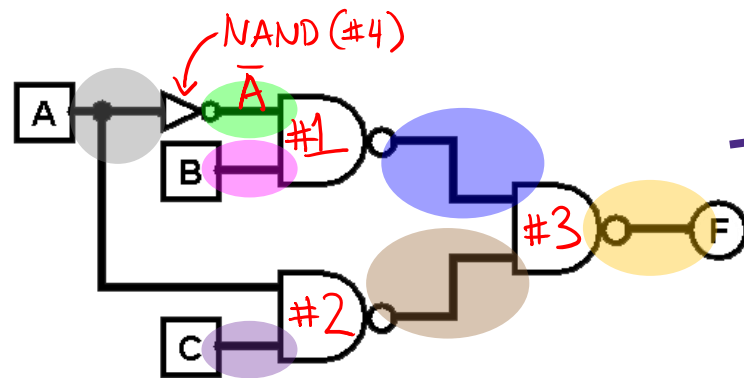
Transistor-Transistor Logic (TTL) Packages



- ❖ Diagrams like these and other useful/helpful information can be found on part **data sheets**
 - It's really useful to learn how to read these

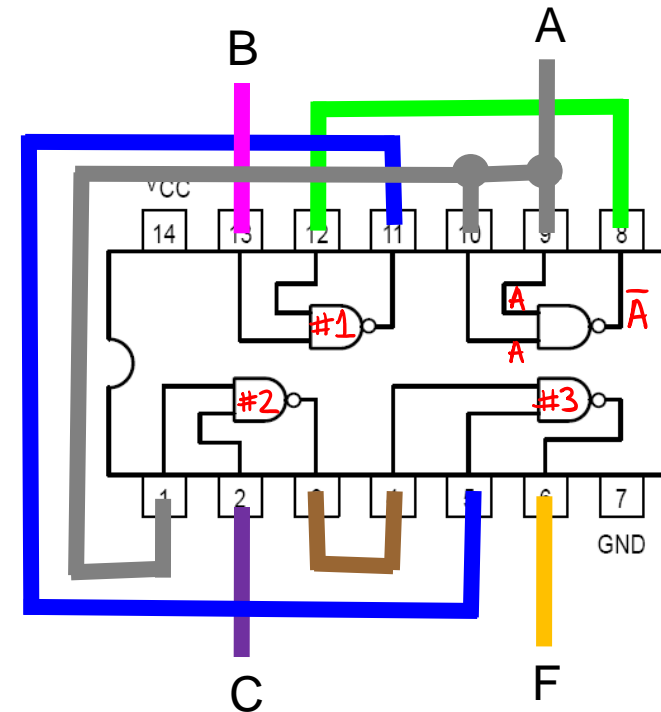
Mapping truth tables to logic gates

- ❖ Given a truth table:
 - 1) Write the Boolean expression
 - 2) Minimize the Boolean expression
 - 3) Draw as gates
 - 4) Map to available gates
 - 5) Determine # of packages and their connections

