Query Processing on Clusters

Communication-Cost Model Multiway Joins Recursion

Environment

Computing cluster with distributed file system.

- E.g., GFS, HDFS.
- Map-reduce implementation.
 - E.g., Hadoop.
 - Or extension to general acyclic workflow, as in Clustera, Hyracks.

Communication Cost

- Assumption: efficiency of an algorithm is tied to the sum of the input sizes to all tasks.
- Justification:
 - 1. Typically simple, main-memory operation at each task, e.g., hash-join.
 - 2. Large outputs are either input to another task or aggregated in final result.

Multiway Join

 Afrati/Ullman EDBT-2010.
Key idea: sometimes, a cascade of 2way joins, each implemented by a mapreduce stage, is less efficient than a single multiway join.

 Intuition: if intermediate result is large, communicating it costs more than replicating tuples of the arguments.

Example: Star-Join



Implementing a Star Join

 Use the A's as a hash key, so the Map tasks hash each fact tuple to one Reduce task corresponding to a bucket.

- Optimization: each A_i is hashed a number of ways inversely proportional to the size of the dimension table (A_i, B_i).
- A tuple of dimension table *i* is sent to all Reduce tasks corresponding to its value of A_i (and any buckets for the other A's).

Example: Star Join

Four dimension tables of equal size. 256 Reduce tasks. \bullet Hash each A_i to 4 buckets. \bullet Tuple (a,b) of dimension table (A₂, B₂) is sent to 64 Reduce tasks corresponding to hash values (*,h(a),*,*).

Aster Data Approach

- Hash and distribute the fact table permanently.
- Replicate the dimension tables as for a join of all relations.
- But they patented a strange approach that is data-dependent.
- We give optimal partitioning independent of data.

Problem with Recursion

 Map-reduce works because task failures can be handled by restarting only the failed task(s).

 Why possible? Because each task delivers output only at the end.

 But recursive tasks must make outputs and then process more input.

Solutions

- 1. HaLoop: use iterated map-reduce, with attention to avoiding redistribution of intermediate results.
- 2. Pregel: use recursive tasks but checkpoint after every few rounds.
 - Rollback on failure, but not too far.

Solutions – (2)

 Work of Afrati, Vinayak Borkar, Mike Carey, Alkis Polyzotis, Ullman.
If operations are idempotent (e.g., Datalog recursions computing sets), then a recursive task can be restarted anyway.

Solutions – (3)

 In general (not idempotent) case, files passed between tasks are replicated anyway.

 Master controller can restart a task, provide previous input files, but throw away output files previously delivered.