

# Intro to Digital Design

## Combinational Logic

**Instructor:** Justin Hsia

**Teaching Assistants:**

Caitlyn Rawlings

Donovan Clay

Emilio Alcantara

Joy Jung

Naoto Uemura

# Introducing Your Course Staff



- ❖ Your Instructor: just call me Justin
  - CSE Associate Teaching Professor
  - From California (UC Berkeley and the Bay Area)
  - Raising a toddler takes up energy and dictates my schedule

- ❖ TAs:



Caitlyn



Donovan



Emilio



Joy



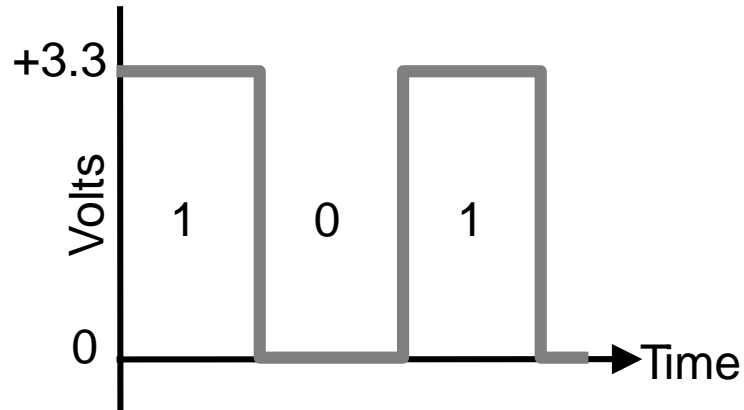
Naoto

- Available in labs, support (office) hours, and on Ed
  - An invaluable source of information and help
- ❖ Get to know us – we are here to help you succeed!

# Course Motivation

- ❖ Electronics an increasing part of our lives
  - Computers & phones
  - Vehicles (cars, planes)
  - Robots
  - Portable & household electronics
  
- ❖ *An introduction* to digital logic design
  - **Lecture:** How to think about hardware, basic higher-level circuit design techniques – preparation for EE/CSE469
  - **Lab:** Hands-on FPGA programming using Verilog – preparation for EE/CSE371

# Digital vs. Analog



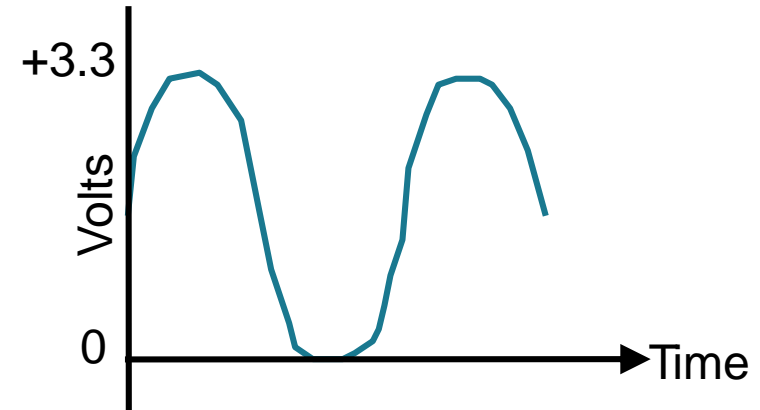
## Digital:

Discrete set of possible values

## Binary (2 values):

On, 3.3 V, high, TRUE, "1"

Off, 0 V, low, FALSE, "0"



## Analog:

Values vary over a continuous range

# Digital vs. Analog Systems

- ❖ Digital systems are more reliable and less error-prone
  - Slight errors can cascade in Analog system
  - Digital systems reject a significant amount of error; easy to cascade
- ❖ Computers use digital circuits internally
  - CPU, memory, I/O
- ❖ Interface circuits with “real world” often analog
  - Sensors & actuators

***This course is about logic design,  
not system design (processor architecture),  
and not circuit design (transistor level)***

# Digital Design: What's It All About?

- ❖ Come up with an implementation using a set of primitives given a functional description and constraints
- ❖ Digital design is in some ways more art than a science
  - The creative spirit is in combining primitive elements and other components in new ways to achieve a desired function
- ❖ However, unlike art, we have objective measures of a design (*i.e.*, constraints):
  - Performance
  - Power
  - Cost

# Digital Design: What's It All About?

- ❖ How do we learn how to do this?
    - Learn about the primitives and how to use them
    - Learn about design representations
    - Learn formal methods and tools to manipulate representations
    - Look at design examples
    - Use trial and error – CAD tools and prototyping (practice!)
- 
- Lecture
- apply
- Lab

# Lecture Outline

- ❖ **Course Logistics**
- ❖ Combinational Logic Review
- ❖ Combinational Logic in the Lab



# Bookmarks

- ❖ Website: <https://cs.uw.edu/369/>
  - Schedule (lecture slides, lab specs), weekly calendar, other useful documents
- ❖ Ed Discussion: <https://edstem.org/us/courses/50615/>
  - Announcements made here
  - Ask and answer questions – staff will monitor and contribute
- ❖ Gradescope: <https://www.gradescope.com/courses/680677/>
  - Lab submissions, Quiz grades, regrade requests
- ❖ Canvas: <https://canvas.uw.edu/courses/1695952/>
  - Grade book, Zoom links, lecture recordings

# Grading

- ❖ Labs (66%)
  - 6 regular labs – 1 week each
    - Labs 3-4: 60 points each, Labs 1&2, 5-7: 100 points each
  - 1 “final project” – 2 weeks
    - Lab 8 Check-In: 10 points, Lab 8: 150 points
- ❖ 3 Quizzes (no final exam)
  - Quiz 1 (10%): 20 min in class on January 30
  - Quiz 2 (10%): 30 min in class on February 20
  - Quiz 3 (14%): 60 min in class on March 5
- ❖ This class uses a straight scale ( > 95% → 4.0 )
  - Extra credit points count the same as regular points