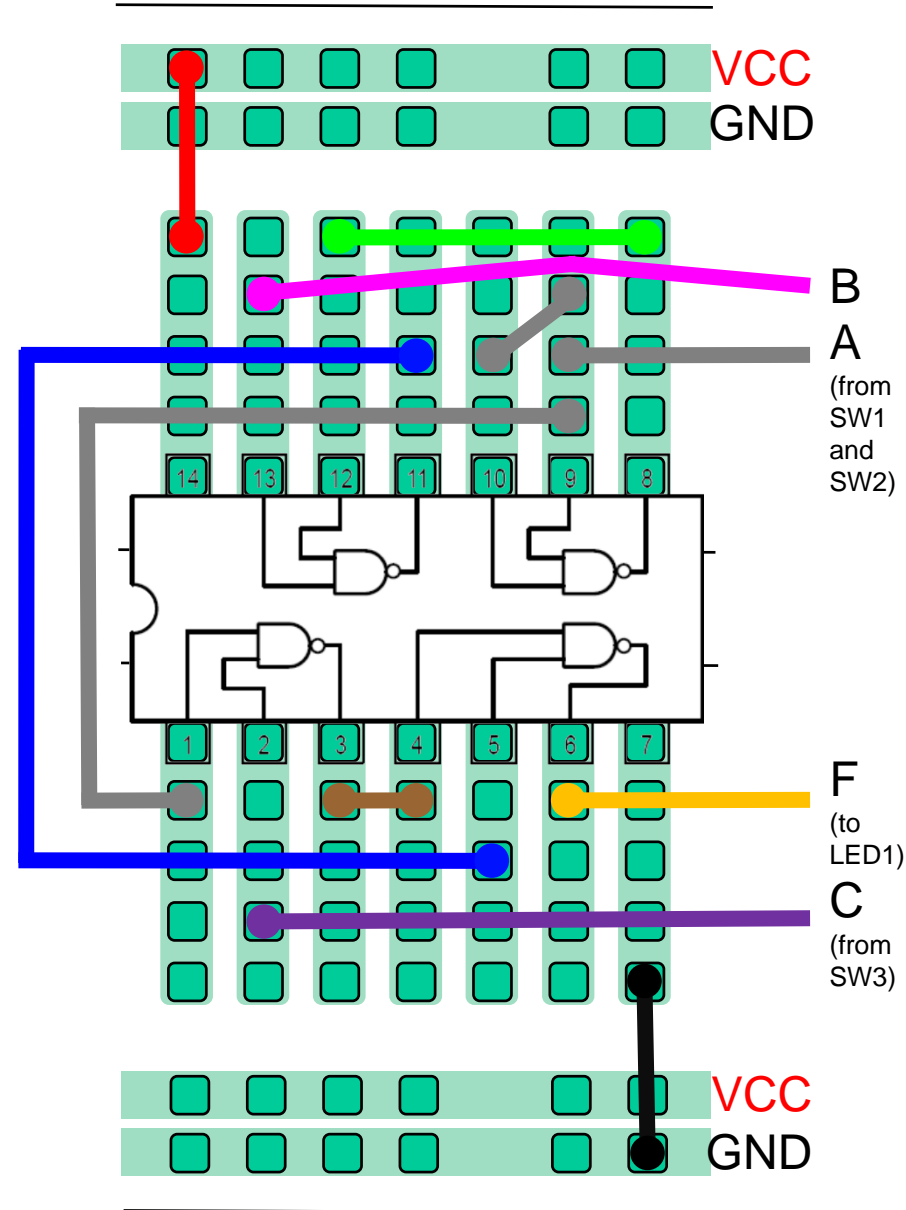
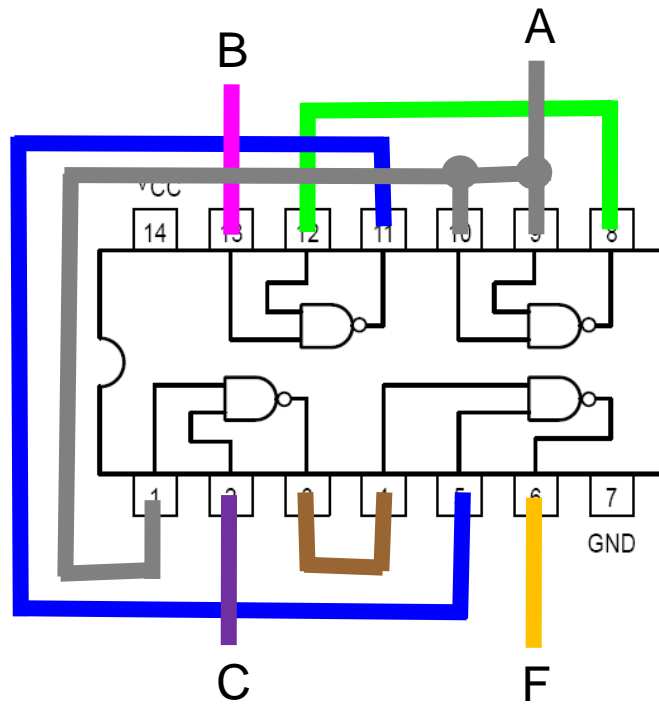


7 nets (wires) in this design

(4) →



Breadboarding circuits



Summary

- ❖ Digital systems are constructed from Combinational and Sequential Logic
- ❖ Logic minimization to create smaller and faster hardware
- ❖ Gates come in TTL packages that require careful wiring

