## CSE 401 - Compilers Section 2

1/24/2013<br>12:30 - MEB 238 1:30 - EE 037

## Regex Exercise

strings of 0s and 1 s such that every sequence of two 1s must be preceded by at least two consecutive Os and followed by at least three

Anyone think about this more?
Have a regex?
Have a DFA/NFA?
Think it's impossible?

## Regex Exercise

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A Key Observation: The validity of the next character depends on at most the four preceding characters

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Suggests that we can build a DFA

- States encode last characters seen


## Regex Exercise

strings of 0s and 1 s such that every sequence of two 1 s must be preceded by at least two consecutive Os and followed by at least three

| seen can see | seen can see |  |  |
| ---: | :--- | ---: | :--- |
| $\wedge$ | 0,1 | $\wedge 1$ | 0 |
| $\wedge 0$ | 0,1 | $\wedge 01$ | 0 |
| 00 | 0,1 | 101 | 0 |
| 10 | 0,1 | 0011 | 0 |
| 001 | 0,1 | 110 | 0 |
|  |  | 1100 | 0 |

## Regex Exercise

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Regular languages: RE <-> NFA <-> DFA
We've seen RE -> NFA -> DFA
DFA -> NFA is trivial
NFA -> RE can be done algorithmically via...

# Generalized Nondeterministic Finite Automaton (GNFA) 

An NFA but:

- One start state
- One accept state
- REs instead of single characters on its edges

NFA -> GNFA:

- Add super-start and super-accept states

GNFA -> RE:

- Remove states one at a time, fixing edges


## Regex Exercise

strings of 0s and 1s such that every sequence of two 1s must be preceded by at least two consecutive Os and followed by at least three

$$
\left((1 ? 0)^{*}\left(1 ?(00110)^{*} 00\right) ?\right)^{*} 1 ?
$$

## Believe me? Questions?

## Project Clarifications

Longest match examples:

- "true;" -> TRUE SEMICOLON
- "truethat;" -> ID(truethat) SEMICOLON
- "verytrue;" -> ID(verytrue) SEMICOLON
- "true that;" -> TRUE ID(that) SEMICOLON

JFLEX tries to all match REs at once

Another case:

- "45true" -> INT(45) TRUE


## Project Questions?

## Parser Ambiguities

```
expr ::= expr + expr | expr - expr |
    expr * expr | expr / expr |
    int
int : := 0 | 1 | | 2 | 3 | 4 | | 5 | | 6 | | 7 | 8 | 9 | | 0
```

1. Find an ambiguous parse
2. Fix the grammar
3. Support parenthesis

## Parser Ambiguities

```
expr ::= expr + term | expr - term | term
term ::= term * factor | term / factor | factor
factor ::= int | ( expr )
int ::= 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 0
```


## Definition Review

Sentential Form: $\alpha$ in $S=>^{*} \alpha$

S =>* ( term * factor ) =>* (2 * 3)

Handle: A position in $\alpha$ and a production that we can "undo"
term : := term * factor at position 4

## Shift-Reduce Exercise

```
expr ::= expr + term | expr - term | term
term ::= term * factor | term / factor | factor
factor ::= int | ( expr )
int ::= 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 0
```

Stack
\$
\$S

Input
$1+2$ * $3 \$$
\$

Action (shift or reduce) shift
accept

## Regular or Context-Free?

1. $L=\left\{0^{n} 1^{n} \mid n>=0\right\}$
2. $L=\left\{0^{n} 1^{m} \mid n>=0, m>n\right\}$
3. $L=\left\{w \mid \#_{0}(w)==\#_{1}(w)\right\}$
4. $L=\left\{w \mid \#_{01}(w)==\#_{10}(w)\right\}$
5. Balanced parenthesis?

Generating regex / DFA / grammar?

## Regular or Context-Free?

1. $L=\left\{0^{n} 1^{n} \mid n>=0\right\} \quad C F$
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## Questions?

Go get a job!