# Labs

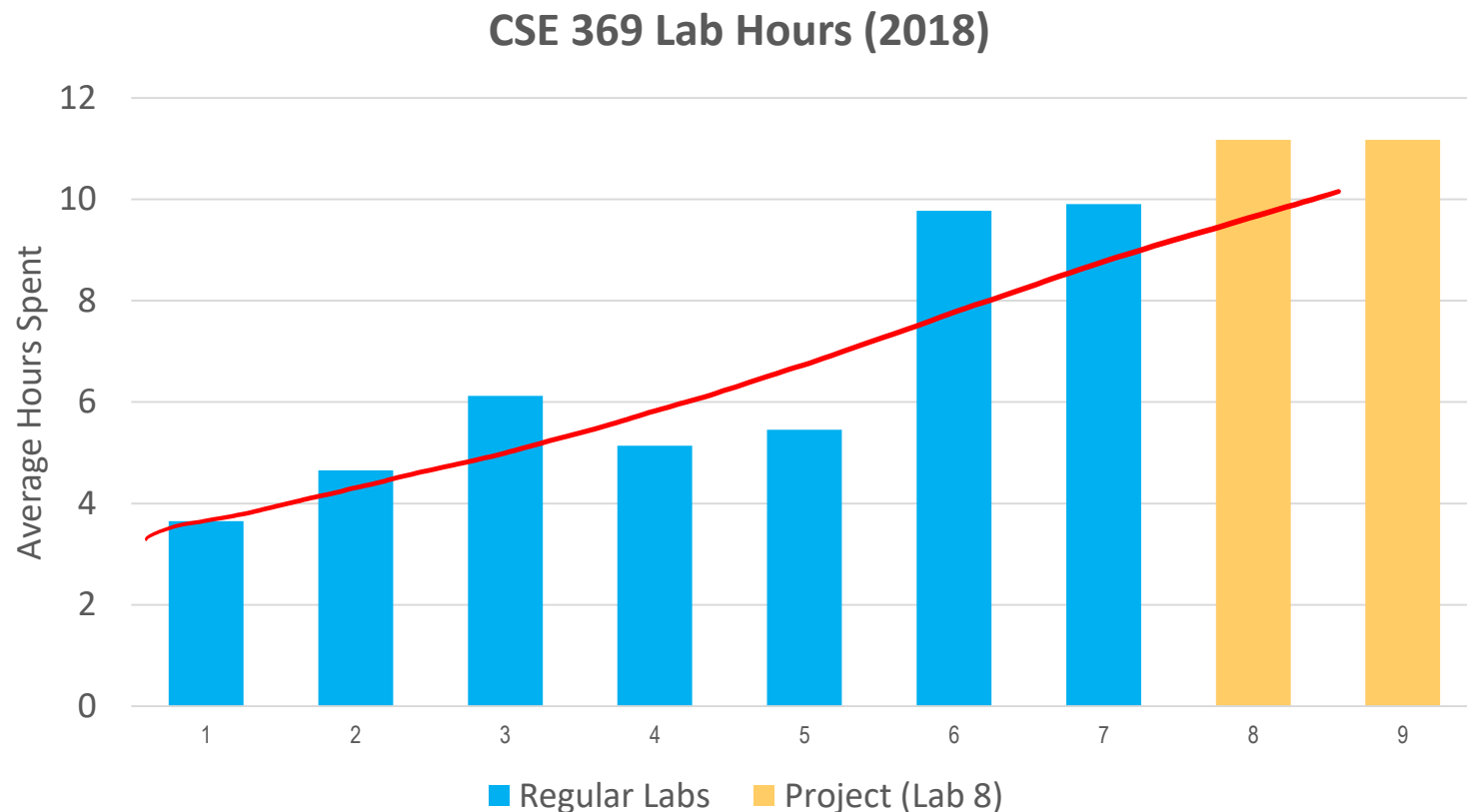
- ❖ Lab Hours: Wed & Thu 2:30-5:20 pm (CSE 003)
- ❖ Each student will get a lab kit for the quarter
  - Lab kit picked up from CSE 003 during labs/OHs this week
  - Install software on laptop (Windows or VM)
- ❖ Labs are combination of report + demo
  - Submit via Gradescope **Wednesdays before 2:30 pm**
  - 10-minute demos done in lab sections (sign-up process)
- ❖ Late penalties:
  - No lab report can be submitted more than two days late
  - 4 late day tokens to prevent penalties, 10%/day after that
  - No penalties on lab demos, but must be done by EOD Friday

# Collaboration Policy

- ❖ Labs and project are to be completed *individually*
  - Goal is to give every student the hands-on experience
  - Violation of these rules is grounds for failing the class
  
- ❖ **OK:**
  - Discussing lectures and/or readings, studying together
  - *High-level* discussion of general approaches
  - Help with debugging, tools peculiarities, etc.
  
