Latency-Tolerant Runtime System for Large Graph Computations

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Example: parallel graph traversal

```c
explore (Person p) {
  for each f in p.friends {
    if (atomic_increment(&f.visited) == 0) {
      /* f not yet visited */
      explore (f)
    }
  }
}
```

For large graphs, this is almost always a cache miss

Modern processors depend on locality for performance

Many important problems have little locality, leading to inefficient execution at scale

Example: Social networks

Multithreading helps hide latencies that lead to inefficient execution

Example: Madduri et al used special-purpose latency-tolerant hardware (500MHz Cray XMT) to do a centrality analysis of the IMDb movie actor database (1.54M vertices, 78M edges)

Result: 4.75x speedup over 2.4GHz Xeon system

Can we get similar speedups with commodity hardware?

Key Components

Lightweight synchronization

Problem domain has high communication to work ratio

Approach: associate a lock with every word; provide operations to manipulate data and lock atomically

Fast Context Switching

Enable memory concurrency by overlapping computation with memory requests from multiple threads

Approach: switch at user level and save minimal state schedule based on memory request completion

Memory Concurrency

Support enough outstanding memory operations to saturate processor memory bandwidth

Approach: prefetch before context switch; use coprocessor to extend prefetch depth

Implementation Tradeoffs

Software-only

- methods to provide lightweight synchronization:
  - non-standard datatype
  - multiple atomic operations per sync op

Hardware/Software

- coprocessor can provide synchronization ~ cost of remote memory reference
- coprocessor can help with prefetch and reschedule, or support more remote outstanding memory ops

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