- ❖ **Not OK:**
  - Developing a lab together
  - Giving away solutions or having someone else do your lab for you

# Course Workload

- ❖ The workload (3 credits) ramps up significantly towards the end of the quarter:

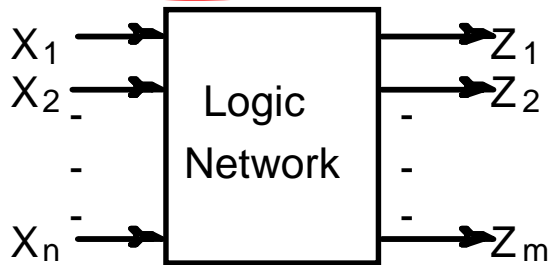


# Lecture Outline

- ❖ Course Logistics
- ❖ **Combinational Logic Review**
- ❖ Combinational Logic in the Lab

# Combinational vs. Sequential Logic

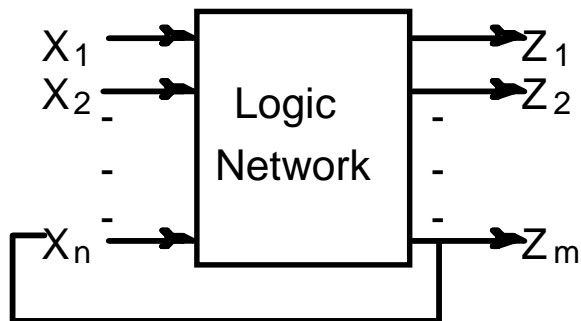
## ❖ **Combinational Logic (CL)**



Network of logic gates without feedback.

Outputs are functions only of inputs.

## ❖ **Sequential Logic (SL)**



The presence of feedback introduces the notion of “state.”  
Circuits that can “remember” or store information.

# Representations of Combinational Logic

- 1 ❖ Text Description
  - 2 ❖ Circuit Description
    - ~~Transistors~~ Not covered in 369
    - Logic Gates
  - 3 ❖ Truth Table
  - 4 ❖ Boolean Expression
- ❖ All are equivalent!



# Example: Simple Car Electronics

- ❖ Door Ajar (DriverDoorOpen, PassengerDoorOpen)

- $DA = DDO + PDO$



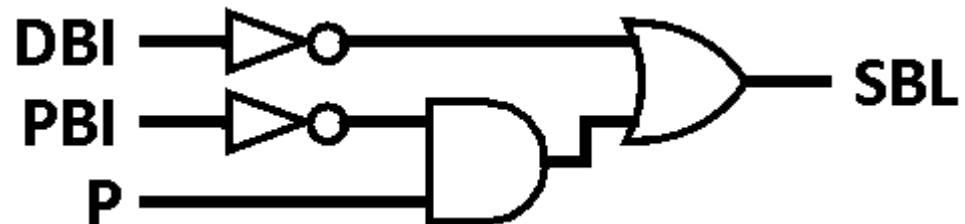
- ❖ High Beam Indicator (LightsOn, HighBeamOn)

- $HBI = LO \cdot HBO$



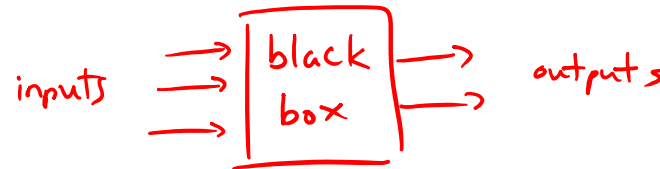
- ❖ Seat Belt Light (DriverBeltIn, PassengerBeltIn, Passenger)

- $SBL = \overline{DBI} + (P \cdot \overline{PBI})$



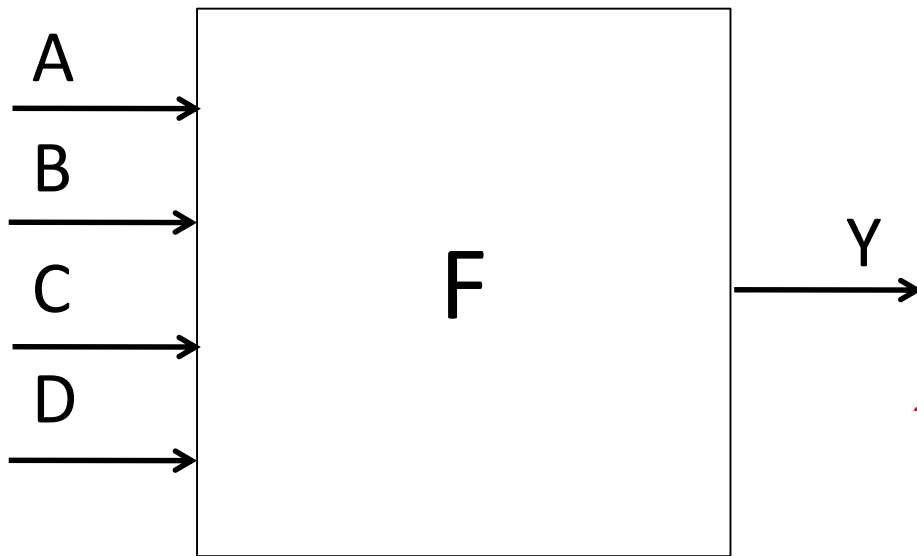
# Truth Tables

- ❖ Table that relates the inputs to a combinational logic (CL) circuit to its output
  - Output *only* depends on current inputs
  - Use abstraction of 0/1 instead of high/low voltage
  - Shows output for every possible combination of inputs (“black box” approach)



- ❖ How big is the table?
  - 0 or 1 for each of  $N$  inputs, so  $2^N$  rows
  - Each output is a separate function of inputs, so don't need to add rows for additional outputs

# CL General Form



If  $N$  inputs, how many distinct functions  $F$  do we have?  $2^N$  output "positions", each being 0/1

Function maps each row to 0 or 1, so  $2^{(2^N)}$  possible functions

a	b	c	d	y
0	0	0	0	$F(0,0,0,0)$
0	0	0	1	$F(0,0,0,1)$
0	0	1	0	$F(0,0,1,0)$
0	0	1	1	$F(0,0,1,1)$
0	1	0	0	$F(0,1,0,0)$
0	1	0	1	$F(0,1,0,1)$
0	1	1	0	$F(0,1,1,0)$
0	1	1	1	$F(0,1,1,1)$
1	0	0	0	$F(1,0,0,0)$
1	0	0	1	$F(1,0,0,1)$
1	0	1	0	$F(1,0,1,0)$
1	0	1	1	$F(1,0,1,1)$
1	1	0	0	$F(1,1,0,0)$
1	1	0	1	$F(1,1,0,1)$
1	1	1	0	$F(1,1,1,0)$
1	1	1	1	$F(1,1,1,1)$

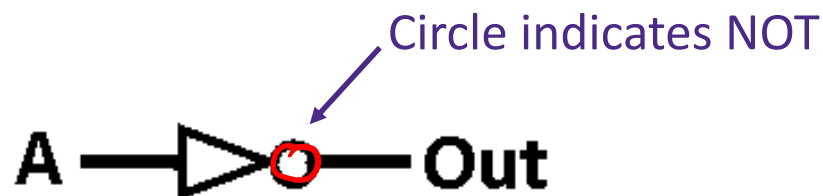
$2^N$  rows  
(16)

0/1

# Logic Gates (1/2)

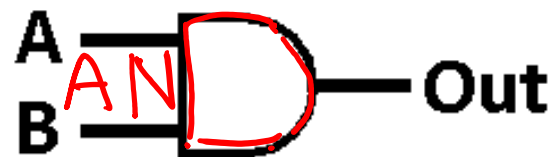
- ❖ Special names and symbols:

**NOT**



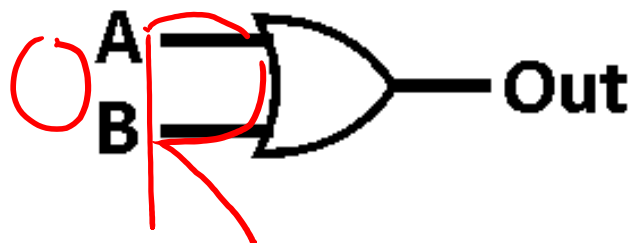
A	Out
0	1
1	0

**AND**



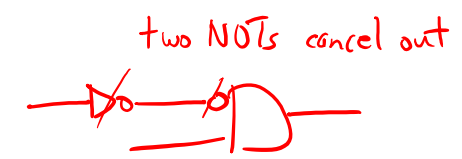
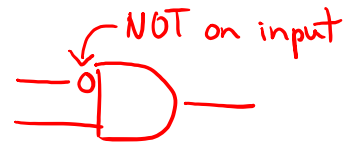
A	B	Out
0	0	0
0	1	0
1	0	0
1	1	1

**OR**



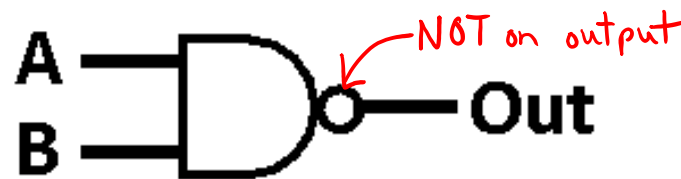
A	B	Out
0	0	0
0	1	1
1	0	1
1	1	1

# Logic Gates (2/2)



❖ Special names and symbols:

**NAND**



A	B	Out
0	0	1
0	1	1
1	0	1
1	1	0

**NOR**



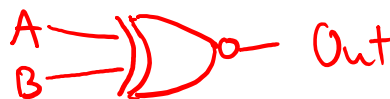
A	B	Out
0	0	1
0	1	0
1	0	0
1	1	0

**XOR**



A	B	Out
0	0	0
0	1	1
1	0	1
1	1	0

**XNOR**



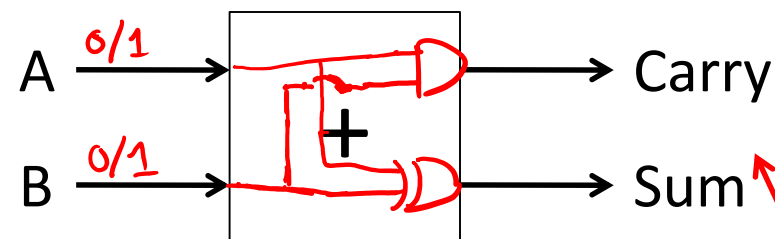
# More Complicated Truth Tables

## 3-Input Majority

How many rows?  $2^3 = 8$  rows

A	B	C	Out
0	0	0	0
0	0	1	0
0	1	0	0
0	<u>1</u>	<u>1</u>	→ 1
1	0	0	0
<u>1</u>	0	<u>1</u>	→ 1
<u>1</u>	<u>1</u>	0	→ 1
<u>1</u>	<u>1</u>	<u>1</u>	→ 1

## 1-bit Adder



A	B	Carry	Sum
0	0	0	0
0	1	0	1
1	0	0	1
1	1	1	0

2 separate functions (columns)

$A \cdot B$     $A \oplus B$

# Boolean Algebra

- ❖ Represent inputs and outputs as variables
  - Each variable can only take on the value 0 or 1
- ¬ ❖ Overbar is NOT: “logical complement”
  - If A is 0, then  $\bar{A}$  is 1 and vice-versa
- ∨ ❖ Plus (+) is 2-input OR: “logical sum”
- ∧ ❖ Product (·) is 2-input AND: “logical product”
- ❖ All other gates and logical expressions can be built from combinations of these
  - *e.g.*,  $A \text{ XOR } B = A \oplus B = \bar{A}B + \bar{B}A$

# Truth Table to Boolean Expression

❖ Read off of table

- For 1, write variable name
- For 0, write complement of variable

a	b	c	row
0	0	0	1
0	1	1	2
1	0	1	3
1	1	0	4

❖ *Sum of Products (SoP)*

- Take rows with 1's in output column, sum products of inputs

*sets to 1 when input combination matches*

■  $C = \overline{A}B + \overline{B}A$

*row 2 row 3*

We can show that these are equivalent!

❖ *Product of Sums (PoS)*

- Take rows with 0's in output column, product the sum of the complements of the inputs

*sets to 0 when input combination matched*

■  $C = (A + B) \cdot (\overline{A} + \overline{B})$

*row 1 row 4*



# Basic Boolean Identities

$$\diamond X + 0 = X$$

$$\diamond X \cdot 1 = X$$

$$\diamond X + 1 = 1$$

$$\diamond X \cdot 0 = 0$$

$$\diamond X + X = X$$

$$\diamond X \cdot X = X$$

$$\diamond X + \bar{X} = 1$$

$$\diamond X \cdot \bar{X} = 0$$

$$\diamond \bar{\bar{X}} = X$$

# Basic Boolean Algebra Laws

## ❖ Commutative Law:

$$X + Y = Y + X$$

$$X \cdot Y = Y \cdot X$$

## ❖ Associative Law:

$$X + (Y + Z) = (X + Y) + Z$$

$$X \cdot (Y \cdot Z) = (X \cdot Y) \cdot Z$$

## ❖ Distributive Law:

$$X \cdot (Y + Z) = X \cdot Y + X \cdot Z$$

$$X + YZ = (X + Y) \cdot (X + Z)$$

# Advanced Laws (Absorption)

$$\diamond X + XY = X$$

$$\diamond XY + X\bar{Y} = X$$

$$\diamond \underline{X + \bar{X}Y} = X + Y$$

$$Y = \bar{X}Y + XY = Y(\bar{X} + X)$$

$$\begin{aligned} X + \bar{X}Y &= X \cdot 1 + \bar{X}Y = X \cdot (1 + Y) + \bar{X}Y \\ &= X + XY + \bar{X}Y \\ &= X + Y \quad \checkmark \end{aligned}$$

$$\diamond X(X + Y) = X$$

$$\diamond (X + Y)(X + \bar{Y}) = X$$

$$\diamond X(\bar{X} + Y) = XY$$

# Practice Problem

❖ Boolean Function:  $F = \underline{\bar{X}YZ} + \underline{XZ}$

Truth Table:

X	Y	Z	F
0	0	0	0
0	0	1	0
0	1	0	0
→ 0	1	1	1
1	0	0	0
⇒ 1	0	1	1
1	1	0	0
⇒ 1	1	1	1

Simplification:

$$= \bar{X}YZ + X\bar{Y}Z + XYZ$$

$$= \bar{X}YZ + XZ$$

*distribution*

$$= (\bar{X}Y + X)Z$$

*absorption*

$$= (X + Y)Z$$

*distribute*

$$= XZ + YZ$$

2 gates (1 OR, 1 AND)

Which of these is "simpler"?

3 gates (2 AND, 1 OR)

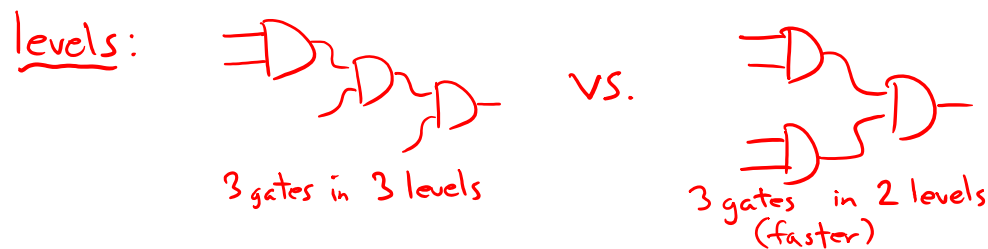
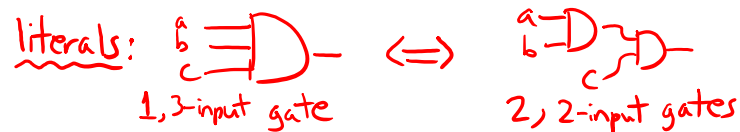
# Technology Break

# Lecture Outline

- ❖ Course Logistics
- ❖ Combinational Logic Review
-  **Combinational Logic in the Lab**

# Why Is This Useful?

- ❖ Logic minimization: reduce complexity at gate level
  - Allows us to build smaller and faster hardware
  - Care about both # of gates, # of literals (gate inputs), # of gate levels, and types of logic gates



types: speed of NOT gate vs. AND gate?

# Why Is This Useful?

- ❖ Logic minimization: reduce complexity at gate level
  - Allows us to build smaller and faster hardware
  - Care about both # of gates, # of literals (gate inputs), # of gate levels, and types of logic gates
- ❖ Faster hardware?
  - Fewer inputs implies faster gates in some technologies
  - Fan-ins (# of gate inputs) are limited in some technologies
  - Fewer levels of gates implies reduced signal propagation delays
  - # of gates (or gate packages) influences manufacturing costs
- ★ ■ Simpler Boolean expressions → smaller transistor networks → smaller circuit delays → faster hardware 😊



# Are Logic Gates Created Equal?

❖ No!

2-Input Gate Type	# of CMOS transistors
NOT	2
AND	6
OR	6
NAND	4
NOR	4
XOR	8
XNOR	8

← simplest, but not too useful

} useful, and simpler than alternatives

❖ Can recreate all other gates using only NAND or only NOR gates

- Called “universal” gates
- e.g.,  $A \text{ NAND } A = \bar{A}$ ,  $B \text{ NOR } B = \bar{B}$
- DeMorgan's Law helps us here!



x	y	NAND
→ 0	0	1
0	1	1
1	0	1
→ 1	1	0

0 → 1  
1 → 0

# DeMorgan's Law

❖  $\overline{X + Y} = \bar{X} \cdot \bar{Y}$

❖  $\overline{X \cdot Y} = \bar{X} + \bar{Y}$

X	Y	$\bar{X}$	$\bar{Y}$	NOR		NAND	
				$\overline{X + Y}$	$\bar{X} \cdot \bar{Y}$	$\overline{X \cdot Y}$	$\bar{X} + \bar{Y}$
0	0	1	1	1	1	1	1
0	1	1	0	0	0	1	1
1	0	0	1	0	0	1	1
1	1	0	0	0	0	0	0

❖ In Boolean Algebra, converts between AND-OR and OR-AND expressions

▪  $Z = \bar{A}\bar{B}C + \bar{A}B\bar{C} + A\bar{B}\bar{C}$

▪  $\bar{Z} = (A + B + \bar{C}) \cdot (A + \bar{B} + \bar{C}) \cdot (\bar{A} + B + \bar{C})$

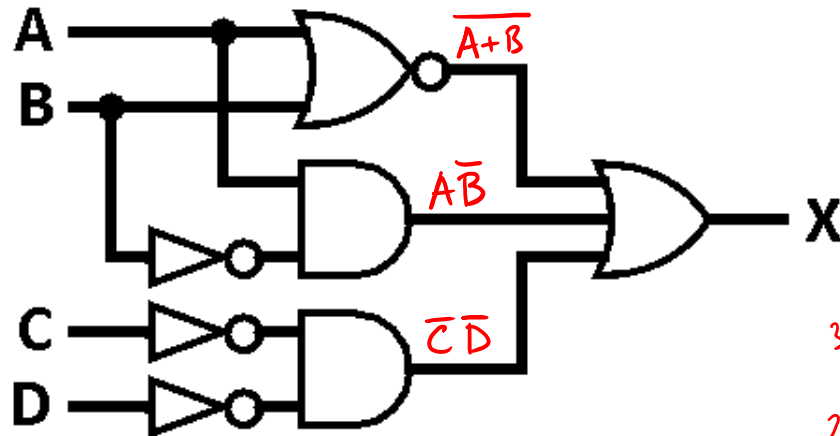
❖ At gate level, can convert from AND/OR to NAND/NOR gates

▪ "Flip" all input/output bubbles and "switch" gate 



# DeMorgan's Law Practice Problem

❖ Simplify the following diagram:



$$X = \overline{A + B} + A\overline{B} + \overline{C}D$$

*DeMorgan's*

$$X = \overline{A}\overline{B} + A\overline{B} + \overline{C}D$$

*5 gates*

$$X = \overline{B} + \overline{C}D$$

*3-4 gates*

$$X = \overline{B} + \overline{C + D}$$

*DeMorgan's*

$$X = \overline{B(C + D)}$$

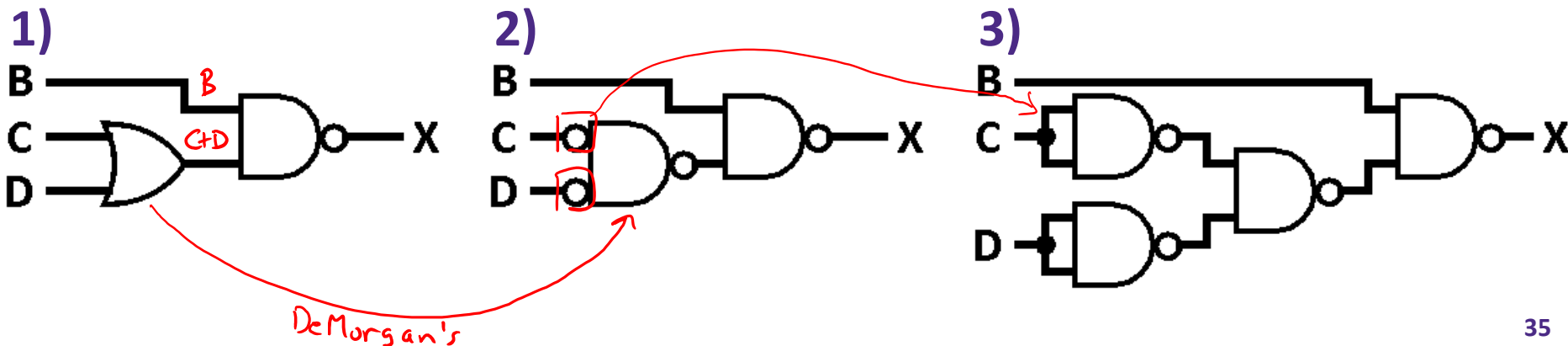
*2-3 gates*

*let E = C+D, so*

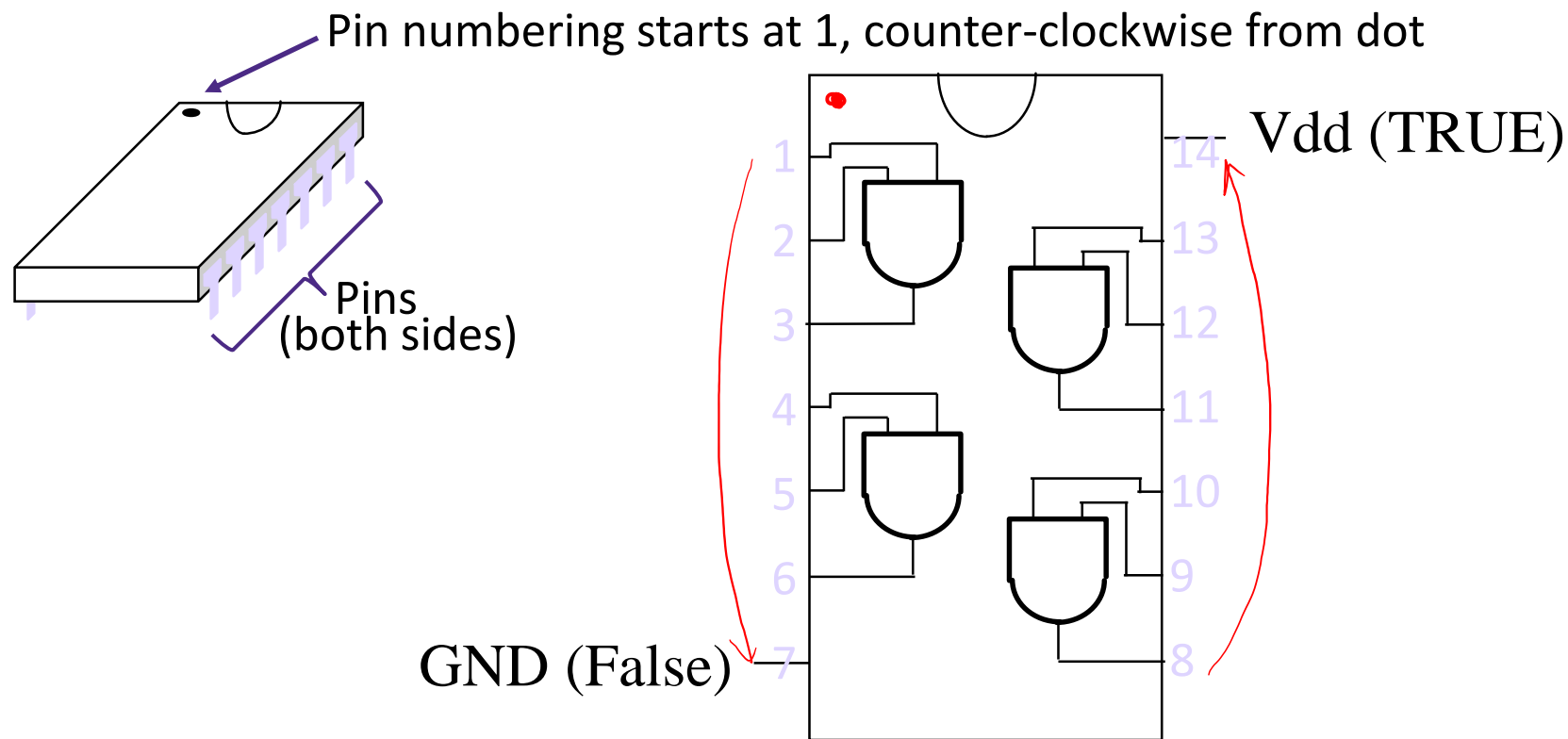
$$X = \overline{B + E}$$

$$X = \overline{BE}$$

❖ Then implement with only NAND gates:



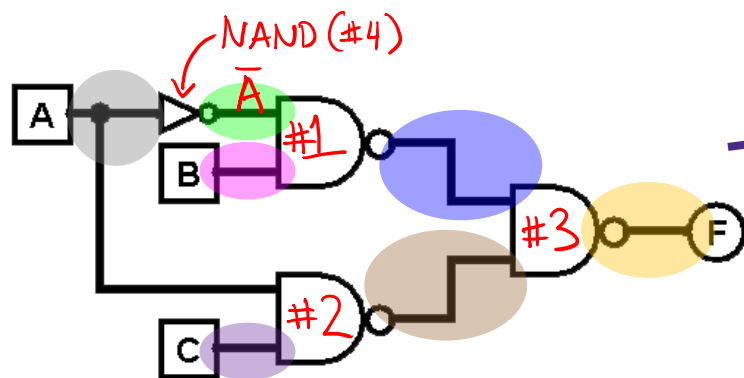
# Transistor-Transistor Logic (TTL) Packages



- ❖ Diagrams like these and other useful/helpful information can be found on part **data sheets**
  - It's really useful to learn how to read these

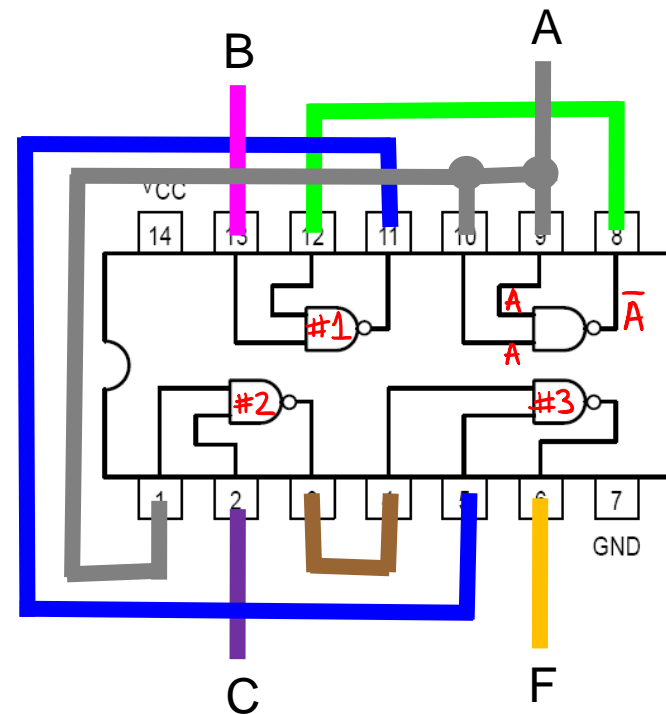
# Mapping truth tables to logic gates

- ❖ Given a truth table:
  - 1) Write the Boolean expression
  - 2) Minimize the Boolean expression
  - 3) Draw as gates
  - 4) Map to available gates
  - 5) Determine # of packages and their connections

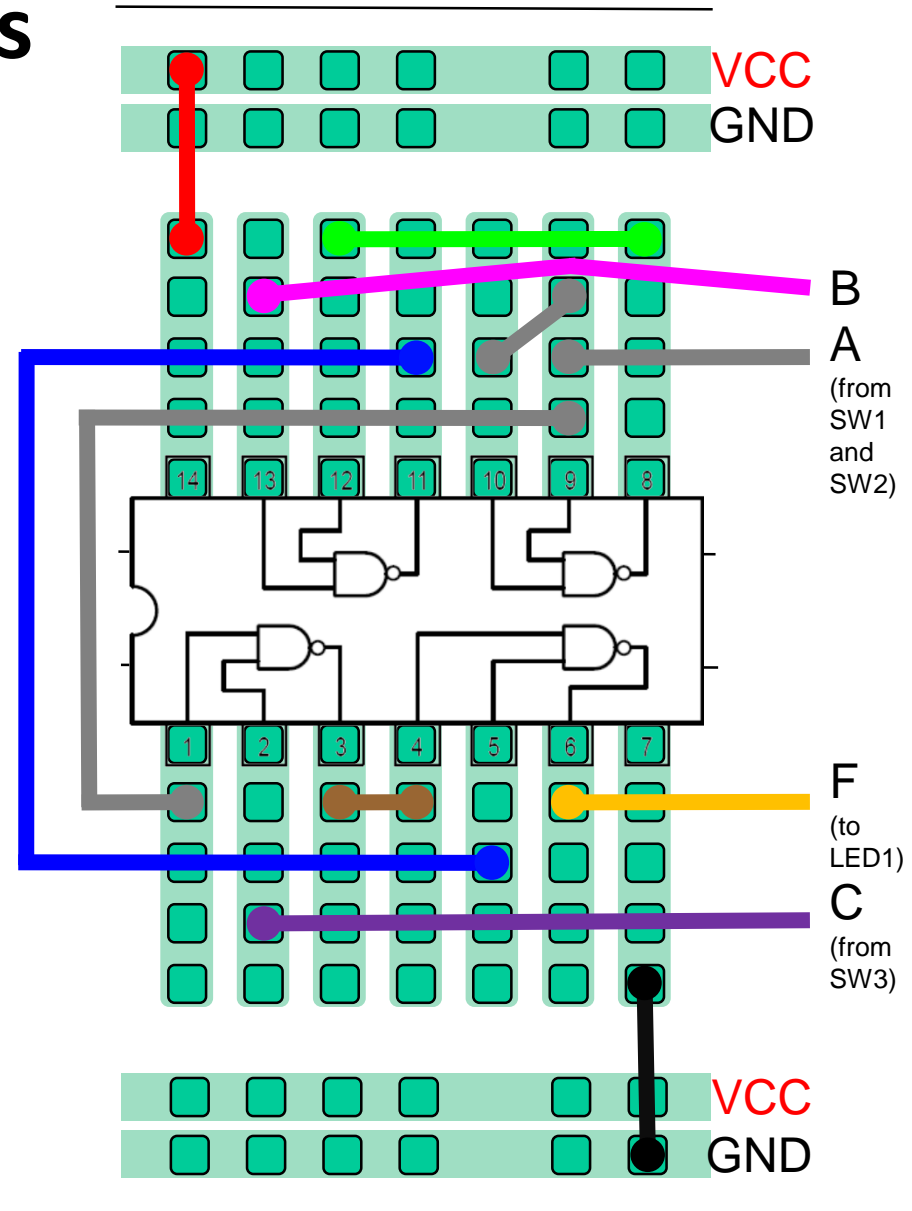
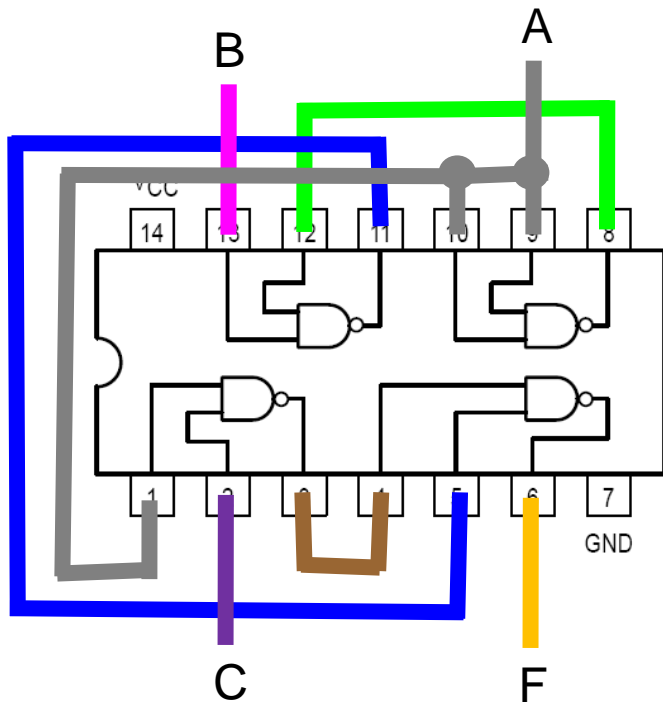


7 nets (wires) in this design

(4) →



# Breadboarding circuits



# Summary

- ❖ Digital systems are constructed from Combinational and Sequential Logic
- ❖ Logic minimization to create smaller and faster hardware
- ❖ Gates come in TTL packages that require careful wiring

